

PIC24FJ1024GA610/GB610 Family Flash Programming Specification

1.0 DEVICE OVERVIEW

This document defines the programming specification for the PIC24FJ1024GA610/GB610 family of 16-bit microcontrollers with dual partition Flash functionality. This programming specification is required only for those developing programming support for the following devices:

- PIC24FJ1024GB610
- PIC24FJ512GB610
- PIC24FJ256GB610
- PIC24FJ128GB610
- PIC24FJ1024GB606
- PIC24FJ512GB606
- PIC24FJ256GB606
- PIC24FJ128GB606

- PIC24FJ1024GA610
- PIC24FJ512GA610
- PIC24FJ256GA610
- PIC24FJ128GA610
- PIC24FJ1024GA606
- PIC24FJ512GA606
- PIC24FJ256GA606
- PIC24FJ128GA606

Topics covered include:

- Section 1.0 "Device Overview"
- Section 2.0 "Programming Overview"
- Section 3.0 "Device Programming ICSP"
- Section 4.0 "Device Programming Enhanced ICSP"
- Section 5.0 "Programming the Programming Executive to Memory"
- Section 6.0 "The Programming Executive"
- Section 7.0 "Dual Partition Flash Programming Considerations"
- Section 8.0 "Device ID"
- Section 9.0 "Checksum Computation"
- Section 10.0 "AC/DC Characteristics and Timing Requirements"

2.0 PROGRAMMING OVERVIEW

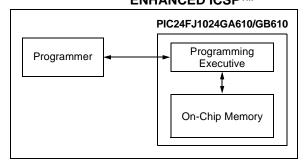
There are two methods of programming that are discussed in this programming specification:

- In-Circuit Serial Programming[™] (ICSP[™])
- · Enhanced In-Circuit Serial Programming

The ICSP programming method is the most direct method to program the device; however, it is also the slower of the two methods. It provides native, low-level programming capability to erase, program and verify the device.

The Enhanced ICSP protocol uses a faster method that takes advantage of the Programming Executive (PE), as illustrated in Figure 2-1. The PE provides all the necessary functionality to erase, program and verify the chip through a small command set. The command set allows the programmer to program a PIC24FJ1024GA610/GB610 family device without dealing with the low-level programming protocols.

FIGURE 2-1: PROGRAMMING SYSTEM OVERVIEW FOR ENHANCED ICSP™



This programming specification is divided into two major sections that describe the programming methods independently. Section 3.0 "Device Programming – ICSP" describes the ICSP method. Section 4.0 "Device Programming – Enhanced ICSP" describes the Enhanced ICSP method.

2.1 Required Connections

These devices require specific connections for programming to take place. These connections include power, VCAP, MCLR and one programming pair (PGEDx/PGECx). Table 2-1 describes these connections (refer to the specific device data sheet for pin descriptions and power connection requirements).

2.2 Power Requirements

All PIC24FJ1024GA610/GB610 family devices power their core digital logic at a nominal 1.8V. To simplify system design, all devices in the PIC24FJ1024GA610/GB610 family incorporate an on-chip regulator that allows the device to run its core logic from VDD.

The regulator provides power to the core from the other VDD pins. A low-ESR capacitor (such as ceramic or tantalum) must be connected to the VCAP pin (see Table 2-1 and Figure 2-2). This helps to maintain the stability of the regulator. The specifications for core voltage and capacitance are listed in Section 10.0 "AC/DC Characteristics and Timing Requirements".

FIGURE 2-2: CONNECTIONS FOR THE ON-CHIP REGULATOR

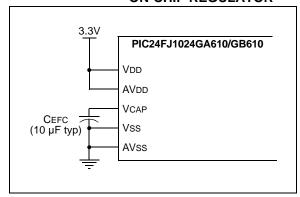


TABLE 2-1: PINS USED DURING PROGRAMMING

Pin Name Pin		Pin Description	
MCLR	I	Programming Enable	
VDD and AVDD ⁽¹⁾	Р	Power Supply ⁽¹⁾	
Vss and AVss ⁽¹⁾	Р	Ground ⁽¹⁾	
VCAP	Р	On-Chip Voltage Regulator Filter Capacitor	
PGECx	I	Programming Pin Pair: Serial Clock	
PGEDx	I/O	Programming Pin Pair: Serial Data	

Legend: I = Input O = Output P = Power

Note 1: All power supply and ground pins must be connected, including AVDD and AVSS.

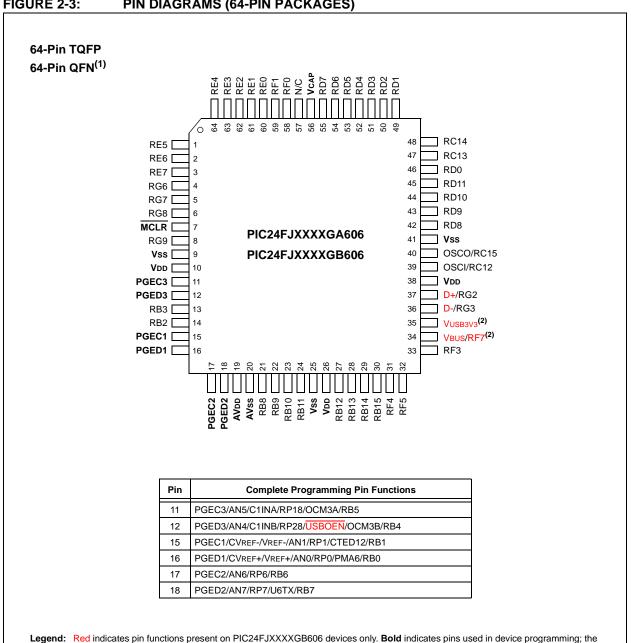
2.3 **Pin Diagrams**

Figure 2-3 through Figure 2-5 show the pin diagrams for the PIC24FJ1024GA610/GB610 family. The pins that are required for programming are listed in Table 2-1 and are indicated in bold text in the figures. Refer to the appropriate device data sheet for complete pin descriptions.

PGECx AND PGEDx PIN PAIRS 2.3.1

All devices in the PIC24FJ1024GA610/GB610 family have three separate pairs of programming pins, labeled as PGEC1/PGED1, PGEC2/PGED2 and PGEC3/ PGED3. Any one of these pin pairs may be used for device programming by either ICSP or Enhanced ICSP. Unlike voltage supply and ground pins, it is not necessary to connect all three pin pairs to program the device. However, the programming method must use both pins of the same pair.





complete list of functions associated with programming/emulation pins is shown in the accompanying table.

Note 1: It is recommended to connect the metal pad on the bottom of the 64-pin QFN package to Vss.

RF2 and RF6 are on Pins 34 and 35 of PIC24FJXXXXGA606 devices.

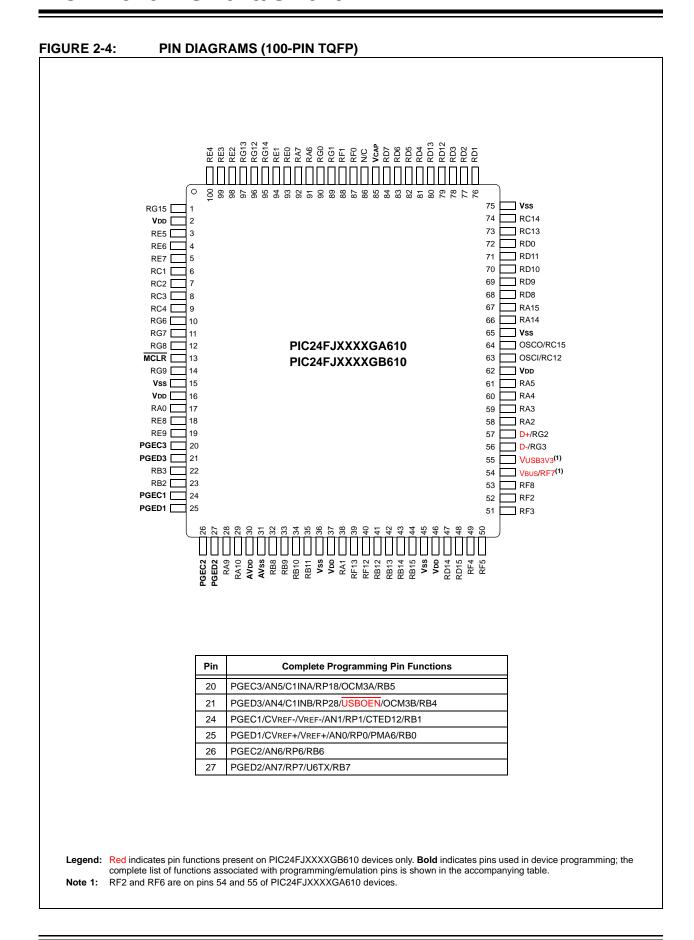


FIGURE 2-5: F	IN DIAGRAMS	(121-PIN BGA)
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	1	2	3	4	5	6	7	8	9	10	11
	0	0	0	0	0	0	0	0	0	0	0
Α	RE4	RE3	RG13	RE0	RG0	RF1	N/C	N/C	RD12	RD2	RD1
В	N/C	RG15	RE2	RE1	RA7	RF0	VCAP	RD5	RD3	Vss	RC14
С	RE6	VDD	RG12	RG14	RA6	N/C	RD7	RD4	N/C	RC13	RD11
D	RC1	C) RE7	C RE5	N/C	N/C	N/C	RD6	C RD13	O RD0	N/C	C RD10
E	C RC4	C RC3	C RG6	C RC2	O N/C	O RG1	O N/C	RA15	C RD8	C RD9	C RA14
_	0	0	\bigcirc	0	\bigcirc	0	0	\circ	\circ	\bigcirc	0
F	MCLR	RG8	RG9	RG7	Vss	N/C	N/C	VDD	RC12	Vss	RC15
G	RE8	RE9	RA0	N/C	VDD	Vss	Vss	N/C	RA5	RA3	RA4
Н	PGEC3	PGED3	N/C	N/C	N/C	VDD	N/C	VBUS/RF7 ⁽¹⁾	_	_	RA2
J	RB3	C) RB2	PGED2	AVDD	C) RB11	C) RA1	RB12	N/C	N/C	C RF8	D-/RG3
K	PGEC1	O PGED1	C RA10	O RB8	O N/C	O RF12	C RB14	○ VDD	C RD15	C RF3	O RF2
	0	\bigcirc	\bigcirc	\bigcirc	\circ	\bigcirc	\bigcirc	\bigcirc	\bigcirc	\bigcirc	\circ
L	PGEC2	RA9	AVss	RB9	RB10	RF13	RB13	RB15	RD14	RF4	RF5

Pin	Complete Programming Pin Functions
H1	PGEC3/AN5/C1INA/RP18/OCM3A/RB5
H2	PGED3/AN4/C1INB/RP28/USBOEN/OCM3B/RB4
K1	PGEC1/CVREF-/VREF-/AN1/RP1/CTED12/RB1
K2	PGED1/CVREF+/VREF+/AN0/RP0/PMA6/RB0
26	PGEC2/AN6/RP6/RB6
27	PGED2/AN7/RP7/U6TX/RB7

Legend: Red indicates pin functions present on PIC24FJXXXXGB610 (BGA) devices only. **Bold** indicates pins used in device programming; the complete list of functions associated with programming/emulation pins is shown in the accompanying table.

Note 1: RF2 and RF6 are on balls H8 and H9 of PIC24FJXXXXGA610 devices.

2.4 Program Memory Write/Erase Requirements

The program Flash memory has a specific write/erase requirement that must be adhered to for proper device operation. The rule is that any given word in memory must not be written without first erasing the page in which it is located. Thus, the easiest way to conform to this rule is to write all the data in a programming block within one write cycle. The programming methods specified in this document comply with this requirement.

2.5 Memory Map

The Flash program memory can be set as 'Single Partition' or 'Dual Partition' Flash mode. Single Partition has one contiguous Flash block with designated areas for Reset vectors, Configuration Words, OTP and Device ID Words. Dual Partition Flash mode splits the Flash memory into two halves, beginning at address, 0x40000. Each half has these same areas, at different addresses, to allow two sets of possible device behavior at boot time, based on which partition is used. These partitions are referred to as the 'Active' or the 'Inactive' Partition.

The Single Partition memory map extends from 000000h to FFFFFEh. Code storage is located at the base of the memory map. The last locations of implemented program memory are reserved for the device Configuration bits.

Locations, 800100h through 800FFEh, are reserved for executive code memory. This region stores the PE and the debugging executive, which are used for device programming. This region of memory cannot be used to store user code. See Section 6.0 "The Programming Executive" for more information.

Locations, 801700h through 8017FEh, are reserved for the customer OTP data. This area can be used for storing product information, such as serial numbers, system manufacturing dates, manufacturing lot numbers and other application-specific information. It is described in Section 2.6.3 "Customer OTP Memory".

Locations, FF0000h and FF0002h, are reserved for the Device ID Word registers. These bits can be used by the programmer to identify which device type is being programmed. They are described in **Section 8.0** "Device ID". The Device ID registers read out normally, even after code protection is applied.

Figure 2-6 and Figure 2-7 show the generic memory maps for the devices described in this specification. See the "Memory Organization" chapter in the specific device data sheet for exact memory addresses.

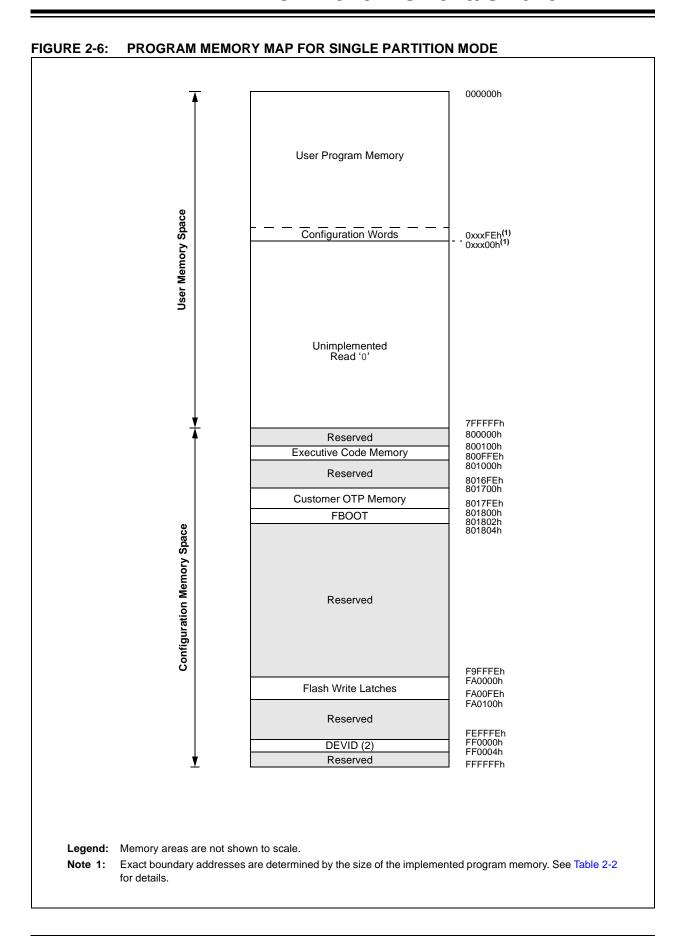
Table 2-2 lists the code memory size, the size of the erase blocks and the number of erase blocks present in each device variant.

TABLE 2-2:	PROGRAM MEMORY SIZES AND BOUN	TUVDIEG

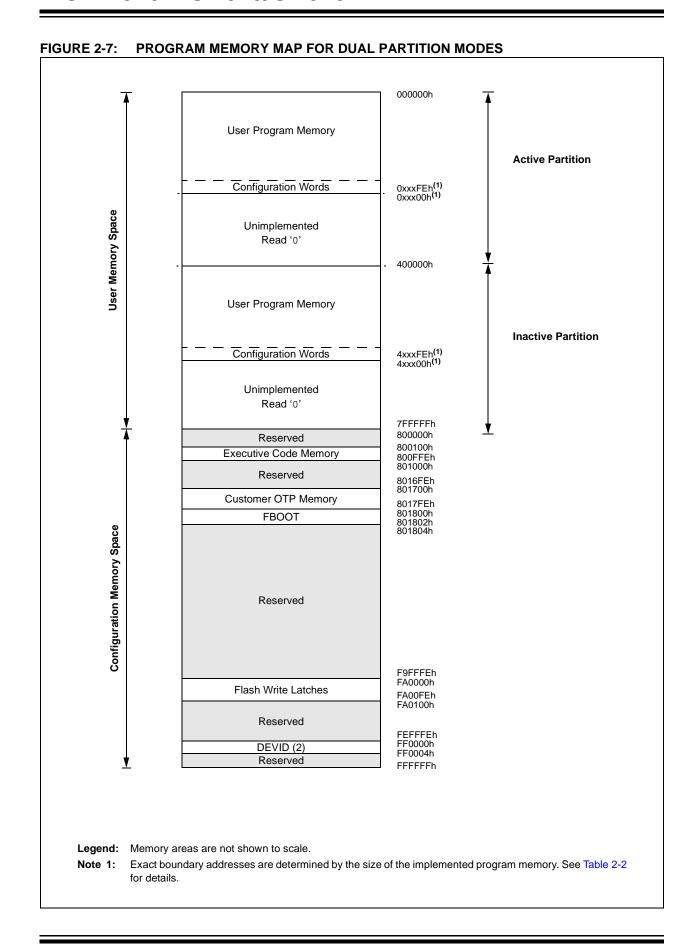
Program Memory Upper Boundary (Instruction Words)					
Device	Single Partition Mode	Dual Partit	Write Blocks ⁽¹⁾	Erase Blocks ⁽¹⁾	
	Single Partition Mode	Active Partition	Inactive Partition		
PIC24FJ1024GX6XX	0ABFFEh (352,256 x 24)	055FFEh (176,128K x 24)	455FFEh (176,128 x 24)	5504	688
PIC24FJ512GX6XX	055FFEh (176,128 x 24)	02AFFEh (88,064 x 24)	42AFFEh (88,064 x 24)	2752	344
PIC24FJ256GX6XX	02AFFEh (88,064 x 24)	0157FEh (44,032 x 24)	4157FEh (44,032 x 24)	1376	172
PIC24FJ128GX6XX	15FFEh (45,056 x 24) ⁽²⁾	0AFFEh (22,528 x 24) ⁽²⁾	40AFFEh (22,528 x 24) ⁽²⁾	704	88

Note 1: 1 Write Block = 128 Instruction Words; 1 Erase Block = 1024 Instruction Words.

^{2:} To maintain integer page sizes, the 128K variant is slightly larger than half of the 256K variant.



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2.6 Configuration Bits

2.6.1 OVERVIEW

The Configuration bits are stored in locations based on if the part is in Single or Dual Partition Flash mode. These bits can be set or cleared to select various device configurations. There are two types of Configuration bits: system operation bits and code-protect bits. The system operation bits determine the power-on settings for system-level components, such as the oscillator and the Watchdog Timer. The code-protect bits prevent program memory from being read and written.

Table 2-3 lists the Configuration register address range for each device in Single and Dual Partition Flash modes. Table 2-4 lists all of the Configuration bits found in PIC24FJ1024GA610/GB610 family devices, as well as their Configuration register locations. Refer to the "Special Features" chapter in the specific device data sheet for the full Configuration register description for a specific device.

2.6.2 CODE-PROTECT CONFIGURATION BITS

The device implements an intermediate security feature defined by the FSEC register. The Boot Segment (BS) is the higher privileged segment and the General Segment (GS) is the lower privileged segment. The total user code memory can be split into BS or GS. The size of the segments is determined by BSLIM<12:0>. The relative location of the segments within user space does not change, such that BS (if present) occupies the memory area just after the Interrupt Vector Table (IVT) and the GS occupies the space just after the BS (or if the Alternate IVT is enabled, just after it).

The Configuration Segment (CS) is a small segment (less than a page, typically just one row) within user Flash address space. It contains all user configuration data that is loaded by the NVM controller during the Reset seguence.

TABLE 2-3: CONFIGURATION WORD ADDRESSES

Single Partition Mode					
PIC24FJ1024GX6XX	PIC24FJ512GX6XX	PIC24FJ256GX6XX	PIC24FJ128GX6XX		
0ABF00h	055F00h	02AF00h	015F00h		
0ABF10h	055F10h	02AF10h	015F10h		
0ABF14h	055F14h	02AF14h	015F14h		
0ABF18h	055F18h	02AF18h	015F18h		
0ABF1Ch	055F1Ch	02AF1Ch	015F1Ch		
0ABF20h	055F20h	02AF20h	015F20h		
0ABF24h	055F24h	02AF24h	015F24h		
0ABF28h	0ABF28h 055F28h		015F28h		
0ABF2Ch	055F2Ch	02AF2Ch	015F2Ch		
801800h					
	Dual Partiti	on Modes ⁽¹⁾			
055F00h/455F00h	02AF00h/42AF00h	015700h/415700h	00AF00h/40AF00h		
055F10h/455F10h	02AF10h/42AF10h	015710h/415710h	00AF10h/40AF10h		
055F14h/455F14h	02AF14h/42AF14h	015714h/ 415714h	00AF14h/40AF14h		
055F18h/455F18h	02AF18h/42AF18h	015718h/415718h	00AF18h/40AF18h		
055F1Ch/455F1Ch	02AF1Ch/42AF1Ch	01571Ch/41571Ch	00AF1Ch/40AF1Ch		
055F20h/455F20h	02AF20h/42AF20h	015720h/415720h	00AF20h/40AF20h		
055F24h/ 455F24h	02AF24h/42AF24h	015724h/415724h	00AF24h/40AF24h		
055F28h/455F28h	02AF28h/42AF28h	015728h/415728h	00AF28h/40AF28h		
055F2Ch/455F2Ch	02AF2Ch/42AF2Ch	01572Ch/41572Ch	00AF2Ch/40AF2Ch		
055FFCh/455FFCh	02AFFCh/42AFFCh	0157FCh/4157FCh	00AFFCh/40AFFCh		
801800h					
	0ABF00h 0ABF10h 0ABF14h 0ABF18h 0ABF1Ch 0ABF20h 0ABF24h 0ABF28h 0ABF2Ch 055F00h/455F00h 055F10h/455F10h 055F14h/455F14h 055F120h/455F1Ch 055F20h/455F20h 055F24h/455F24h 055F28h/455F2Ch 055F5Ch/455FCCh	PIC24FJ1024GX6XX PIC24FJ512GX6XX 0ABF00h 055F00h 0ABF10h 055F10h 0ABF14h 055F14h 0ABF18h 055F18h 0ABF1Ch 055F1Ch 0ABF20h 055F20h 0ABF24h 055F24h 0ABF28h 055F28h 0ABF2Ch 055F2Ch Bual Partiti 055F00h/455F00h 02AF00h/42AF00h 055F10h/455F10h 02AF10h/42AF10h 055F14h/455F14h 02AF14h/42AF14h 055F18h/455F16h 02AF1Ch/42AF1Ch 055F20h/455F20h 02AF20h/42AF20h 055F24h/455F24h 02AF24h/42AF24h 055F28h/455F28h 02AF28h/42AF28h 055F2Ch/455F2Ch 02AF2Ch/42AF2Ch 055FFCh/455FFCh 02AFFCh/42AFFCh	PIC24FJ1024GX6XX PIC24FJ512GX6XX PIC24FJ256GX6XX 0ABF00h 055F00h 02AF00h 0ABF10h 055F10h 02AF10h 0ABF14h 055F14h 02AF14h 0ABF18h 055F18h 02AF18h 0ABF1Ch 055F1Ch 02AF1Ch 0ABF20h 055F20h 02AF20h 0ABF24h 055F24h 02AF24h 0ABF28h 055F28h 02AF28h 0ABF2Ch 055F2Ch 02AF2Ch B01800h Dual Partition Modes ⁽¹⁾ Dual Partition Modes ⁽¹⁾ 055F00h/455F00h 02AF00h/42AF00h 015700h/415700h 055F10h/455F10h 02AF10h/42AF10h 015710h/415710h 055F14h/455F18h 02AF18h/42AF18h 015718h/415718h 055F1Ch/455F1Ch 02AF10h/42AF1Ch 015710h/415710h 055F20h/455F20h 02AF20h/42AF20h 015720h/415720h 055F24h/ 455F24h 02AF28h/42AF28h 015724h/415728h 055F2Ch/455F2Ch 02AF28h/42AF2Ch 01572Ch/4157FCh 055FFCh/455FFCh		

- Note 1: Addresses shown for Dual Partition modes are for the Active/Inactive Partitions, respectively.
 - **2:** Changes to these Inactive Partition Configuration Words affect how the Active Partition accesses the Inactive Partition.
 - 3: FBTSEQ is a 24-bit Configuration Word that uses all three bytes of the program memory width.

TABLE 2-4: CONFIGURATION BITS FOR PIC24FJ1024GA610/GB610 FAMILY DEVICES

Bit Field	Register	Description		
AIVTDIS	FSEC<15>	Alternate Interrupt Vector Table Disable bit 1 = Disables AIVT; INTCON2<8> (AIVTEN) bit is not available 0 = Enables AIVT; INTCON2<8> (AIVTEN) bit is available		
ALTCMPI	FDEVOPT1<1>	Alternate Comparator Input Enable bit 1 = C1INC, C2INC and C3INC are on their standard pin locations 0 = C1INC, C2INC and C3INC are on RG9		
ALTVREF	FDEVOPT1<14>	Alternate Voltage Reference Location Enable bit (100-pin and 121-pin devices only) 1 = VREF+ and CVREF+ on RA10, VREF- and CVREF- on RA9 0 = VREF+ and CVREF+ on RB0, VREF- and CVREF- on RB1		
BOREN<1:0>	FPOR<1:0>	Brown-out Reset Enable bits 11 = Brown-out Reset is enabled in hardware; SBOREN bit is disabled 10 = Brown-out Reset is enabled only while device is active and disabled in Sleep; SBOREN bit is disabled 01 = Brown-out Reset is controlled with the SBOREN bit setting 00 = Brown-out Reset is disabled in hardware; SBOREN bit is disabled		
BSLIM<12:0>	FBSLIM<12:0>	Boot Segment Code Flash Page Address Limit bits Specifies the last Boot Segment page + 1 (the first page of the General Segment). As this value is inverted, the default unprogrammed value (1FFFh) sets the Boot Segment size to 0.		
BSS<1:0>	FSEC<2:1>	Boot Segment Code Protection Level bits 11 = No protection (other than BWRP) 10 = Standard security 0x = High security		
BSEN	FSEC<3>	Boot Segment Control bit 1 = No Boot Segment is enabled 0 = Boot Segment size is determined by BSLIM<12:0>		
BTMODE<1:0>	FBOOT<1:0>	Device Partition Mode Configuration bits 11 = Single Partition Flash mode 10 = Dual Partition Flash mode 01 = Protected Dual Partition Flash mode (Partition 1 is write-protected when inactive) 00 = Reserved; do not use		
BSEQ<11:0>	FBTSEQ<11:0>	Boot Sequence Number bits (Dual Partition Flash modes only) Relative value defining which partition will be active after a device Reset; the partition containing a lower boot number will be active.		
BTSWP	FICD<15>	BOOTSWP Instruction Enable bit 1 = BOOTSWP instruction is disabled 0 = BOOTSWP instruction is enabled		
BWRP	FSEC<0>	Boot Segment Program Write Protection bit 1 = Boot Segment can be written 0 = Boot Segment is write-protected		
CSS<2:0>	FSEC<11:9>	Configuration Segment Code Protection Level bits 111 = No protection (other than CWRP) 110 = Standard security 10x = Enhanced security 0xx = High security		
CWRP	FSEC<8>	Configuration Segment Program Write Protection bit 1 = Configuration Segment is not write-protected 0 = Configuration Segment is write-protected		

TABLE 2-4: CONFIGURATION BITS FOR PIC24FJ1024GA610/GB610 FAMILY DEVICES (CONTINUED)

Bit Field	Register	Description
DEBUG	FICD<7>	Background Debugger Enable bit
		1 = Background debugger is disabled
		0 = Background debugger is enabled
DNVPEN	FPOR<3>	Downside Voltage Protection Enable bit
		1 = Downside protection is enabled using BOR; BOR re-arms POR as needed0 = Downside protection is disabled when BOR is inactive
FCKSM<1:0>	FOSC<7:6>	Clock Switching and Monitor Selection bits
		1x = Clock switching and the Fail-Safe Clock Monitor are disabled 01 = Clock switching is enabled, Fail-Safe Clock Monitor is disabled 00 = Clock switching and the Fail-Safe Clock Monitor are enabled
FNOSC<2:0>	FOSCSEL<2:0>	Oscillator Selection bits
		111 = Oscillator with Frequency Divider (OSCFDIV) 110 = Digitally Controlled Oscillator (DCO) 101 = Low-Power RC Oscillator (LPRC) 100 = Secondary Oscillator (SOSC) 011 = Primary Oscillator with PLL (XTPLL, HSPLL, ECPLL) 010 = Primary Oscillator (XT, HS, EC) 001 = Fast RC Oscillator with PLL (FRCPLL) 000 = Fast RC Oscillator (FRC)
FWDTEN<1:0>	FWDT<6:5>	Watchdog Timer Enable bits
		11 = WDT is enabled 10 = WDT is disabled (control is placed on the SWDTEN bit) 01 = WDT is enabled only while device is active and disabled in Sleep; SWDTEN bit is disabled 00 = WDT and SWDTEN are disabled
FWPSA	FWDT<4>	Watchdog Timer Prescaler bit
		1 = WDT prescaler ratio of 1:128 0 = WDT prescaler ratio of 1:32
GCP<1:0>	FSEC<7:6>	General Segment Code Protection Level bits
		11 = No protection (other than GWRP) 10 = Standard security 0x = High security
GWRP	FSEC<5>	General Segment Program Write Protection bit
		1 = General Segment is not write-protected 0 = General Segment is write-protected
IBSEQ<11:0>	FBTSEQ<23:12>	Inverse Boot Sequence Number bits (Dual Partition Flash modes only)
		The one's complement of BSEQ<11:0>; must be calculated by the user and written into device programming.
ICS<1:0>	FICD<1:0>	ICD Communication Channel Select bits
		11 = Communicates on PGEC1/PGED1 10 = Communicates on PGEC2/PGED2 01 = Communicates on PGEC3/PGED3 00 = Reserved; do not use
IESO	FOSCSEL<7>	Two-Speed Oscillator Start-up Enable bit
		 1 = Starts up the device with FRC, then automatically switches to the user-selected oscillator when ready 0 = Starts up the device with the user-selected oscillator source
IOL1WAY	FOSC<5>	Peripheral Pin Select Configuration bit
		1 = The IOLOCK bit can be set only once (with unlock sequence). 0 = The IOLOCK bit can be set and cleared as needed (with unlock sequence)

TABLE 2-4: CONFIGURATION BITS FOR PIC24FJ1024GA610/GB610 FAMILY DEVICES (CONTINUED)

Bit Field	Register	Description
JTAGEN	FICD<5>	JTAG Enable bit (not used for checksum calculation)
		1 = JTAG port is enabled
		0 = JTAG port is disabled
LPCFG	FPOR<2>	Low-Power Regulator Control bit
		1 = Retention feature is not available 0 = Retention feature is available and controlled by RETEN during Sleep
OSCIOFNC	FOSC<2>	CLKO Enable Configuration bit
		 1 = CLKO output signal is active on the OSCO pin (when the Primary Oscillator is disabled or configured for EC mode) 0 = CLKO output is disabled
PLLMODE<3:0>	FOSCSEL<6:3>	Frequency Multiplier Select bits
		1111 = No PLL is used (PLLEN bit is unavailable)
		1110 = 8x PLL is selected
		1101 = 6x PLL is selected 1100 = 4x PLL is selected
		0111 = 96 MHz USB PLL is selected (Input Frequency = 48 MHz)
		0110 = 96 MHz USB PLL is selected (Input Frequency = 32 MHz)
		0101 = 96 MHz USB PLL is selected (Input Frequency = 24 MHz)
		0100 = 96 MHz USB PLL is selected (Input Frequency = 20 MHz)
		0011 = 96 MHz USB PLL is selected (Input Frequency = 16 MHz) 0010 = 96 MHz USB PLL is selected (Input Frequency = 12 MHz)
		0001 = 96 MHz USB PLL is selected (Input Frequency = 8 MHz)
		0000 = 96 MHz USB PLL is selected (Input Frequency = 4 MHz)
PLLSS	FOSC<4>	PLL Secondary Selection Configuration bit
		This Configuration bit only takes effect when PLL is NOT being used by the system (i.e., selected as the system clock source).
		1 = PLL is driven by the Primary Oscillator 0 = PLL is driven by the FRC Oscillator
POSCMOD<1:0>	FOSC<1:0>	Primary Oscillator Configuration bits
		11 = Primary Oscillator is disabled
		10 = HS Oscillator mode is selected (10 MHz-2 MHz) 01 = XT Oscillator mode is selected (3.5 MHz-10 MHz)
		00 = EC Oscillator mode is selected
(Reserved)	FSIGN<15>	Reserved; always maintain as '0'.
(110001100)	1 01011 1102	
(Reserved)	FOSCSEL<9:8>	Reserved; always maintain as '0'.
SOSCHP	FDEVOPT1<3>	SOSC High-Power Enable bit (valid only when SOSCSEL = 1)
		1 = SOSC High-Power mode is enabled 0 = SOSC Low-Power mode is enabled
SOSCSEL	FOSC<3>	SOSC Selection Configuration bit
		1 = Crystal (SOSCI/SOSCO) mode 0 = Digital (CLKI) mode
TMPRPIN	FDEVOPT1<2>	Tamper Pin Enable bit
		1 = TMPR pin function is disabled
		$0 = \overline{\text{TMPR}}$ pin function is enabled

TABLE 2-4: CONFIGURATION BITS FOR PIC24FJ1024GA610/GB610 FAMILY DEVICES (CONTINUED)

Bit Field	Register	Description
WDTCLK<1:0>	FWDT<14:13>	Watchdog Timer Clock Select bits 11 = Always uses INTOSC/LPRC 10 = Uses FRC when WINDIS = 0, system clock is not INTOSC/LPRC and device is not in Sleep; otherwise, uses INTOSC/LPRC 01 = Always uses SOSC 00 = Uses peripheral clock when system clock is not INTOSC/LPRC and device is not in Sleep; otherwise, uses INTOSC/LPRC
WDTCMX	FWDT<11>	WDT Clock MUX Control bit 1 = Enables WDT clock MUX 0 = WDT clock MUX is disabled
WDTPS<3:0>	FWDT<3:0>	Watchdog Timer Postscaler bits 1111 = 1:32,768 1110 = 1:16,384 1101 = 1:8,192 1100 = 1:4,096 1011 = 1:2,048 1010 = 1:512 1000 = 1:256 0111 = 1:128 0110 = 1:64 0101 = 1:32 0100 = 1:16 0011 = 1:8 0010 = 1:4 0001 = 1:2 0000 = 1:1
WDTWIN<1:0>	FWDT<9:8>	Watchdog Timer Window Width bits 11 = WDT window is 25% of the WDT period 10 = WDT window is 37.5% of the WDT period 01 = WDT window is 50% of the WDT period 00 = WDT window is 75% of the WDT period
WINDIS	FWDT<7>	Windowed Watchdog Timer Disable bit 1 = Windowed WDT is disabled 0 = Windowed WDT is enabled

2.6.3 CUSTOMER OTP MEMORY

PIC24FJ1024GA610/GB610 family devices provide 384 bytes of One-Time-Programmable (OTP) memory, located at addresses, 801700h through 8017FEh. This memory can be used for persistent storage of application-specific information that will not be erased by reprogramming the device. This includes many types of information, such as:

- · Application checksums
- · Code revision information
- Product information
- Serial numbers
- System manufacturing dates
- · Manufacturing lot numbers

Customer OTP memory may be programmed in any mode, including user RTSP mode, but it cannot be erased. Data is not cleared by a Chip Erase.

Note: OTP memory should only be written to once or ECC errors/traps may occur.

3.0 DEVICE PROGRAMMING - ICSP

ICSP mode is a special programming protocol that allows you to read and write to device memory. The ICSP mode is the most direct method used to program the device, which is accomplished by applying control codes and instructions serially to the device, using the PGECx and PGEDx pins. ICSP mode also has the ability to read executive memory to determine if the Programming Executive (PE) is present and to write the PE to executive memory if Enhanced ICSP mode will be used.

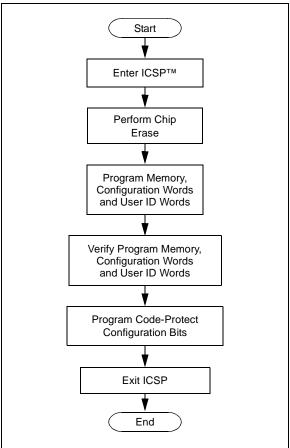
In ICSP mode, the system clock is taken from the PGECx pin, regardless of the device's oscillator Configuration bits. All instructions are shifted serially into an internal buffer, then loaded into the Instruction Register (IR) and executed. No program fetching occurs from internal memory. Instructions are fed in 24 bits at a time. PGEDx is used to shift data in, and PGECx is used as both the serial shift clock and the CPU execution clock.

- Note 1: During ICSP operation, the operating frequency of PGECx must not exceed 5 MHz.
 - **2:** ICSP mode is slower than Enhanced ICSP mode for programming.

3.1 Overview of the Programming Process

Figure 3-1 illustrates the high-level overview of the programming process. After entering ICSP mode, the first action is to Chip Erase program memory. Next, the code memory is programmed, followed by the device Configuration bits. Code memory (including the Configuration bits) is then verified to ensure that programming was successful. Then, the code-protect Configuration bits are programmed, if required.

FIGURE 3-1: HIGH-LEVEL ICSP™ PROGRAMMING FLOW



3.2 Entering ICSP Mode

As shown in Figure 3-2, entering ICSP Program/Verify mode requires three steps:

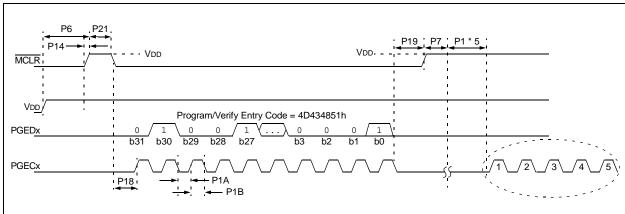
- 1. MCLR is briefly driven high, then low (P21).
- 2. A 32-bit key sequence is clocked into PGEDx. The interval of at least P18 must elapse before presenting the key sequence on PGEDx.
- 3. MCLR is held low during a specified period, P19, and then driven high.
- 4. After a P7 + 5 * P1 delay, five clock pulses must be generated on the PGECx pin.

Note: If a capacitor is present on the MCLR pin, the high time for entering ICSP mode can vary.

The key sequence is a specific 32-bit pattern, '0100 1101 0100 0011 0100 1000 0101 0001' (more easily remembered as 4D434851h in hexadecimal). The device will enter ICSP mode only if the sequence is valid. The Most Significant bit (MSb) of the most significant nibble must be shifted in first.

On successful entry, the program memory can be accessed and programmed in serial fashion.

FIGURE 3-2: ENTERING ICSP™ MODE



3.3 ICSP Operation

After entering into ICSP mode, the CPU is Idle. Execution of the CPU is governed by an internal state machine. A 4-bit control code is clocked in using PGECx and PGEDx, and this control code is used to command the CPU (see Table 3-1).

The SIX control code is used to send instructions to the CPU for execution and the REGOUT control code is used to read data out of the device through the VISI register.

TABLE 3-1: CPU CONTROL CODES IN ICSP™ MODE

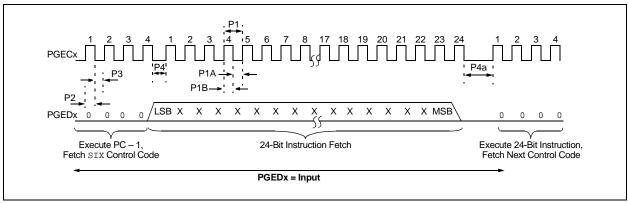
4-Bit Control Code	Mnemonic	Description
0000	SIX	Shift in 24-bit instruction and execute.
0001	REGOUT	Shift out the VISI register.
0010-1111	N/A	Reserved.

3.3.1 SIX SERIAL INSTRUCTION EXECUTION

The SIX control code allows execution of assembly instructions. When the SIX code is received, the CPU is suspended for 24 clock cycles as the instruction is then clocked into the internal buffer. Once the instruction is shifted in, the state machine allows it to be executed over the next four clock cycles. While the received instruction is executed, the state machine simultaneously shifts in the next 4-bit command (see Figure 3-3).

- Note 1: The device will latch input PGEDx data on the rising edge of PGECx. For all data transmissions, the Least Significant bit (LSb) is transmitted first.
 - 2: TBLRDH, TBLRDL, TBLWTH and TBLWTL instructions must be followed by two NOP instructions.
 - **3:** During ISCP programming, the CLKO pin will toggle. If external logic is connected to CLKO, be sure that this toggling will not effect the circuitry during the programming sequence.

FIGURE 3-3: SIX SERIAL EXECUTION



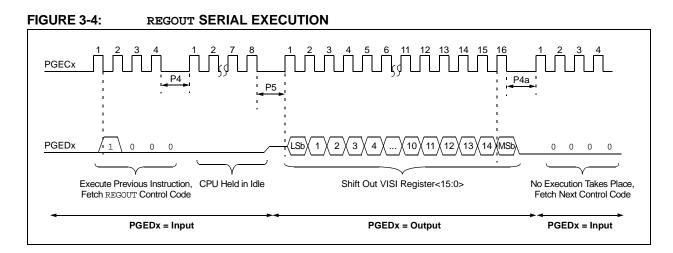
3.3.2 REGOUT SERIAL INSTRUCTION EXECUTION

The REGOUT control code allows for data to be extracted from the device in ICSP mode. It is used to clock the contents of the VISI register, out of the device, over the PGEDx pin. After the REGOUT control code is received, the CPU is held Idle for eight cycles. After these eight cycles, an additional 16 cycles are required to clock the data out (see Figure 3-4).

The REGOUT code is unique because the PGEDx pin is an input when the control code is transmitted to the device. However, after the control code is processed, the PGEDx pin becomes an output as the VISI register is shifted out.

Note:

The device will output data on the PGEDx line on the rising edge of PGECx. For all data transmissions, the Least Significant bit (LSb) is transmitted first.



3.4 Flash Memory Programming in ICSP Mode

3.4.1 PROGRAMMING OPERATIONS

Flash memory write/erase operations are controlled by the NVMCON register. Programming is performed by setting NVMCON to select the type of erase operation (Table 3-2) or write operation (Table 3-3) and initiating the programming by setting the WR control bit (NVMCON<15>).

In ICSP mode, all programming operations are self-timed. The WR control bit is cleared by hardware when the programming operation is complete. The ICSP programmer must supply enough clock pulses on the PGECx pin to complete the erase or program operation. Refer to Section 10.0 "AC/DC Characteristics and Timing Requirements" for detailed information about the maximum number of clock pulses required for erase or write operations.

TABLE 3-2: NVMCON ERASE OPERATIONS

NVMCON Value	Erase Operation
400Eh	Chip Erase (erases user memory, does not erase executive memory, Device ID or customer OTP).
4003h	Erases a page of program or executive memory.
4004h	Erases user memory and Configuration Words in the Inactive Partition (Dual Partition Flash modes only).

TABLE 3-3: NVMCON WRITE OPERATIONS

NVMCON Value	Write Operation
4001h	Double-word program operation.
4002h	Row programming operation.
4008h	The next WR command will perform a Boot mode program (writing to FBOOT) and then program the Dual Partition Signature (FSIGN) bit. Device must be reset before the newly programmed mode can take effect.

3.4.2 STARTING AND STOPPING A PROGRAMMING CYCLE

For protection against accidental operations, the erase/write initiate sequence must be written to the NVMKEY register to allow any erase or program operation to proceed. The three instructions following the start of the programming sequence should be NOPs. To start an erase or write sequence, the following steps must be completed:

- 1. Write 55h to the NVMKEY register.
- 2. Write AAh to the NVMKEY register.
- 3. Set the WR bit in the NVMCON register.
- 4. Execute three NOP instructions.

All erase and write cycles are self-timed. The WR bit can be polled to supply enough PGECx clock pulses for the operation and determine if the erase or write cycle has been completed.

3.5 Erasing Program Memory

The general procedure for erasing user memory is shown in Figure 3-5. The process for Chip Erase, Page Erase and Inactive Partition Erase are all similar, and are described in Table 3-4 through Table 3-6.

The last row of the last page of program memory contains the Configuration Words. Before programming these Words, they must be erased. If they are erased with a Page Erase operation, all other rows in the page will also be erased. Users may want to either avoid using the rest of this page for application code or ensure that the non-configuration data in the CS page is copied before the erase and reprogrammed afterwards.

- **Note 1:** Program memory must be erased before writing any data to program memory.
 - 2: For Page Erase operations, the NVMADR/NVMADRU registers must also be loaded with the address of the page to be erased.

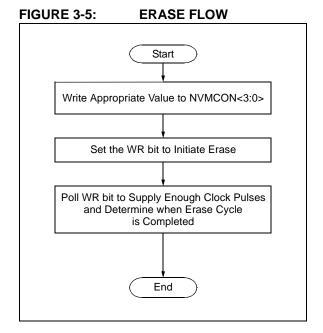


TABLE 3-4: SERIAL EXECUTION FOR CHIP ERASE

Command (Binary)	Data (Hex)		Description		
Step 1: Exit the R	Step 1: Exit the Reset vector.				
0000	000000	NOP			
0000	040200	GOTO	0x200		
0000	000000	NOP			
Step 2: Configure	the NVMCON reg	ister to p	perform a Chip Erase.		
0000	2400E0	MOV	#0x400E, W0		
0000	883B00	MOV	WO, NVMCON		
Step 3: Set the W	R bit.	•			
0000	200550	MOV	#0x55, W0		
0000	883B30	MOV	WO, NVMKEY		
0000	200AA0	MOV	#0xAA, W0		
0000	883B30	MOV	WO, NVMKEY		
0000	A8E761	BSET	NVMCON, #WR		
0000	000000	NOP			
0000	000000	NOP			
0000	000000	NOP			
Step 4: Repeat this step to poll the WR bit until it is cleared by hardware.					
0000	040200	GOTO	0x200		
0000	000000	NOP			
0000	803B02	MOV	NVMCON, W2		
0000	000000	NOP			
0000	883C22	MOV	W2, VISI		
0000	000000	NOP			
0001	<visi></visi>	; Cloc	ck out the contents of the VISI register.		
0000	000000	NOP			
Step 5: Clear the	WREN bit.				
0000	200000	MOV	#0000, W0		
0000	883B00	MOV	WO, NVMCON		

TABLE 3-5: SERIAL EXECUTION FOR PAGE ERASE

IADLL 3-3.	SENIAL EXECU		SK I AGE EKAGE
Command (Binary)	Data (Hex)		Description
Step 1: Exit the F	Reset vector.		
0000	000000	NOP	
0000	040200	GOTO	0x200
0000	000000	NOP	
Step 2: Set the N	VMCON register to	erase a	page.
0000	240030	MOV	#0x4003, W0
0000	883B00	MOV	WO, NVMCON
Step 3: Load the	address of the pag	e to be e	rased into the NVMADR register pair.
0000	200000	MOV	#PAGE_ADDR_LO, WO
0000	883B10	MOV	WO, NVMADR
0000	200000	MOV	<pre>#PAGE_ADDR_HI, W0</pre>
0000	883B20	MOV	WO, NVMADRU
Step 4: Set the W	/R bit.		
0000	200550	MOV	#0x55, W0
0000	883B30	MOV	WO, NVMKEY
0000	200AA0	MOV	#0xAA, W0
0000	883B30	MOV	WO, NVMKEY
0000	A8E761	BSET	NVMCON, #WR
0000	000000	NOP	
0000	000000	NOP	
0000	000000	NOP	
Step 5: Repeat th	nis step to poll the V	VR bit ur	ntil it is cleared by hardware.
0000	040200	GOTO	0x200
0000	000000	NOP	
0000	803B02	MOV	NVMCON, W2
0000	000000	NOP	
0000	883C22	MOV	W2, VISI
0000	000000	NOP	
0001	<visi></visi>		ck out the contents of the VISI register.
0000	000000	NOP	
Step 6: Clear the	WREN bit.		
0000	200000	MOV	#0000, W0
0000	883B00	MOV	WO, NVMCON

TABLE 3-6: SERIAL EXECUTION FOR INACTIVE PARTITION ERASE

Command (Binary)	Data (Hex)		Description
Step 1: Exit the R	eset vector.		
0000	000000	NOP	
0000	040200	GOTO	0x200
0000	000000	NOP	
Step 2: Set the N	VMCON register to	erase a	page.
0000	240040	MOV	#0x4004, W0
0000	883B00	MOV	WO, NVMCON
Step 3: Set the W	R bit.		
0000	200550	MOV	#0x55, W0
0000	883B30	MOV	WO, NVMKEY
0000	200AA0	MOV	#0xAA, W0
0000	883B30	MOV	WO, NVMKEY
0000	A8E761	BSET	NVMCON, #WR
0000	000000	NOP	
0000	000000	NOP	
0000	000000	NOP	
Step 4: Repeat th	is step to poll the V	VR bit ur	ntil it is cleared by hardware.
0000	040200	GOTO	0x200
0000	000000	NOP	
0000	803B02	MOV	NVMCON, W2
0000	883C22	MOV	W2, VISI
0000	000000	NOP	
0001	<visi></visi>	; Clo	ck out the contents of the VISI register.
0000	000000	NOP	
Step 5: Clear the	WREN bit.	•	
0000	200000	MOV	#0000, W0
0000	883B00	MOV	WO, NVMCON
	l	L	

3.6 Writing Code Memory

For PIC24FJ1024GA610/GB610 devices, code memory is written in three steps:

- Writing the data to memory-mapped write latches (located in the configuration memory space at addresses, FA0000h through FA00FEh);
- · Setting a destination address
- Initiating the memory write sequence

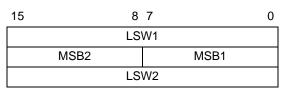
There are two methods available for writing to code memory: double-word writes using the write latches or 128-word row writes. Figure 3-7 provides a high-level description of the two methods.

Double-word writes program code memory with two instruction words at a time. Two words are loaded into the write latches. Next, the write sequence is initiated, and finally, the WR bit is checked for the sequence to be complete. This process continues for all the data to be programmed.

Table 3-7 provides an example of ICSP programming for a double-word write operation.

The data loaded into the programming latches must be in the packed format, as shown in Figure 3-6.

FIGURE 3-6: PACKED INSTRUCTION WORD FORMAT



LSWx: Least Significant 16 bits of instruction word MSBx: Most Significant Byte of instruction word

Note: When the number of instruction words transferred is odd, MSB2 is zero and LSW2 cannot be transmitted.

FIGURE 3-7: PROGRAM CODE MEMORY FLOW

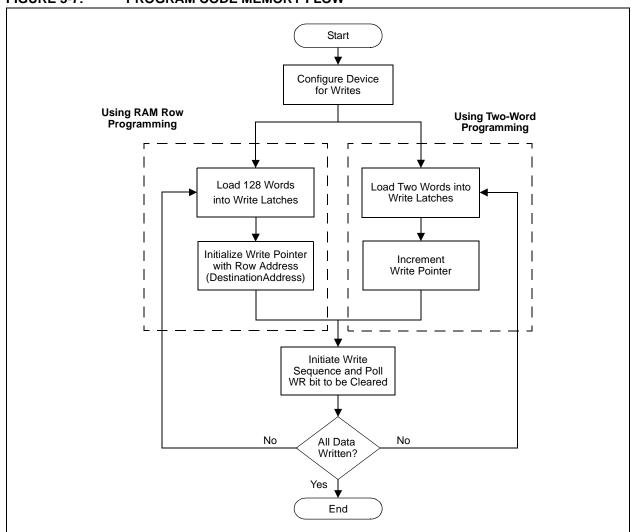


TABLE 3-7: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CODE MEMORY: TWO-WORD LATCH WRITES

Command	Data		Description	
(Binary)	(Hex)		Description	
Step 1: Exit the	e Reset vector.			
0000	000000	NOP		
0000	040200	GOTO	0x200	
0000	000000	NOP		
Step 2: Initializ	e the TBLPAG r	egister for w	rriting to the latches.	
0000	200FAC	MOV	#0xFA, W12	
0000	8802AC	MOV	W12, TBLPAG	
Step 3: Load V	V0:W2 with the r	next two pac	ked instruction words to program.	
0000	2xxxx0	MOV	# <lsw0>, W0</lsw0>	
0000	2xxxx1	MOV	# <msb1:msb0>, W1</msb1:msb0>	
0000	2xxxx2	MOV	# <lsw1>, W2</lsw1>	
Step 4: Set the	Read Pointer (W6) and Wri	ite Pointer (W7), and load the (next set of) write latches.	
0000	EB0300	CLR	W6	
0000	000000	NOP		
0000	EB0380	CLR	W7	
0000	000000	NOP		
0000	BB0BB6	TBLWTL	[W6++], [W7]	
0000	000000	NOP		
0000	000000	NOP		
0000	BBDBB6	TBLWTH.B	[W6++], [W7++]	
0000	000000	NOP		
0000	000000	NOP		
0000	BBEBB6		[W6++], $[++W7]$	
0000	000000	NOP		
0000	000000	NOP		
0000	BB1BB6		[W6++], [W7++]	
0000	000000	NOP		
0000	000000	NOP		
Step 5: Set the NVMADRU/NVMADR register pair to point to the correct address.				
0000	2xxxx3	MOV	<pre>#DestinationAddress<15:0>, W3</pre>	
0000	2xxxx4	MOV	<pre>#DestinationAddress<23:16>, W4</pre>	
0000	883B13	MOV	W3, NVMADR	
0000	883B24	MOV	W4, NVMADRU	
			am two instruction words.	
0000	24001A		#0x4001, W10	
0000	883B0A	MOV	W10, NVMCON	
0000	000000	NOP		
Step 7: Initiate	the write cycle.	1		
0000	200551	MOV	#0x55, W1	
0000	883B31	MOV	W1, NVMKEY	
0000	200AA1	MOV	#0xAA, W1	
0000	883B31	MOV	W1, NVMKEY	
0000	A8E761	BSET	NVMCON, #WR	
0000	000000	NOP		
0000	000000	NOP		
0000	000000	NOP		

TABLE 3-7: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CODE MEMORY: TWO-WORD LATCH WRITES (CONTINUED)

Command (Binary)	Data (Hex)	Description				
Step 8: Wait fo	or the program o	peration to complete and make sure the WR bit is clear.				
0000	803B00	MOV NVMCON, WO				
0000	883C20	MOV W0, VISI				
0000	000000	NOP				
0001	<visi></visi>	; Clock out the contents of the VISI register.				
0000	000000	NOP				
0000	040200	GOTO 0x200				
0000	000000	NOP				
_	_	; Repeat until the WR bit is clear.				
Step 9: Repea	Step 9: Repeat Steps 3-8 until all code memory is programmed.					
Step 10: Clear the WREN bit.						
0000	200000	MOV #0000, WO				
0000	883B00	MOV W0, NVMCON				

Row writes program one row (128 instruction words) at a time. First, the Table Pointer is initialized to point to the program latches and data is written into them with Table Writes. Next, the Write Pointer is initialized (NVMADRU and NVMADR register pair) with the row address (DestinationAddress). Finally, the write sequence is initiated and the WR bit is checked for the row programming to be complete. This process is repeated for all data to be programmed. Table 3-8 shows the ICSP programming details for row writes.

To minimize programming time, the data to be programmed is stored in the W0:W5 registers in a packed data format (Figure 3-8). This is the same packed format used by the PE. See Section 6.2.2 "Packed Data Format" for additional information.

FIGURE 3-8: PACKED INSTRUCTION WORD STORAGE IN W0:W5

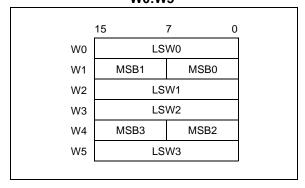


TABLE 3-8: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CODE MEMORY: ROW WRITES

Command (Binary)	Data (Hex)		Description
Step 1: Exit the F	Reset vector.		
0000	000000	NOP	
0000	040200	GOTO	0x200
0000	000000	NOP	
Step 2: Set the N	IVMCON register t	o program	128 instruction words.
0000	240020	MOV	#0X4002, W0
0000	883B00	MOV	WO, NVMCON
Step 3: Initialize	the TBLPAG regist	er for writin	ng to the latches.
0000	200FAC	MOV	#0xFA, W12
0000	8802AC	MOV	W12, TBLPAG
Step 4: Load W0	:W5 with the next	4 instructio	n words to program.
0000	2xxxx0	MOV	# <lsw0>, W0</lsw0>
0000	2xxxx1	MOV	# <msb1:msb0>, W1</msb1:msb0>
0000	2xxxx2	MOV	# <lsw1>, W2</lsw1>
0000	2xxxx3	MOV	# <lsw2>, W3</lsw2>
0000	2xxxx4	MOV	# <msb3:msb2>, W4</msb3:msb2>
0000	2xxxx5	MOV	# <lsw3>, W5</lsw3>
Step 5: Set the R	Read Pointer (W6)	and load th	e (next set of) write latches.
0000	EB0300	CLR	W6
0000	000000	NOP	
0000	EB0380	CLR	W7
0000	000000	NOP	
0000	BB0BB6	TBLWTL	[W6++], [W7]
0000	000000	NOP	
0000	000000	NOP	

TABLE 3-8: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CODE MEMORY: ROW WRITES (CONTINUED)

	KOW WKIIES	(00.11.11.02.2)			
Command (Binary)	Data (Hex)	Description			
Step 5: Set the R	ead Pointer (W6)	and load the (next set of) write latches. (Continued)			
0000	BBDBB6	TBLWTH.B [W6++], [W7++]			
0000	000000	NOP			
0000	000000	NOP			
0000	BBEBB6	TBLWTH.B [W6++], [++W7]			
0000	000000	NOP			
0000	000000	NOP			
0000	BB1BB6	TBLWTL [W6++], [W7++]			
0000	000000	NOP			
0000	000000	NOP			
0000	BB0BB6	TBLWTL [W6++], [W7]			
0000	000000	NOP			
0000	000000	NOP			
0000	BBDBB6	TBLWTH.B [W6++], [W7++]			
0000	000000	NOP			
0000	000000	NOP			
0000	BBEBB6	TBLWTH.B [W6++], [++W7]			
0000	000000	NOP			
0000	000000	NOP			
0000	BB1BB6	TBLWTL [W6++], [W7++]			
0000	000000	NOP			
0000	000000	NOP			
Step 6: Repeat S	teps 4 and 5, for a	a total of 32 times, to load the write latches with 128 instructions.			
Step 7: Set the N	VMADRU/NVMAI	DR register pair to point to the correct address.			
0000	2xxxx3	MOV #DestinationAddress<15:0>, W3			
0000	2xxxx4	MOV #DestinationAddress<23:16>, W4			
0000	883B13	MOV W3, NVMADR			
0000	883B24	MOV W4, NVMADRU			
Step 8: Execute the WR bit unlock sequence and initiate the write cycle.					
0000	200550	MOV #0x55, W0			
0000	883B30	MOV WO, NVMKEY			
0000	200AA0	MOV #0xAA, W0			
0000	883B30	MOV WO, NVMKEY			
0000	A8E761	BSET NVMCON, #WR			
0000	000000	NOP			
0000	000000	NOP			
0000	000000	NOP			
Step 9: Repeat th	is step to poll the	WR bit until it is cleared by hardware.			
0000	040200	GOTO 0x200			
0000	000000	NOP			
0000	803B02	MOV NVMCON, W2			
0000	883C22	MOV W2, VISI			
0000	000000	NOP			
0001	<visi></visi>	; Clock out the contents of the VISI register.			
0000	000000	NOP			
Step 10: Reset th	e device's interna	al Program Counter (PC).			
0000	040200	GOTO 0x200			
0000	000000	NOP			
Step 11: Repeat	Steps 3 through 9	until all code memory is programmed.			
Step 12: Clear the	e WREN bit.				
0000	200000	MOV #0000, W0			
0000	883B00	MOV W0, NVMCON			

3.7 Writing Configuration Bits

The procedure for writing Configuration bits is similar to the procedure for writing code memory. Table 3-9 shows the ICSP programming details for writing the Configuration bits.

For all Configuration Words, except FBTSEQ, only the lower two bytes of the word contain configuration data; the upper byte is unused. It is recommended that the upper byte be programmed with FFFFh, as indicated in Step 3. FBTSEQ (implemented only in Dual Partition Flash modes) uses all 24 bits of the program memory's width to store the Boot Sequence Number bits, BSEQ<11:0> (FBTSEQ<11:0>), and its one's complement, IBSEQ<11:0> (FBTSEQ<23:12>).

To change the values of the Configuration bits once they have been programmed, the device must be erased, as described in **Section 3.5 "Erasing Program Memory"**, and reprogrammed to the desired value. Note that it is only possible to program a Configuration bit from '1' to '0' to enable code protection; it is not possible to program it from '0' to '1'.

TABLE 3-9: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CONFIGURATION WORDS: TWO-WORD LATCH WRITES

	1110-1101	D LATCH WRITES
Command (Binary)	Data (Hex)	Description
Step 1: Exit th	e Reset vector.	
0000	000000	NOP
0000	040200	GOTO 0x200
0000	000000	NOP
Step 2: Initializ	ze the TBLPAG	register for writing to the latches.
0000	200FAC	MOV #0xFA, W12
0000	8802AC	MOV W12, TBLPAG
Step 3: Load	W0:W1 with the	next two Configuration Words to program.
0000	2xxxx0	MOV # <config data="" lower="" word="">, W0</config>
0000	2FFxx1	MOV # <config data="" upper="" word="">, W1</config>
_	_	; Upper word is OxFFFF for all Configuration Words except FBTSEQ
0000	2FFFF2	MOV #0xffff, W2
Step 4: Set the	e Read Pointer	(W6) and Write Pointer (W7), and load the (next set of) write latches.
0000	EB0300	CLR W6
0000	000000	NOP
0000	EB0380	CLR W7
0000	000000	NOP
0000	BB0BB6	TBLWTL [W6++], [W7]
0000	000000	NOP
0000	000000	NOP
0000	BBDBB6	TBLWTH.B [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BBEBB6	TBLWTH.B [W6++], [++W7]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL.W [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
Step 5: Set the	e NVMADRU/N	VMADR register pair to point to the correct address.
0000	2xxxx3	MOV #DestinationAddress<15:0>, W3
0000	2xxxx4	MOV #DestinationAddress<23:16>, W4
0000	883B13	MOV W3, NVMADR
0000	883B24	MOV W4, NVMADRU

TABLE 3-9: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CONFIGURATION WORDS: TWO-WORD LATCH WRITES (CONTINUED)

	TWO-WORD EATON WRITES (GONTHROLD)				
Command (Binary)	Data (Hex)		Description		
Step 6: Set the	e NVMCON reg	jister to pr	ogram two instruction words.		
0000	24001A	MOV	#0x4001, W10		
0000	883B0A	MOV	W10, NVMCON		
0000	000000	NOP			
Step 7: Initiate	e the write cycle).			
0000	200551	MOV	#0x55, W1		
0000	883B31	MOV	W1, NVMKEY		
0000	200AA1	MOV	#0xAA, W1		
0000	883B31	MOV	W1, NVMKEY		
0000	A8E761	BSET	NVMCON, #WR		
0000	000000	NOP			
0000	000000	NOP			
0000	000000	NOP			
Step 8: Wait for program operation to complete and make sure the WR bit is clear.					
0000	803B00	MOV	NVMCON, WO		
0000	883C20	MOV	WO, VISI		
0000	000000	NOP			
0001	<visi></visi>	; Clock	out the contents of the VISI register.		
0000	000000	NOP			
0000	040200	GOTO	0x200		
0000	000000	NOP			
_	_	; Repea	at until the WR bit is clear.		
Step 9: Repea	at Steps 3-8 unt	il all code	memory is programmed.		
Step 10: Clea	r the WREN bit.				
0000	200000	MOV	#0000, W0		
0000	883B00	MOV	WO, NVMCON		

3.8 Reading Code Memory

Reading from code memory is performed by executing a series of TBLRD instructions and clocking out the data using the REGOUT command.

Table 3-10 shows the ICSP programming details for reading code memory.

To minimize reading time, the same packed data format that the PE uses is utilized. See **Section 6.2 "Programming Executive Commands"** for more details on the packed data format.

3.9 Reading Configuration Words

The procedure for reading Configuration Words is identical to the procedure for reading code memory, shown in Table 3-10. Since there are multiple Configuration Words, they are read one at a time.

TABLE 3-10: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY

Command (Binary)	Data (Hex)	Description
Step 1: Exit the Re	eset vector.	
0000	000000	NOP
0000	040200	GOTO 0x200
0000	000000	NOP
Step 2: Initialize th	e Write Pointer (W	7) to point to the VISI register.
0000	207847	MOV # VISI, W7
0000	000000	NOP
Step 3: Initialize th	e TBLPAG register	and the Read Pointer (W6) for a TBLRD instruction.
0000	200xx0	MOV # <sourceaddress23:16>, W0</sourceaddress23:16>
0000	8802A0	MOV W0, TBLPAG
0000	2xxxx6	MOV # <sourceaddress15:0>, W6</sourceaddress15:0>
-		nts of the next two locations of code memory, through the VISI register, using
the REGOUT comm	and.	
0000	BA0B96	TBLRDL [W6], [W7]
0000	000000	NOP
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	BADBB6	TBLRDH.B [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BAD3D6	TBLRDH.B [++W6], [W7]
0000	000000	NOP
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	BA0BB6	TBLRDL [W6++], [W7]
0000	000000	NOP
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
Step 5: Reset the	device's internal Pr	ogram Counter.
0000	040200	GOTO 0x200
0000	000000	NOP
Sten 6: Repeat St	ens 3 through 5 unt	il all desired code memory is read (note that "Reset the device's internal

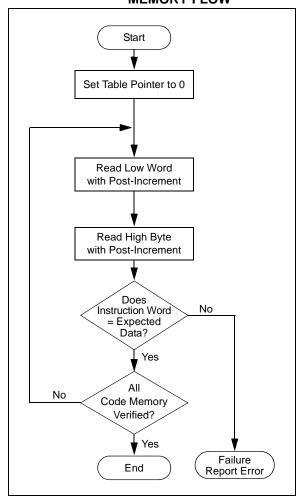
3.10 Verify Code Memory and Configuration Bits

The verify step involves reading back the code memory space and comparing it against the copy held in the programmer's buffer. The Configuration Words are verified with the rest of the code.

The verify process is illustrated in Figure 3-9. The lower word of the instruction is read, and then the lower byte of the upper word is read and compared against the instruction stored in the programmer's buffer. Refer to **Section 3.8** "Reading Code Memory" for implementation details of reading code memory.

Note: Because the Configuration bytes include the device code protection bit, code memory should be verified immediately after writing if the code protection is to be enabled. This is because the device will not be readable or verifiable if a device Reset occurs after the code-protect bit has been cleared.

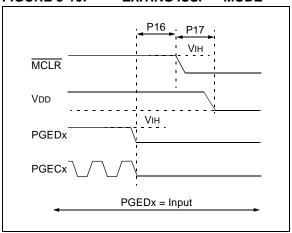
FIGURE 3-9: VERIFY CODE MEMORY FLOW



3.11 Exiting ICSP Mode

Exiting Program/Verify mode is done by removing VIH from MCLR, as illustrated in Figure 3-10. The only requirement for exit is that an interval, P16, should elapse between the last clock, and program signals on PGECx and PGEDx, before removing VIH.

FIGURE 3-10: EXITING ICSP™ MODE



4.0 DEVICE PROGRAMMING – ENHANCED ICSP

This section discusses programming the device through Enhanced ICSP and the Programming Executive (PE). The PE resides in executive memory (separate from code memory) and is executed when Enhanced ICSP Programming mode is entered. The PE provides the mechanism for the programmer (host device) to program and verify the PIC24FJ1024GA610/GB610 family devices, using a simple command set and communication protocol. There are several basic functions provided by the PE:

- · Read Memory
- · Erase Memory
- · Program Memory
- · Blank Check

The PE performs the low-level tasks required for erasing, programming and verifying a device. This allows the programmer to program the device by issuing the appropriate commands and data. A detailed description for each command is provided in Section 6.2 "Programming Executive Commands".

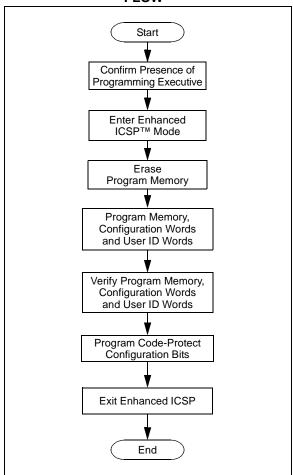
Note:

The PE uses the device's data RAM for variable storage and program execution. After running the PE, no assumptions should be made about the contents of data RAM.

4.1 Overview of the Programming Process

Figure 4-1 shows the high-level overview of the programming process. First, it must be determined if the PE is present in executive memory, and then, Enhanced ICSP mode is entered. The program memory is then erased, and the program memory and Configuration Words are programmed and verified. Last, the code-protect Configuration bits are programmed (if required) and Enhanced ICSP mode is exited.

FIGURE 4-1: HIGH-LEVEL ENHANCED ICSP™ PROGRAMMING FLOW



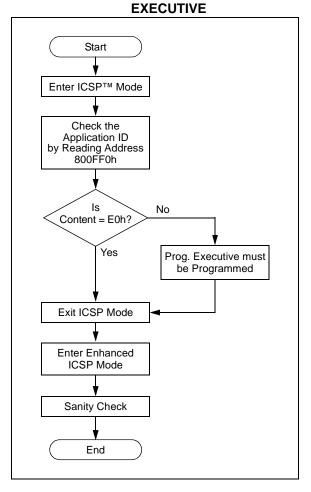
4.2 Confirming the Presence of the Programming Executive

Before programming, the programmer must confirm that the PE is stored in executive memory. The procedure for this task is illustrated in Figure 4-2.

First, ICSP mode is entered. Then, the unique Application ID Word, stored in executive memory, is read. If the Application ID has the value, E0h, the Programming Executive is resident in memory and the device can be programmed. However, if the Application ID Word is not present, the PE must be programmed to executive code memory using the method described in Section 5.0 "Programming the Programming Executive to Memory".

Section 3.0 "Device Programming – ICSP" describes the ICSP programming method. Section 4.3 "Reading the Application ID Word" describes the procedure for reading the Application ID Word in ICSP mode.

FIGURE 4-2: CONFIRMING PRESENCE OF PROGRAMMING



4.3 Reading the Application ID Word

The Application ID Word is stored at address, 800FF0h, in executive code memory. To read this memory location, you must use the SIX control code to move this program memory location to the VISI register. Then, the REGOUT control code must be used to clock the contents of the VISI register out of the device. The corresponding control and instruction codes that must be serially transmitted to the device to perform this operation are shown in Table 4-1.

If the Application ID has the value, E0h, the Programming Executive is resident in memory and the device can be programmed using the mechanism described in this section. However, if the Application ID has any other value, the PE is not resident in memory; it must be loaded to memory before the device can be programmed. The procedure for loading the PE to memory is described in **Section 5.0 "Programming the Programming Executive to Memory"**.

TABLE 4-1: SERIAL INSTRUCTION EXECUTION FOR READING THE APPLICATION ID WORD

===	0 =: \; \; .= :: \ \ .	
Command (Binary)	Data (Hex)	Description
Step 1: Exit the Reset vector.		
0000	000000	NOP
0000	040200	GOTO 0x200
0000	000000	NOP
Step 2: Initialize TBLPAG and the Read Pointer (W0) for the TBLRD instruction.		
0000	200800	MOV #0x80, W0
0000	8802A0	MOV W0, TBLPAG
0000	20FF00	MOV #0xFF0, W0
0000	207841	MOV #VISI, W1
0000	000000	NOP
0000	BA0890	TBLRDL [W0], [W1]
0000	000000	NOP
0000	000000	NOP
0000	000000	NOP
Step 3: Output the VISI register using the REGOUT command.		
0001	<visi></visi>	; Clock out the contents of the VISI register.

4.4 Entering Enhanced ICSP Mode

As illustrated in Figure 4-3, entering Enhanced ICSP Program/Verify mode requires three steps:

- The MCLR pin is briefly driven high and then low.
- 2. A 32-bit key sequence is clocked into PGEDx.
- 3. MCLR is then driven high within a specified period of time and held.

The programming voltage applied to \overline{MCLR} is VIH, which is essentially VDD in PIC24FJ1024GA610/GB610 family devices. There is no minimum time requirement for holding at VIH. After VIH is removed, an interval of at least P18 must elapse before presenting the key sequence on PGEDx.

The key sequence is a specific 32-bit pattern, '0100 1101 0100 0011 0100 1000 0101 0000' (more easily remembered as 4D434850h in hexadecimal format). The device will enter Program/Verify mode only if the key sequence is valid. The Most Significant bit (MSb) of the most significant nibble must be shifted in first.

Once the key sequence is complete, VIH must be applied to MCLR and held at that level for as long as Program/Verify mode is to be maintained. An interval time of at least P19, P7 and P1 * 5 must elapse before presenting data on PGEDx. During the P7 interval, the programmer's PGEDx and PGECx lines must be tri-stated.

On successful entry, the program memory can be accessed and programmed in serial fashion. While in the Program/Verify mode, all unused I/Os are placed in the high-impedance state.

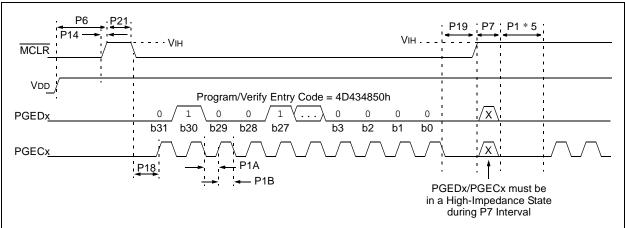
4.5 Blank Check

The term, "Blank Check", implies verifying that the device has been successfully erased and has no programmed memory locations. A blank or erased memory location is always read as '1'.

The Device ID registers (FF0000h:FF0002h) can be ignored by the Blank Check, since this region stores device information that cannot be erased. Additionally, all unimplemented memory space should be ignored by the Blank Check.

The QBLANK command is used for the Blank Check. It determines if the code memory is erased by testing these memory regions. A 'BLANK' or 'NOT BLANK' response is returned. If it is determined that the device is not blank, it must be erased before attempting to program the chip.





4.6 Code Memory Programming

4.6.1 PROGRAMMING METHODOLOGY

There are two commands that can be used for programming code memory when utilizing the PE. The PROG2W command programs and verifies two 24-bit instruction words into the program memory, starting at the address specified. The second and faster command, PROGP, allows up to 128 instruction words (each 24 bits) to be programmed and verified into program memory, starting at the address specified. See Section 6.0 "The Programming Executive" for a full description for each of these commands.

Figure 4-4 and Figure 4-5 show the programming methodology for the PROG2W and PROGP commands. In both instances, 87552 instruction words of the device are programmed.

Note: If a bootloader needs to be programmed, its code must not be programmed into the first page of code memory. For example, if a bootloader, located at address, 200h, attempts to erase the first page, it would inadvertently erase itself. Instead, program the bootloader into the second page (e.g., 400h).

FIGURE 4-4: FLOWCHART FOR DOUBLE-WORD PROGRAMMING

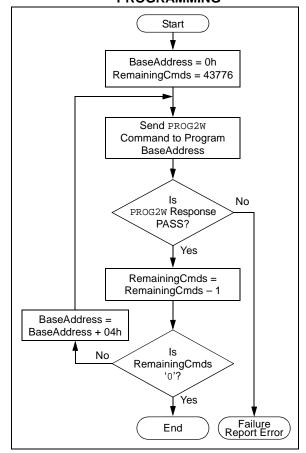
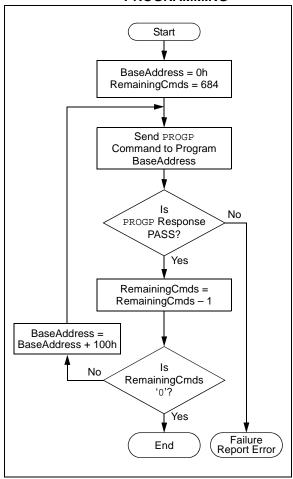


FIGURE 4-5: FLOWCHART FOR MULTIPLE WORD PROGRAMMING

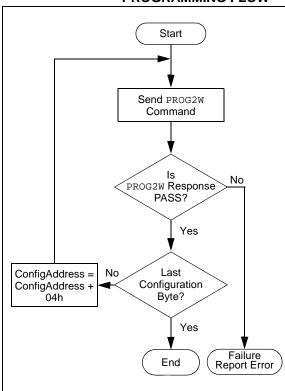


4.7 Configuration Bit Programming

Configuration bits are programmed one at a time using the PROG2W command. This command specifies the configuration data and address. When Configuration bits are programmed, any unimplemented bits must be programmed with a '1'.

Multiple PROG2W commands are required to program all Configuration bits. A flowchart for Configuration bit programming is shown in Figure 4-6.

FIGURE 4-6: CONFIGURATION BIT PROGRAMMING FLOW



4.8 Programming Verification

After code memory is programmed, the contents of memory can be verified to ensure that programming was successful. Verification requires code memory to be read back and compared against the copy held in the programmer's buffer.

The READP command can be used to read back all the programmed code memory and Configuration Words.

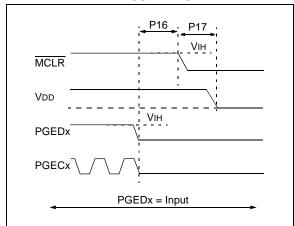
Alternatively, you can have the programmer perform the verification after the entire device is programmed using a checksum computation.

See **Section 9.0 "Checksum Computation"** for more information on calculating the checksum.

4.9 Exiting Enhanced ICSP Mode

Exiting Program/Verify mode is done by removing VIH from MCLR, as illustrated in Figure 4-7. The only requirement for exit is that an interval, P16, should elapse between the last clock, and program signals on PGECx and PGEDx, before removing VIH.

FIGURE 4-7: EXITING ENHANCED ICSP™ MODE



5.0 PROGRAMMING THE PROGRAMMING EXECUTIVE TO MEMORY

Note:

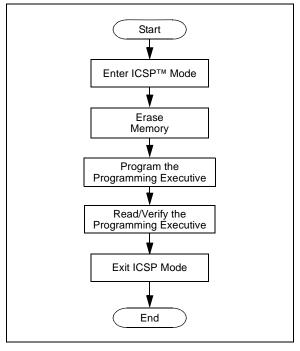
The Programming Executive (PE) can be obtained from each device page on the Microchip web site: www.microchip.com.

5.1 Overview

If it is determined that the PE is not present in executive memory (as described in **Section 4.2 "Confirming the Presence of the Programming Executive"**), the PE must be programmed to executive memory.

Figure 5-1 shows the high-level process of programming the PE into executive memory. First, ICSP mode must be entered, and executive memory and user memory are erased; then, the PE is programmed and verified. Finally, ICSP mode is exited.

FIGURE 5-1: HIGH-LEVEL PROGRAMMING EXECUTIVE PROGRAM FLOW



5.2 Erasing Executive Memory

Executive memory can be erased through a series of Page Erase operations, as shown in Figure 5-2. This consists of setting NVMCON to 4003h, executing the programming cycle and repeating for the rest of the pages of executive memory.

Table 5-1 illustrates the ICSP programming process for Bulk Erasing memory.

Note:

The PE must always be erased before it is programmed, as described in Figure 5-1.

FIGURE 5-2: BULK ERASE FLOW

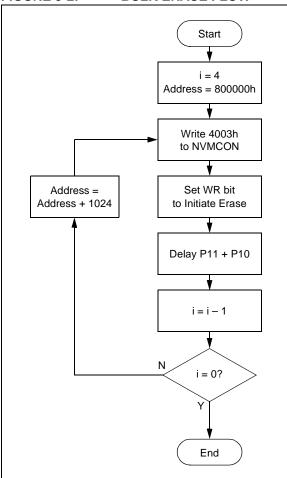


TABLE 5-1: SERIAL INSTRUCTION EXECUTION FOR ERASING EXECUTIVE MEMORY

TABLE 5-1: SERIAL INSTRUCTION EXECUTION FOR ERASING EXECUTIVE MEMORY					
Command (Binary)	Data (Hex)		Description		
Step 1: Exit the	e Reset vector.				
0000	000000	NOP			
0000	040200	GOTO	0x200		
0000	000000	NOP			
0000	000000	NOP			
Step 2: Set the	NVMCON registe	er to eras	se a page.		
0000	240030	MOV	#0x4003, W0		
0000	883B00	MOV	WO, NVMCON		
Step 3: Load to	he address of the	page to b	be erased into the NVMADR register pair.		
0000	200004	MOV	#0000, W4		
0000	883B14	MOV	W4, NVMADR		
0000	200800	MOV	#0080, W0		
0000	883B20	MOV	WO, NVMADRU		
Step 4: Set the	e WR bit.				
0000	200550	MOV	#0x55, W0		
0000	883B30	MOV	WO, NVMKEY		
0000	200AA0	MOV	#0xAA, W0		
0000	883B30	MOV	WO, NVMKEY		
0000	A8E761	BSET	NVMCON, #WR		
0000	000000	NOP			
0000	000000	NOP			
0000	000000	NOP			
Step 5: Repea	t this step to poll t	he WR b	it until it is cleared by hardware.		
0000	040200	GOTO	0x200		
0000	000000	NOP			
0000	803B02	MOV	NVMCON, W2		
0000	883C22	MOV	W2, VISI		
0000	000000	NOP			
0001	<visi></visi>		ck out the contents of the VISI register.		
0000	000000	NOP			
	nent W4 by 1024 (1			
0000	204003	MOV	#400, W3		
0000	418204	ADD	W3, W4, W4		
0000	883B14	MOV	W4, NVMADR		
Step 7: Repea	Step 7: Repeat Steps 4-6 until the entire test memory has been erased.				
Step 8: Clear t	he WREN bit.				
0000	200000	MOV#	0000, W0		
0000	883B00	MOV	WO, NVMCON		

5.3 Program the Programming Executive

Storing the PE to executive memory is similar to normal programming of code memory. The executive memory must first be erased and then programmed, using either two-word writes (two instruction words) or row writes (128 instruction words). The control flow for both methods is identical to that for programming code memory, as shown in Figure 3-7.

Table 5-2 and Table 5-3 illustrate the ICSP programming processes for PE memory. To minimize programming time, the same packed data format that the PE uses is utilized. See **Section 6.2 "Programming Executive Commands"** for more details on the packed data format.

TABLE 5-2: PROGRAMMING THE PROGRAMMING EXECUTIVE (TWO-WORD LATCH WRITES)

(IWO-WORD LATCH WRITES)				
Command (Binary)	Data (Hex)		Description	
Step 1: Exit th	e Reset vector.			
0000	000000	NOP		
0000	040200	GOTO	0x200	
0000	000000	NOP		
Step 2: Initializ	ze the TBLPAG regi	ster for writir	ng to the latches.	
0000	200FAC	MOV	#0xFA, W12	
0000	8802AC	MOV	W12, TBLPAG	
Step 3: Load V	W0:W2 with the nex	t two packed	I instruction words to program.	
0000	2xxxx0	MOV	# <lsw0>, W0</lsw0>	
0000	2xxxx1	MOV	# <msb1:msb0>, W1</msb1:msb0>	
0000	2xxxx2	MOV	# <lsw1>, W2</lsw1>	
Step 4: Set the	e Read Pointer (W6	and the Wr	ite Pointer (W7), and load the write latches.	
0000	EB0300	CLR	W6	
0000	000000	NOP		
0000	EB0380	CLR	w7	
0000	000000	NOP		
0000	BB0BB6	TBLWTL	[W6++], [W7]	
0000	000000	NOP		
0000	000000	NOP		
0000	BBDBB6	TBLWTH.B	[W6++], [W7++]	
0000	000000	NOP		
0000	000000	NOP		
0000	BBEBB6	TBLWTH.B	[W6++], [++W7]	
0000	000000	NOP		
0000	000000	NOP		
0000	BB1BB6	TBLWTL.W	[W6++], [W7++]	
0000	000000	NOP		
0000	000000	NOP		
Step 5: Set the	e NVMADRU/NVMA	DR register	pair to point to the correct row.	
0000	2xxxx3	MOV	<pre>#DestinationAddress<15:0>, W3</pre>	
0000	2xxxx4	MOV	<pre>#DestinationAddress<23:16>, W4</pre>	
0000	883B13	MOV	W3, NVMADR	
0000	883B24	MOV	W4, NVMADRU	
Step 6: Set the	e NVMCON register	to program	two instruction words.	
0000	24001A	MOV	#0x4001, W10	
0000	000000	NOP	•	
0000	883B0A	MOV	W10, NVMCON	
0000	000000	NOP		
0000	000000	NOP		
0000	000000	NOP		

TABLE 5-2: PROGRAMMING THE PROGRAMMING EXECUTIVE (TWO-WORD LATCH WRITES) (CONTINUED)

Command (Binary)	Data (Hex)		Description		
	. ,				
Step 7: Initiate	the write cycle.				
0000	200551	MOV	#0x55, W1		
0000	883B31	MOV	W1, NVMKEY		
0000	200AA1	MOV	#0xAA, W1		
0000	883B31	MOV	W1, NVMKEY		
0000	A8E761	BSET	NVMCON, #WR		
0000	000000	NOP			
0000	000000	NOP			
0000	000000	NOP			
Step 8: Wait fo	or program operation	on to comp	olete and make sure the WR bit is clear.		
0000	000000	NOP			
0000	803B00	MOV	NVMCON, W0		
0000	000000	NOP			
0000	883C20	MOV	WO, VISI		
0000	000000	NOP			
0001	<visi></visi>	; Cloc	k out the contents of the VISI register.		
0000	000000	NOP			
0000	040200	GOTO	0x200		
0000	000000	NOP			
_	_	; Repe	at until the WR bit is clear.		
Step 9: Repea	at Steps 3-8 until all	code mer	mory is programmed.		
Step 10: Clear	r the WREN bit.				
0000	200000	MOV	#0000, W0		
0000	883B00	MOV	WO, NVMCON		

TABLE 5-3: PROGRAMMING THE PROGRAMMING EXECUTIVE (ROW WRITES)

IABLE 5-3:		G THE PROGRAMMING EXECUTIVE (ROW WRITES)
Command (Binary)	Data (Hex)	Description
Step 1: Exit the	Reset vector.	
0000	000000	NOP
0000	040200	GOTO 0x200
0000	000000	NOP
Step 2: Set the N	NVMCON register t	o program 128 instruction words.
0000	240020	MOV #0X4002, W0
0000	883B00	MOV W0, NVMCON
Step 3: Initialize	the TBLPAG regist	er for writing to the latches.
0000	200FAC	MOV #0xFA, W12
0000	8802AC	MOV W12, TBLPAG
		4 instruction words to program.
0000	2xxxx0	MOV # <lsw0>, W0</lsw0>
0000	2xxxx1	MOV # <msb1:msb0>, W1</msb1:msb0>
0000	2xxxx2	MOV # <lsw1>, W2</lsw1>
0000	2xxxx3	MOV # <lsw2>, W3</lsw2>
0000	2xxxx4	MOV # <msb3:msb2>, W4</msb3:msb2>
0000	2xxxx5	MOV # <lsw3>, W5</lsw3>
Step 5: Set the F	Read Pointer (W6)	and load the (next set of) write latches.
0000	EB0300	CLR W6
0000	000000	NOP
0000	EB0380	CLR W7
0000	000000	NOP
0000	BB0BB6	TBLWTL [W6++], [W7]
0000	000000	NOP
0000	000000	NOP
0000	BBDBB6	TBLWTH.B [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BBEBB6	TBLWTH.B [W6++], [++W7]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BB0BB6	TBLWTL [W6++], [W7]
0000	000000	NOP
0000	000000	NOP
0000	BBDBB6	TBLWTH.B [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BBEBB6	TBLWTH.B [W6++], [++W7]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
	1	total of 32 times, to load the write latches with 128 instructions.

TABLE 5-3: PROGRAMMING THE PROGRAMMING EXECUTIVE (ROW WRITES) (CONTINUED)

Command (Binary)	Data (Hex)	Description
Step 7: Set the N	IVMADRU/NVMAD	DR register pair to point to the correct address.
0000	2xxxx3	MOV #DestinationAddress<15:0>, W3
0000	2xxxx4	MOV #DestinationAddress<23:16>, W4
0000	883B13	MOV W3, NVMADR
0000	883B24	MOV W4, NVMADRU
Step 8: Execute	the WR bit unlock	sequence and initiate the write cycle.
0000	200550	MOV #0x55, W0
0000	883B30	MOV W0, NVMKEY
0000	200AA0	MOV #0xAA, W0
0000	883B30	MOV WO, NVMKEY
0000	A8E761	BSET NVMCON, #WR
0000	000000	NOP
Step 9: Repeat th	his step to poll the	WR bit until it is cleared by hardware.
0000	040200	GOTO 0x200
0000	000000	NOP
0000	803B00	MOV NVMCON, W2
0000	A8E761	MOV W2, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
Step 10: Reset th	ne device's interna	Il Program Counter.
0000	040200	GOTO 0x200
0000	000000	NOP
Step 11: Repeat	Steps 3 through 9	until all code memory is programmed.
Step 12: Clear th	e WREN bit.	
0000	200000	MOV #0000, W0
0000	883B00	MOV W0, NVMCON
		·

5.4 Reading Executive Memory

Reading from executive memory is performed by executing a series of TBLRD instructions and clocking out the data using the REGOUT command.

Table 5-4 shows the ICSP programming details for reading executive memory.

To minimize reading time, the same packed data format that the PE uses is utilized. See **Section 6.2** "**Programming Executive Commands**" for more details on the packed data format.

TABLE 5-4: SERIAL EXECUTION FOR READING EXECUTIVE MEMORY

TABLE 3-4. 5	LINAL LALCOTIC	IN I OK KL	ADING EXECUTIVE MEMORY
Command (Binary)	Data (Hex)		Description
Step 1: Exit the Re	set vector.		
0000	000000	NOP	
0000	040200	GOTO	0x200
0000	000000	NOP	
Step 2: Initialize the	e TBLPAG register a	nd the Read	Pointer (W6) for the TBLRD instruction.
0000	200xx0	MOV	# <sourceaddress23:16>, W0</sourceaddress23:16>
0000	8802A0	MOV	WO, TBLPAG
0000	2xxxx6	MOV	<pre>#<sourceaddress15:0>, W6</sourceaddress15:0></pre>
Step 3: Initialize the	e Write Pointer (W7)	and store the	next four locations of code memory in W0:W5.
0000	EB0380	CLR	W7
0000	000000	NOP	
0000	BA1B96	TBLRDL	[W6], [W7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BADBB6	TBLRDH.B	[W6++], [W7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BADBD6	TBLRDH.B	[++\\6], [\\7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BA1BB6	TBLRDL	[W6++], [W7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BA1B96	TBLRDL	[W6], [W7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BADBB6	TBLRDH.B	[W6++], [W7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BADBD6	TBLRDH.B	[++W6], [W7++]
0000	000000	NOP	
0000	000000	NOP	
0000	BA0BB6	TBLRDL	[W6++], [W7]
0000	000000	NOP	
0000	000000	NOP	

TABLE 5-4: SERIAL EXECUTION FOR READING EXECUTIVE MEMORY (CONTINUED)

Command (Binary)	Data (Hex)	Description
Step 4: Output Wo	:W5 using the VISI r	egister and the REGOUT command.
0000	883C20	MOV W0, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	883C21	MOV W1, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	883C22	MOV W2, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	883C23	MOV W3, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	883C24	MOV W4, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
0000	883C25	MOV W5, VISI
0000	000000	NOP
0001	<visi></visi>	; Clock out the contents of the VISI register.
0000	000000	NOP
Step 5: Reset the o	levice's internal Pro	gram Counter.
0000	000000	NOP
0000	040200	GOTO 0x200
0000	000000	NOP

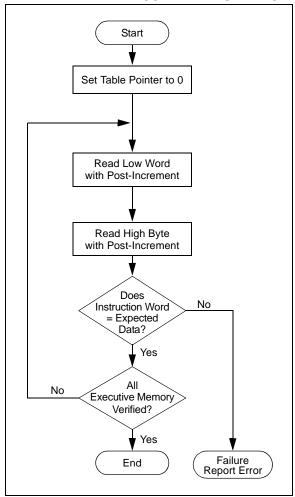
Step 6: Repeat Steps 3 through 5 until all desired code memory is read (note that "Reset the device's internal Program Counter" will be Step 5).

5.5 Verify Programming Executive

The verify step involves reading back the executive memory space and comparing it against the copy held in the programmer's buffer.

The verify process is illustrated in Figure 5-3. The lower word of the instruction is read, and then the lower byte of the upper word is read and compared against the instruction stored in the programmer's buffer. Refer to **Section 5.4** "**Reading Executive Memory**" for implementation details of reading executive memory.

FIGURE 5-3: VERIFY PROGRAMMING EXECUTIVE MEMORY FLOW



6.0 THE PROGRAMMING EXECUTIVE

6.1 Programming Executive Communication

The programmer and PE have a master-slave relationship, where the programmer is the master programming device and the PE is the slave.

All communication is initiated by the programmer in the form of a command. Only one command at a time can be sent to the PE. In turn, the PE only sends one response to the programmer after receiving and processing a command. The PE command set is described in **Section 6.2 "Programming Executive Commands"**. The response set is described in **Section 6.3 "Programming Executive Responses"**.

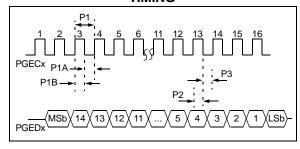
6.1.1 COMMUNICATION INTERFACE AND PROTOCOL

The ICSP/Enhanced ICSP interface is a 2-wire SPI, implemented using the PGECx and PGEDx pins. The PGECx pin is used as a clock input pin and the clock source must be provided by the programmer. The PGEDx pin is used for sending command data to, and receiving response data from, the PE.

Note: For Enhanced ICSP, all serial data is transmitted on the falling edge of PGECx and latched on the rising edge of PGECx. All data transmissions are sent to the MSb first, using 16-bit mode (see Figure 6-1).

Since a 2-wire SPI is used, and data transmissions are bidirectional, a simple protocol is used to control the direction of PGEDx. When the programmer completes a command transmission, it releases the PGEDx line and allows the PE to drive this line high. The PE keeps the PGEDx line high to indicate that it is processing the command.

FIGURE 6-1: PROGRAMMING EXECUTIVE SERIAL TIMING



After the PE has processed the command, it brings PGEDx low (P9b) to indicate to the programmer that the response is available to be clocked out. The programmer can begin to clock out the response after a maximum wait (P9b) and it must provide the necessary amount of clock pulses to receive the entire response from the PE.

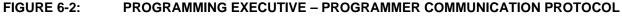
After the entire response is clocked out, the programmer should terminate the clock on PGECx until it is time to send another command to the PE. This protocol is illustrated in Figure 6-2.

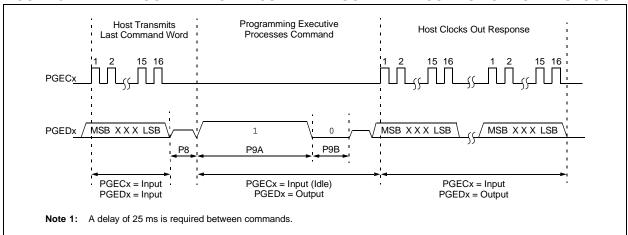
6.1.2 SPI RATE

In Enhanced ICSP mode, the PIC24FJ1024GA610/ GB610 family devices operate from the Fast Internal RC (FRC) Oscillator, which has a nominal frequency of 8 MHz. This oscillator frequency yields an effective system clock frequency of 4 MHz. To ensure that the programmer does not clock too fast, it is recommended that a 2 MHz clock be provided by the programmer.

6.1.3 SPI PORT SCANNING

In order to detect the proper SPI port used for the PE, a specific software routine must be used. This software is included in **Appendix A: "SPI Port Scanning Software Routine"** and must be used for all PIC24FJ1024GA610/GB610 family devices.





6.1.4 TIME-OUTS

The PE uses no Watchdog Timer or time-out for transmitting responses to the programmer. If the programmer does not follow the flow control mechanism using PGECx, as described in **Section 6.1.1** "Communication Interface and Protocol", it is possible that the PE will behave unexpectedly while trying to send a response to the programmer. Since the PE has no time-out, it is imperative that the programmer correctly follow the described communication protocol.

As a safety measure, the programmer should use the command time-outs identified in Table 6-1. If the command time-out expires, the programmer should reset the PE and start programming the device again.

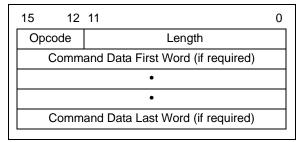
6.2 Programming Executive Commands

The PE command set is shown in Table 6-1. This table contains the opcode, mnemonic, length, time-out and description for each command. Functional details on each command are provided in the command descriptions (Section 6.2.4 "Command Descriptions").

6.2.1 COMMAND FORMAT

All PE commands have a general format, consisting of a 16-bit header and any required data for the command (see Figure 6-3). The 16-bit header consists of a 4-bit opcode field, which is used to identify the command, followed by a 12-bit command length field.

FIGURE 6-3: COMMAND FORMAT



The command opcode must match one of those in the command set. Any command that is received, which does not match the list in Table 6-1, will return a "NACK" response (see Section 6.3.1.1 "Opcode Field").

The command length is represented in 16-bit words since the SPI operates in 16-bit mode. The PE uses the command length field to determine the number of words to read from the SPI port. If the value of this field is incorrect, the command will not be properly received by the PE.

TABLE 6-1: PROGRAMMING EXECUTIVE COMMAND SET

Opcode	Mnemonic	Length (16-bit words)	Time-out	Description		
0x0	SCHECK	1	1 ms	Sanity check.		
0x1	READC	3	1 ms	Read an 8-bit word from the specified Configuration register or Device ID register.		
0x2	READP	4	1 ms/row	Read 'N' 24-bit instruction words of primary Flash memory, starting from the specified address.		
0x3	PROG2W	6	5 ms	Program a double instruction word of code memory at the specified address and verify.		
0x4	Reserved	N/A	N/A	This command is reserved; it will return a NACK.		
0x5	PROGP	99	5 ms	Program 128 words of program memory at the specified starting address, then verify.		
0x6	Reserved	N/A	N/A	This command is reserved; it will return a NACK.		
0x7	ERASEB	1	125 ms	Chip Erase the device.		
0x8	Reserved	N/A	N/A	This command is reserved; it will return a NACK.		
0x9	ERASEP	3	25 ms/page	Command to erase a page.		
0xA	Reserved	N/A	N/A	This command is reserved; it will return a NACK.		
0xB	QVER	1	1 ms	Query the PE software version.		
0xC	CRCP	5	1s	Performs a CRC-16 on the specified range of memory.		
0xD	Reserved	N/A	N/A	This command is reserved; it will return a NACK.		
0xE	QBLANK	5	700 ms	Query to check whether the code memory is blank.		

6.2.2 PACKED DATA FORMAT

When 24-bit instruction words are transferred across the 16-bit SPI interface, they are packed to conserve space using the format illustrated in Figure 6-4. This format minimizes traffic over the SPI and provides the PE with data that is properly aligned for performing Table Write operations.

FIGURE 6-4: PACKED INSTRUCTION WORD FORMAT

15		8	7	0	
	LSW1				
	MSB2 MSB1				
	LSW2				

LSWx: Least Significant 16 bits of instruction word MSBx: Most Significant Byte of instruction word

Note: When the number of instruction words transferred is odd, MSB2 is zero and LSW2 cannot be transmitted.

6.2.3 PROGRAMMING EXECUTIVE ERROR HANDLING

The PE will "NACK" all unsupported commands. Additionally, due to the memory constraints of the PE, no checking is performed on the data contained in the programmer command. It is the responsibility of the programmer to command the PE with valid command arguments or the programming operation may fail. Additional information on error handling is provided in Section 6.3.1.3 "QE_Code Field".

6.2.4 COMMAND DESCRIPTIONS

All commands supported by the PE are described in Section 6.2.4.1 "SCHECK Command" through Section 6.2.4.10 "QBLANK Command".

6.2.4.1 SCHECK Command

15	12	11 0
	Opcode	Length

Table 6-2 shows the description for the SCHECK command.

TABLE 6-2: COMMAND DESCRIPTION

Field	Description
Opcode	0x0
Length	0x1

The SCHECK command instructs the PE to do nothing but generate a response. This command is used as a "Sanity Check" to verify that the PE is operational.

Expected Response (2 words):

0x1000 0x0002

Note:	This i	nstructio	n is	not	required	for
	prograr	but	is	provided	for	
	develor	oment pu	ırpose	s onl	у.	

6.2.4.2 READC Command

15 12 11 8 7 0

Opcode Length

N Addr_MSB

Addr_LS

Table 6-3 shows the description for the READC command.

TABLE 6-3: COMMAND DESCRIPTION

Field	Description
Opcode	0x1
Length	0x3
N	Number of 8-bit Configuration registers or Device ID registers to read (maximum of 256)
Addr_MSB	MSB of 24-bit source address
Addr_LS	Least Significant 16 bits of 24-bit source address

The READC command instructs the PE to read N Configuration registers or Device ID registers, starting from the 24-bit address specified by Addr_MSB and Addr_LS. This command can only be used to read 8-bit or 16-bit data.

When this command is used to read Configuration registers, the upper byte in every data word returned by the PE is 0x00 and the lower byte contains the Configuration register value.

Expected Response ((2 + N) words):

0x1100

2 + N

Configuration Register or Device ID Register 1

...

Configuration Register or Device ID Register N

Note: Reading unimplemented memory will cause the PE to reset. To prevent this from occurring, ensure that only memory locations present on a particular device are accessed.

6.2.4.3 READP Command

 15
 12
 11
 8
 7
 0

 Opcode
 Length

 N

 Reserved
 Addr_MSB

 Addr_LS

Table 6-4 shows the description for the READP command.

TABLE 6-4: COMMAND DESCRIPTION

Field	Description
Opcode	0x2
Length	0x4
N	Number of 24-bit instructions to read (maximum of 32768)
Reserved	0x0
Addr_MSB	MSB of 24-bit source address
Addr_LS	Least Significant 16 bits of 24-bit source address

The READP command instructs the PE to read N 24-bit words of code memory, starting from the 24-bit address specified by Addr_MSB and Addr_LS. This command can only be used to read 24-bit data. All data returned in response to this command uses the packed data format described in Section 6.2.2 "Packed Data Format".

Expected Response (2 + 3 * N/2 words for N even):

0x1200

2 + 3 * N/2

Least Significant Program Memory Word 1

...

Least Significant Data Word N

Expected Response (4 + 3 * (N - 1)/2 words) for N odd):

0x1200

4 + 3 * (N - 1)/2

Least Significant Program Memory Word 1

...

MSB of Program Memory Word N (zero-padded)

Note: Reading unimplemented memory will cause the PE to reset. To prevent this from occurring, ensure that only memory locations present on a particular device are accessed.

6.2.4.4 PROG2W Command

15	12	11	8	7		0
Opc	ode			Ler	ngth	
	Reserved Addr_MSB					
	Addr_LS					
			DataL	_LS		
	DataH_MSB DataL_MSB					
	DataH_LS					

Table 6-5 shows the description for the PROG2W command.

TABLE 6-5: COMMAND DESCRIPTION

Field	Description
Opcode	0x3
Length	0x6
DataL_MSB	MSB of 24-bit data for low instruction word
DataH_MSB	MSB of 24-bit data for high instruction word
Addr_MSB	MSB of 24-bit destination address
Addr_LS	Least Significant 16 bits of 24-bit destination address
DataL_LS	Least Significant 16 bits of 24-bit data for low instruction word
DataH_LS	Least Significant 16 bits of 24-bit data for high instruction word

The PROG2W command instructs the PE to program two instruction words of code memory (6 bytes) to the specified memory address.

After the words have been programmed to code memory, the PE verifies the programmed data against the data in the command.

Expected Response (2 words):

0x1300

0x0002

6.2.4.5 PROGP Command

15 12 11 8 7 0

Opcode		Length	
Rese	rved	Addr_MSB	
Addr_LS			
D_1			
D_2			
	D_N		

Table 6-6 shows the description for the PROGP command.

TABLE 6-6: COMMAND DESCRIPTION

Field	Description
Opcode	0x5
Length	0xC3
Reserved	0x0
Addr_MSB	MSB of 24-bit destination address
Addr_LS	Least Significant 16 bits of 24-bit destination address
D_1	16-bit Data Word 1
D_2	16-bit Data Word 2
•••	16-bit Data Word 3 through 191
D_192	16-bit Data Word 192

The PROGP command instructs the PE to program one row of code memory (128 instruction words) to the specified memory address. Programming begins with the row address specified in the command. The destination address should be a multiple of 0x100.

The data to program the memory, located in command words, D_1 through D_192, must be arranged using the packed instruction word format illustrated in Figure 6-4.

After all data has been programmed to code memory, the PE verifies the programmed data against the data in the command.

Expected Response (2 words):

0x1500 0x0002

Note: Refer to Table 2-2 for code memory size information.

6.2.4.6 **ERASEB** Command

15	12 11	8 7	0
Opco	ode	Leng	th

Table 6-7 shows the description for the ERASEB command.

TABLE 6-7: COMMAND DESCRIPTION

Field	Description
Opcode	0x7
Length	0x1

The ERASEB command instructs the PE to perform a Chip Erase (i.e., erase all of the primary Flash memory and code-protect bits).

Expected Response (2 words):

0x1700 0x0002

6.2.4.7 **ERASEP** Command

0 15 12 11 8 7

Opcode		Length
NUM_PAGES		Addr_MSB
Addr_LS		S

Table 6-6 shows the description for the ERASEP command.

TABLE 6-8: COMMAND DESCRIPTION

Field	Description
Opcode	0x9
Length	0x3
NUM_PAGES	Up to 255
Addr_MSB	Most Significant Byte of the 24-bit address
Addr_LS	Least Significant 16 bits of the 24-bit address

The ERASEP command instructs the PE to Page Erase [NUM_PAGES] of code memory. The code memory must be erased at an "even" 1024 instruction words address boundary.

Expected Response (2 words):

0x1900 0x0002 6.2.4.8 **OVER Command**

15	12	11	0
Opcode	!	Length	

Table 6-9 shows the description for the QVER command.

TABLE 6-9: COMMAND DESCRIPTION

Field	Description
Opcode	0xB
Length	0x1

The QVER command queries the version of the PE software stored in test memory. The "version.revision" information is returned in the response's QE_Code, using a single byte with the following format: main version in upper nibble and revision in the lower nibble (i.e., 0x23 means Version 2.3 of PE software).

Expected Response (2 words):

0x1BMN (where "MN" stands for Version M.N) 0x0002

6.2.4.9 CRCP Command

 15
 12
 11
 8
 7
 0

 Opcode
 Length

 Reserved
 Addr_MSB

 Addr_LSW

 Reserved
 Size_MSB

 Size LSW

Table 6-11 shows the description for the CRCP command.

TABLE 6-10: COMMAND DESCRIPTION

Field	Description
Opcode	Ch
Length	5h
Addr_MSB	Most Significant Byte of 24-bit address
Addr_LSW	Least Significant 16 bits of 24-bit address
Size	Number of 24-bit locations (address range divided by 2)

The CRCP command performs a CRC-16 on the range of memory specified. This command can substitute for a full chip verify. Data is shifted in a packed method, as demonstrated in Figure 6-4, byte-wise, Least Significant Byte (LSB) first.

Example:

CRC-CCITT-16 with test data of "123456789" becomes 29B1h

Expected Response (3 words):

QE_Code: 0x1C00 Length: 0x0003 CRC Value: 0xXXXX

6.2.4.10 QBLANK Command

 15
 12
 11
 0

 Opcode
 Length

 Reserved
 Size_MSB

 Size_LSW

 Reserved
 Addr_MSB

 Addr_LSW

Table 6-11 shows the description for the \mathtt{QBLANK} command.

TABLE 6-11: COMMAND DESCRIPTION

Field	Description					
Opcode	0xE					
Length	0x5					
Size	Length of program memory to check (in 24-bit words) + Addr_MS					
Addr_MSB	Most Significant Byte of the 24-bit address					
Addr_LSW	Least Significant 16 bits of the 24-bit address					

The QBLANK command queries the PE to determine if the contents of code memory are blank (contains all '1's). The size of code memory to check must be specified in the command.

The Blank Check for code memory begins at [Addr] and advances toward larger addresses for the specified number of instruction words.

QBLANK returns a QE_Code of F0h if the specified code memory is blank; otherwise, QBLANK returns a QE_Code of 0Fh.

Expected Response (2 words for blank device):

0x1EF0 0x0002

Expected Response (2 words for non-blank device):

0x1E0F 0x0002

Note: The QBLANK command does not check the system operation Configuration bits, since these bits are not set to '1' when a Chip Erase is performed.

6.3 Programming Executive Responses

The PE sends a response to the programmer for each command that it receives. The response indicates if the command was processed correctly. It includes any required response data or error data.

The PE response set is shown in Table 6-12. This table contains the opcode, mnemonic and description for each response. The response format is described in **Section 6.3.1** "Response Format".

TABLE 6-12: PROGRAMMING EXECUTIVE RESPONSE OPCODES

Opcode	Mnemonic	Description							
0x1	PASS	Command successfully processed.							
0x2	FAIL	Command unsuccessfully processed.							
0x3	NACK	Command not known.							

6.3.1 RESPONSE FORMAT

All PE responses have a general format, consisting of a two-word header and any required data for the command.

15	12	11	8	7		0		
Орс	ode	Last	_Cmd		QE_Code			
	Length							
	D_1 (if applicable)							
	D_N (if applicable)							

Table 6-13 shows the description of the response format.

TABLE 6-13: RESPONSE FORMAT DESCRIPTION

Field	Description
Opcode	Response opcode.
Last_Cmd	Programmer command that generated the response.
QE_Code	Query code or error code.
Length	Response length in 16-bit words (includes 2 header words).
D_1	First 16-bit data word (if applicable).
D_N	Last 16-bit data word (if applicable).

6.3.1.1 Opcode Field

The opcode is a 4-bit field in the first word of the response. The opcode indicates how the command was processed (see Table 6-12). If the command was processed successfully, the response opcode is PASS. If there was an error in processing the command, the response opcode is FAIL and the QE_Code indicates the reason for the failure. If the command sent to the PE is not identified, the PE returns a NACK response.

6.3.1.2 Last Cmd Field

The Last_Cmd is a 4-bit field in the first word of the response and indicates the command that the PE processed. Since the PE can only process one command at a time, this field is technically not required. However, it can be used to verify that the PE correctly received the command that the programmer transmitted.

6.3.1.3 QE Code Field

The QE_Code is a byte in the first word of the response. This byte is used to return data for query commands and error codes for all other commands.

When the PE processes one of the two query commands (QBLANK or QVER), the returned opcode is always PASS and the QE_Code holds the query response data. The format of the QE_Code for both queries is shown in Table 6-14.

TABLE 6-14: QE Code FOR QUERIES

Query	QE_Code					
QBLANK	0x0F = Code memory is NOT blank 0xF0 = Code memory is blank					
QVER	0xMN, where PE Software Version = M.N (i.e., 0x32 means Software Version 3.2).					

When the PE processes any command other than a query, the QE_Code represents an error code. Supported error codes are shown in Table 6-15. If a command is successfully processed, the returned QE_Code is set to 0x0, which indicates that there is no error in the command processing. If the verify of the programming for the PROGW command fails, the QE_Code is set to 0x1. For all other PE errors, the QE_Code is set to 0x02.

TABLE 6-15: QE_Code FOR NON-QUERY COMMANDS

QE_Code	Description				
0x0	No error				
0x1	Verify failed				
0x2	Other error				

6.3.1.4 Response Length

The response length indicates the length of the PE's response in 16-bit words. This field includes the 2 words of the response header.

With the exception of the response for the read commands, the length of each response is only 2 words.

The response to the READP commands uses the packed instruction word format described in Section 6.2.2 "Packed Data Format". When reading an odd number of program memory words (N odd), the response to the READP command is (3 * (N + 1)/2 + 2) words. When reading an even number of program memory words (N even), the response to the READP command is (3 * N/2 + 2) words.

7.0 DUAL PARTITION FLASH PROGRAMMING CONSIDERATIONS

The PIC24FJ1024GB610 devices supports a Single Partition Flash mode and two Dual Partition Flash modes. The Dual Partition modes allow the device to be programmed with two separate applications to facilitate bootloading or to allow an application to be programmed at run time without stalling the CPU.

The part's Partition Flash mode is determined by the BTMODE<1:0> bits in the FBOOT Configuration register (Table 7-1). The device will automatically check FBOOT on Reset and determine the appropriate Partition mode.

TABLE 7-1: PARTITION MODE SELECT

FBOOT<1:0>	Partition Mode			
00	Reserved			
01	Protected Dual Partition Flash mode			
10	Dual Partition Flash mode			
11	Single Partition Flash mode (default)			

7.1 Dual Participation Memory Organization

In the Dual Partition modes, the device's memory is divided evenly into two physical sections, known as Partition 1 and Partition 2. Each of these partitions contains its own program memory and Configuration Words. During program execution, the code on only one of these partitions is executed; this is the Active Partition. The other partition, or the Inactive Partition, is not used, but can be programmed.

The Active Partition is always mapped to logical address, 000000h, while the Inactive Partition will always be mapped to logical address, 400000h. Note that even when the code partitions are switched between Active and Inactive by the user, the address of the Active Partition will still be 000000h and the address of the Inactive Partition will still be at 400000h.

The Boot Sequence Configuration Word (FBTSEQ) determines whether Partition 1 or Partition 2 will be active after Reset. If the part is operating in Dual Partition mode, the partition with the lower Boot Sequence Number will operate as the Active Partition (FBTSEQ is unused in Single Partition mode). The partitions can be switched between Active and Inactive by reprogramming their Boot Sequence Numbers, but the Active Partition will not change until a device Reset is performed. If both Boot Sequence Numbers are the same, or if both are corrupted, the part will use Partition 1 as the Active Partition. If only one Boot Sequence Number is corrupted, the device will use the partition without a corrupted Boot Sequence Number as the Active Partition.

The user can also change which partition is active at run time using the BOOTSWP instruction. Issuing a BOOTSWP instruction does not affect which partition will be the Active Partition after a Reset. Figure 7-1 demonstrates how the relationship between Partitions 1 and 2, shown in red and blue, respectively, and the Active and Inactive Partitions are affected by reprogramming the Boot Sequence Number or issuing a BOOTSWP instruction.

The P2ACTIV bit (NVMCON<10>) can be used to determine which physical partition is the Active Partition. If P2ACTIV = 1, Partition 2 is Active; if P2ACTIV = 0, Partition 1 is the Active Partition.

7.2 Erase Operations with Dual Partition

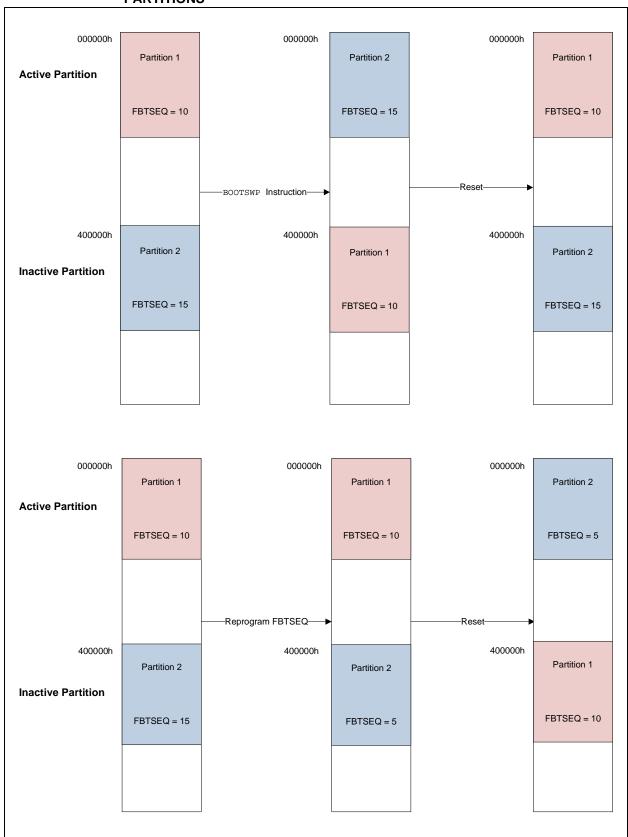
PIC24FJ1024GB610 devices support three erase operations: Chip Erase, Inactive Partition Erase and Page Erase.

A Chip Erase operation erases all of user memory, including the Flash Configuration Words and the FBOOT Configuration register. This resets the Partition mode of the device to its default, Single Partition mode. Chip Erase is not available during run-time operation.

An Inactive Partition Erase operation can be executed at run time from the Active Partition. It will erase all user memory and Flash Configuration Words in the Inactive Partition. Note that the Inactive Partition Erase command is only functional when the device is in one of the Dual Partition modes.

Since the Flash Configuration Words reside in the last locations of user program space, they may be erased using a Chip Erase, an Inactive Partition Erase or a Page Erase that targets the last page of user Flash memory.

FIGURE 7-1: RELATIONSHIP BETWEEN PARTITIONS 1 AND 2 AND ACTIVE/INACTIVE PARTITIONS



7.3 Dual Partition Configuration Words

In Dual Partition modes, each partition has its own set of Flash Configuration Words. The full set of Configuration registers in the Active Partition is used to determine the device's configuration; the Configuration Words in the Inactive Partition are used to determine the device's configuration when that partition becomes active. However, some of the Configuration registers in the Inactive Partition (FSEC, FBSLIM and FSIGN) may be used to determine how the Active Partition is able or allowed to access the Inactive Partition.

7.4 Programming in Dual Partition Mode

When programming a PIC24FJ1024GB610 device with multiple applications, the following sequence of steps is recommended to program the device correctly:

- 1. Perform a Chip Erase on the device.
- Program the FBOOT register to one of the Dual Partition modes, then reset the device to put it in Dual Partition mode. At this point, the device's program memory is functionally divided into two partitions.
- 3. Program the first application (including its Boot Sequence Number and Configuration Words) into the Active Partition at address, 000000h.
- Program the second application, its Boot Sequence Number and its Configuration Words into the Inactive Partition at address, 400000h.
- Finally, perform a device Reset. The Active/ Inactive Partitions may switch at this point, depending on the values that the user used for the Boot Sequence Numbers for each application.

8.0 DEVICE ID

The Device ID region of memory can be used to determine variant and manufacturing information about the chip. This region of memory is read-only and can be read when code protection is enabled. The DEVID register (FF0000h) identifies the specific part number of the device, while DEVREV (FF0002h) shows the silicon revision level.

Table 8-1 lists the identification information for each device. Table 8-2 shows the Device ID registers and Table 8-3 describes the bit field of each register.

TABLE 8-1: DEVICE IDs

Device	DEVID
PIC24FJ128GA606	6000h
PIC24FJ256GA606	6008h
PIC24FJ512GA606	6010h
PIC24FJ1024GA606	6018h
PIC24FJ128GA610	6001h
PIC24FJ256GA610	6009h
PIC24FJ512GA610	6011h
PIC24FJ1024GA610	6019h
PIC24FJ128GB606	6004h
PIC24FJ256GB606	600Ch
PIC24FJ512GB606	6014h
PIC24FJ1024GB606	601Ch
PIC24FJ128GB610	6005h
PIC24FJ256GB610	600Dh
PIC24FJ512GB610	6015h
PIC24FJ1024GB610	601Dh

TABLE 8-2: PIC24FJ1024GA610/GB610 FAMILY DEVICE ID REGISTERS

Address	Name		Bit														
Address	Name	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
FF0000h	DEVID		FAMID<7:0> DEV<					/<7:0>									
FF0002h	DEVREV		_					REV.	<3:0>								

TABLE 8-3: DEVICE ID BIT FIELD DESCRIPTIONS

Bit Field	Register	Description		
FAMID<7:0>	DEVID Encodes the family ID of the device.			
DEV<7:0>	DEVID	Encodes the individual ID of the device.		
REV<3:0>	DEVREV	Encodes the sequential (numerical) revision identifier of the device.		

8.1 Unique Device Identifier (UDID)

All PIC24FJ1024GA610/GB610 family devices are individually encoded during final manufacturing with a Unique Device Identifier, or UDID. The UDID cannot be erased by a Bulk Erase command or any other user-accessible means. This feature allows for manufacturing traceability of Microchip Technology devices in applications where this is a requirement. It may also be used by the application manufacturer for any number of things that may require unique identification, such as:

- Tracking the device
- Unique serial number
- Unique security key

The UDID comprises five 24-bit program words. When taken together, these fields form a unique 120-bit identifier.

The UDID is stored in five read-only locations, located between 801600h and 801608h in the device Configuration space. Table 8-4 lists the addresses of the identifier words and shows their contents.

TABLE 8-4: UDID ADDRESSES

UDID	Address	Description
UDID1	801600	UDID Word 1
UDID2	801602	UDID Word 2
UDID3	801604	UDID Word 3
UDID4	801606	UDID Word 4
UDID5	801608	UDID Word 5

9.0 CHECKSUM COMPUTATION

Checksums for devices are 16 bits in size. The checksum is calculated by summing the following:

- · Contents of code memory locations
- Contents of Flash Configuration Words

All memory locations, including Configuration Words, are summed by adding all three bytes of each memory address.

Since the configuration memory space in this device is Flash-based, the erased value of all memory bits, except FSIGN, is a '1'. Certain Configuration bits will be reserved and these bits may require a '0' or a '1' to be written in order for the device to operate properly. If a Configuration bit is defaulted to a '0' (for example, JTAGEN), the erased state is still a '1'. The bit is set to '0' by the device programmer and is masked when computing the checksum. Therefore, the term, 'blank checksum', when used by MPLAB X IDE or other programmers, must take masked Configuration bits into account. There is no error checking in either hardware or in MPLAB® X IDE to prevent the user from improperly programming these bits. The device checksum is calculated based on the contents of the Hex file, even if the reserved bits are left unprogrammed.

This device uses a Dual Partition Flash memory whose behavior is controlled by the FBOOT register. Table 9-1 shows a Single Partition set of checksums and Table 9-2 shows a Dual Partition set of checksums. The difference in the two is the inclusion of a second FSIGN register in Dual Partition mode and the values for FBOOT; this will alter the computed checksum.

The entire Flash-based Configuration Word area is used when computing the checksum, with the exception of the Configuration Word, FBTSEQ. FBTSEQ is not used in the checksum calculation since it can be altered by a bootloader.

The checksum is computed "bytewise", with the final result truncated to 16 bits. In the PIC24 architecture, each Flash memory address contains two bytes (if an even address) or one byte (if an odd address, since the upper byte is always 0x00). When computing the checksum, both the upper and lower bytes of the word at a given address should be added to the running sum separately as bytes, instead of as a single 16-bit word.

For example, in a program that contains two words with contents such as:

Contents at address 0x0000 = 0xABCD Contents at address 0x0001 = 0x00EF

The checksum over the [0x0000:0x0001] region is: 0xCD + 0xAB + 0xEF + 0x00 = 0x0267.

The CFGB block checksum is also a "bytewise" sum of all contents of the configuration block address region, and is computed identically to the PROG region, with the exception of the FBTSEQ, FICD and FSIGN Configuration Words. The FSIGN and FICD Configuration Words contain a reserved bit that should be AND masked out and excluded during the checksum computation.

The checksum mask values are as follows:

- FSIGN FF7FFFh, which masks bit 15
- FICD FFFFDFh, which masks bit 5 (JTAGEN bit)

TABLE 9-1: CHECKSUM COMPUTATION (SINGLE PARTITION)

Device	Read Code Protection	Checksum Computation	Blank Value	Value with 0xAAAAAA at 0x0 and Last Code Address
PIC24FJ1024GB610	Disabled	PROG[0:ABEFEh] + CFGB[ABF00h:ABFFEh]	DC63h	DA65h
	Enabled	0	0000h	0000h
PIC24FJ512GB610	Disabled	PROG[0:55EFEh] + CFGB[55F00h:55FFEh]	EC63h	EA65h
	Enabled	0	0000h	0000h
PIC24FJ256GB610	Disabled	PROG[0:2AEFEh] + CFGB[2AF00h:2AFFEh]	F463h	F265h
	Enabled	0	0000h	0000h
PIC24FJ128GB610	Disabled	PROG[0:15EFEh] + CFGB[15F00h:15FFEh]	EC63h	EA65h
	Enabled	0	0000h	0000h

Legend: PROG[a:b] = Program memory byte sum of locations, a to b inclusive (all 3 bytes of code memory)

CFGB[c:d] = Configuration memory byte sum of locations, c to d inclusive, with FSIGN<15> and FICD<5>
masked, and FBTSEQ omitted

TABLE 9-2: CHECKSUM COMPUTATION (DUAL PARTITION)

Device	Read Code Protection	Checksum Computation	Blank Value	Value with 0xAAAAAA at 0x0 and Last Code Address
PIC24FJ1024GB610	Disabled	PROG[0:55EFEh] + CFGB[55F00h:55FFEh] + PROG[400000:455EFEh] + CFGB[455F00h:455FFEh]	D8C6h	D4CAh
	Enabled	0	0000h	0000h
PIC24FJ512GB610	Disabled	PROG[0:2AEFEh] + CFGB[2AF00h:2AFFEh] + PROG[400000:42AEFEh] + CFGB[42AF00h:42AFFEh]	E8C6h	E4CAh
	Enabled	0	0000h	0000h
PIC24FJ256GB610	Disabled	PROG[0:156FEh] + CFGB[15700h:157FEh] + PROG[400000:4156FEh] + CFGB[415700h:4157FEh]	F0C6h	ECCAh
	Enabled	0	0000h	0000h
PIC24FJ128GB610	Disabled	PROG[0:0AEFEh] + CFGB[0AF00h:0AFFEh] + PROG[400000:40AEFEh] + CFGB[40AF00h:40AFFEh]	E8C6h	E4CAh
	Enabled	0	0000h	0000h

Legend: PROG[a:b] = Program memory byte sum of locations, a to b inclusive (all 3 bytes of code memory)

CFGB[c:d] = Configuration memory byte sum of locations, c to d inclusive, with FSIGN<15> and FICD<5>
masked, and FBTSEQ omitted

10.0 AC/DC CHARACTERISTICS AND TIMING REQUIREMENTS

Table 10-1 lists the AC/DC characteristics and timing requirements.

TABLE 10-1: AC/DC CHARACTERISTICS AND TIMING REQUIREMENTS

Standard Operating Conditions

Operating Temperature: -40°C to +85°C. Programming at +25°C is recommended.

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Param. No.	Symbol	Characteristic	Min.	Max.	Units	Conditions	
D111	VDD	Supply Voltage During Programming	2.0	3.6	V	See Note 1 and Note 2	
D113	IDDP	Supply Current During Programming	_	8	mA	See Note 2	
D114	IPEAK	Instantaneous Peak Current During Start-up	_	_	mA	See Note 2	
D031	VIL	Input Low Voltage	_	0.2 VDD	V	See Note 2	
D041	VIH	Input High Voltage	0.8 VDD	VDD	V	See Note 2	
D080	Vol	Output Low Voltage	_	0.8	V	See Note 2	
D090	Vон	Output High Voltage	0.8 VDD	VDD	V	See Note 2	
D012	Сю	Capacitive Loading on I/O Pin (PGEDx)	_	50	pF	See Note 2	
P1	TPGC	Serial Clock (PGECx) Period (ICSP™)	200	_	ns		
P1	TPGC	Serial Clock (PGECx) Period (Enhanced ICSP)	500	_	ns		
P1A	TPGCL	Serial Clock (PGECx) Low Time (ICSP)	80	_	ns		
P1A	TPGCL	Serial Clock (PGECx) Low Time (Enhanced ICSP)	200	_	ns		
P1B	Трдсн	Serial Clock (PGECx) High Time (ICSP)	80	_	ns		
P1B	TPGCH	Serial Clock (PGECx) High Time (Enhanced ICSP)	200	_	ns		
P2	TSET1	Input Data Setup Time to Serial Clock ↓	15	_	ns		
P3	THLD1	Input Data Hold Time from PGECx ↓	15	_	ns		
P4	TDLY1	Delay Between 4-Bit Command and Command Operand	40	_	ns		
P4A	TDLY1A	Delay Between Command Operand and Next 4-Bit Command	40	_	ns		
P5	TDLY2	Delay Between Last PGECx ↓ of Command to First PGECx ↑ of Read of Data Word	20	_	ns		
P6	TSET2	VDD ↑ Setup Time to MCLR ↑	100	_	ns		
P7	THLD2	Input Data Hold Time from MCLR ↑	50	_	ms		
P8	TDLY3	Delay Between Last PGECx ↓ of Command Byte to PGEDx ↑ by PE	12	_	μS		

Note 1: VDD must also be supplied to the AVDD pins during programming. AVDD and AVss should always be within ±0.3V of VDD and Vss, respectively.

^{2:} Time depends on the FRC accuracy and the value of the FRC Oscillator Tuning register. Refer to the "Electrical Characteristics" chapter in the specific device data sheet.

^{3:} This time applies to Program Memory Words, Configuration Words and User ID Words.

TABLE 10-1: AC/DC CHARACTERISTICS AND TIMING REQUIREMENTS (CONTINUED)

Standard Operating Conditions

Operating Temperature: -40°C to +85°C. Programming at +25°C is recommended.

Param. No.	Symbol	Characteristic	Min.	Max.	Units	Conditions
P9A	TDLY4	PE Command Processing Time	10	_	μS	
P9B	TDLY5	Delay Between PGEDx ↓ by PE to PGEDx Released by PE	15	23	μS	
P10	TDLY6	PGECx Low Time After Programming	400	_	ns	
P11	TDLY7	Chip Erase Time	16	25	ms	
P12	TDLY8	Page Erase Time	16	25	ms	See Note 2
P13	TDLY9	Double-Word Programming Time	16	20	μS	See Note 2 and Note 3
P14	TR	MCLR Rise Time to Enter ICSP mode	_	1.0	μS	
P15	TVALID	Data Out Valid from PGECx ↑	10	_	ns	
P16	TDLY10	Delay Between Last PGECx \downarrow and $\overline{\text{MCLR}}$ \downarrow	0	_	S	
P17	THLD3	MCLR ↓ to VDD ↓	100	_	ns	
P18	TKEY1	Delay from First MCLR ↓ to First PGECx ↑ for Key Sequence on PGEDx	1	_	ms	
P19	TKEY2	Delay from Last PGECx ↓ for Key Sequence on PGEDx to Second MCLR ↑	25	_	ns	
P21	TMCLRH	MCLR High Time	_	500	μS	

Note 1: VDD must also be supplied to the AVDD pins during programming. AVDD and AVSS should always be within ±0.3V of VDD and VSS, respectively.

^{2:} Time depends on the FRC accuracy and the value of the FRC Oscillator Tuning register. Refer to the "**Electrical Characteristics**" chapter in the specific device data sheet.

^{3:} This time applies to Program Memory Words, Configuration Words and User ID Words.

APPENDIX A: SPI PORT SCANNING SOFTWARE ROUTINE

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```
/***********************
   © 2015 Microchip Technology Inc.
  FileName:
                 spi_search.s
  Dependencies: none
  Processor:
                 PIC24
  Compiler:
                  XC16
                  MPLAB® X
Description:
                 This file performs a search of the correct
                 PGEC/PGED port and connect SPI2 slave to it.
*************************
/***********************
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************************
```

APPENDIX A: SPI PORT SCANNING SOFTWARE ROUTINE (CONTINUED)

```
.global
      _sd_Port_Search
.section.text
_sd_Port_Search:
  ; Unlock PPS
  ; place zero in w0 to unlock PPS
     #0, w0
  mov
  ; OSCCONL (low byte) unlock sequence
     #OSCCONL, w1
  mov
  mov
      #0x46, w2
  mov
     #0x57, w3
  mov.b w2, [w1]
                  ; write 0x46
  mov.b w3, [w1]
                 ; write 0x9A
  ; unlock PPS (IOLOCK bit = 0)
  mov.b w0, [w1]
; Initialize SPI1 as master and SPI2 as slave
  ; for loopback test on each programming port
  SPI1BRGL
                  ; clear SPI1BRGL
  clr
  clr SPI1CON1H
                 ; clear SPI1CON1H
  clr SPI1CON1L
                 ; clear SPI1CON1L
  mov #0x8120, w0
                 ; set SPIEN, CKE and MSTEN (master)
     w0, SPI1CON1L
  mov
  clr SPI2CON1H
                 ; clear SPI2CON1H
  clr SPI2CON1L
                 ; clear SPI2CON1L
  mov #0x0100, w0
                 ; set CKE (slave)
  mov w0, SPI2CON1L
; Connect SPI2 to the 1st port,
  ; if loopback test works then port is found.
  mov \#0x0100, w0 ; SCK2R to RP1 and SDI2R to RP0
     w0, RPINR22
  mov
                  ; connect SPI2
  mov
     #0x0b0a, w0
  mov w0, RPOR0
  rcall _sd_Loopback_Test
     #1, w1
  mov
  cp0
     wΩ
  bra
     z, port_found
; Connect SPI2 to the 2nd port,
  ; if loopback test works then port is found.
  mov \#0x0607, w0 ; SCK2R to RP6 and SDI2R to RP7
     w0, RPINR22
                 ; connect SPI2
     #0x0a0b, w0
  mov
     w0, RPOR3
  rcall _sd_Loopback_Test
  mov
     #2, w1
  cp0
     wΩ
     z, port_found
  bra
```

APPENDIX A: SPI PORT SCANNING SOFTWARE ROUTINE (CONTINUED)

```
; Connect SPI2 to the 3rd port,
  ; if loopback test works then port is found.
  mov \#0x121c, w0 ; SCK2R to RP18 and SD12R to RP28 mov \#0, RPINR22 ; connect SP12
     #0x0b, w0
  mov
     w0, RPOR9
  mov
  mov
     #0x0a, w0
  mov w0, RPOR14
  rcall _sd_Loopback_Test
     #3, w1
  mov
  cp0 w0
  {\tt bra} {\tt z, port\_found}
  bclr SPI1CON1L, #15 ; disable SPI1
  clr w0
  return
port_found:
 bclr SPI1CON1L, #15 ; disable SPI1
     w1, w0
  return
_sd_Loopback_Test:
  bset SPI2CON1L, #15
                ; enable SPI2
     #0x00a5, w0
                ; 0xa5 pattern
  mov
     w0, SPI1BUFL
  mov
                 ; send
wait data:
 btss SPI1STATL, #0
                ; wait for the data
  bra wait_data
  mov SPI1BUFL, w1
  mov SPI2BUFL, w1
                ; read data
     w0, w1
  CD
  bra z, found
     #1, w0
  mov
  bclr SPI2CON1L, #15 ; disable SPI2
 return
found:
 clr w0
 bclr SPI2CON1L, #15 ; disable SPI2
.end
```

APPENDIX B: REVISION HISTORY

Revision A (April 2014)

Original version of the programming specification.

Revision B (February 2015)

The terms for Dual Boot and Single Boot modes have changed to Dual Partition and Single Partition modes, respectively. All checksum values have changed and there are minor edits throughout text.

Revision C (November 2015)

Updated Table 2-2, Table 2-3, Table 3-3, Table 5-3, Table 6-1, Table 6-6.

In Section 6.2.4.2 "READC Command", the Expected Response was changed to (2 + N) words.

In Section 6.2.4.2 "READC Command", the Expected Response for a blank device was changed to 0x1EF0, and the Expected Response for a non-blank device was changed to 0x1E0F.

Added Section 6.1.3 "SPI Port Scanning", Section 8.1 "Unique Device Identifier (UDID)" and Appendix A: "SPI Port Scanning Software Routine".

Revision D (January 2016)

Updated Table 2-4, Table 9-1 and Table 9-2.

Updated Section 9.0 "Checksum Computation".

Revision E (March 2016)

Updated Section 9.0 "Checksum Computation".

Updated Table 9-1.

Revision F (September 2016)

Revised last sentence in Section 4.4 "Entering Enhanced ICSP Mode".

Updated Figure 4-3.

NOTES:			

Note the following details of the code protection feature on Microchip devices:

- Microchip products meet the specification contained in their particular Microchip Data Sheet.
- Microchip believes that its family of products is one of the most secure families of its kind on the market today, when used in the
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