
**dsPIC33CDVL64MC106 Family
Flash Programming Specification**

1.0 DEVICE OVERVIEW

This document defines the programming specification for the dsPIC33CDVL64MC106 Digital Signal Controller (DSC) device. This programming specification is required only for those developing programming support for the dsPIC33CDVL64MC106 device.

Customers only using this device for application development should use development tools that already provide support for device programming.

Topics covered include:

- [Section 1.0 “Device Overview”](#)
- [Section 2.0 “Programming Overview”](#)
- [Section 3.0 “Device Programming – ICSP”](#)
- [Section 4.0 “Device Programming – Enhanced ICSP”](#)
- [Section 5.0 “The Programming Executive”](#)
- [Section 6.0 “Device ID/Unique ID”](#)
- [Section 7.0 “CRC Checksum Computation”](#)
- [Section 8.0 “AC/DC Characteristics and Timing Requirements”](#)

2.0 PROGRAMMING OVERVIEW

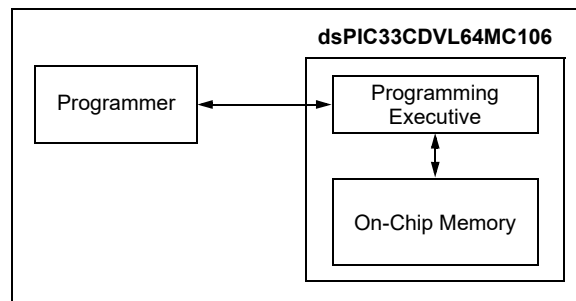
The following are the two methods of programming that are discussed in this programming specification:

- In-Circuit Serial Programming™ (ICSP™)
- Enhanced In-Circuit Serial Programming

The ICSP programming method is the most direct method to program the device; however, it is also the slower of the two methods. It provides native, low-level programming capability to erase, program and verify the device.

The Enhanced ICSP protocol uses a faster method that takes advantage of the Programming Executive (PE), as illustrated in [Figure 2-1](#). The PE provides all the necessary functionality to erase, program and verify the chip through a small command set. The command set allows the programmer to program a dsPIC33CDVL64MC106 device without dealing with the low-level programming protocols.

FIGURE 2-1: PROGRAMMING SYSTEM OVERVIEW FOR ENHANCED ICSP™



This programming specification is divided into two major sections that describe the programming methods independently. [Section 3.0 “Device Programming – ICSP”](#) describes the ICSP method. [Section 4.0 “Device Programming – Enhanced ICSP”](#) describes the Enhanced ICSP method.

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2.1 Required Connections

This device requires specific connections for programming to take place. These connections include power, $\overline{\text{MCLR}}$ and one programming pair (PGEDx/PGECx). Table 2-1 describes these connections (refer to the specific device data sheet for pin descriptions and power connection requirements).

2.2 Power Requirements

The dsPIC33CDVL64MC106 device powers its core digital logic at a nominal 1.2V and incorporates a capacitor-free, on-chip regulator that allows the device to run its core logic from VDD. The regulator provides power to the core from the other VDD pins and does not require an external CPU Logic Filter Capacitor (VCAP) connection. The dsPIC33CDVL64MC106 device also incorporates an on-chip 3.3V regulator. This regulator outputs 3.3V on the VREG pin when 6-28V are supplied to the HVDD pin.

VDD pins may be powered by either an external power supply or the VREG output pin. If an external power supply is used to power the VDD pins directly, the HVDD pin does not need to be powered to program the dsPIC33CDVL64MC106 device. The specifications for core voltage and capacitance are listed in Section 8.0 “AC/DC Characteristics and Timing Requirements”.

FIGURE 2-2: CONNECTIONS FOR THE ON-CHIP REGULATOR

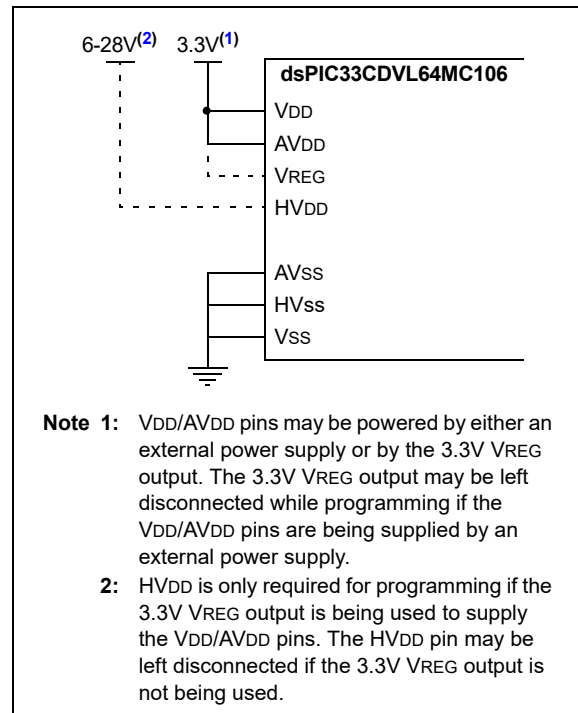


TABLE 2-1: PINS USED DURING PROGRAMMING

| Pin Name | Pin Type | Pin Description |
|-----------------------------|----------|--|
| $\overline{\text{MCLR}}$ | I | Programming Enable |
| HVDD | P | High-Voltage Power Supply ⁽¹⁾ |
| VREG | P | 3.3V Regulator Output |
| VDD and AVDD ⁽¹⁾ | P | Power Supply ⁽²⁾ |
| VSS and AVSS ⁽¹⁾ | P | Ground ⁽²⁾ |
| PGECx | I | Programming Pin Pair: Serial Clock |
| PGEDx | I/O | Programming Pin Pair: Serial Data |

Legend: I = Input O = Output P = Power

Note 1: HVDD and VREG may be left disconnected during programming if the 3.3V VREG output is not being used to supply the VDD and AVDD pins.

2: All VDD, VSS, AVDD and AVSS power supply and ground pins must be connected during programming.

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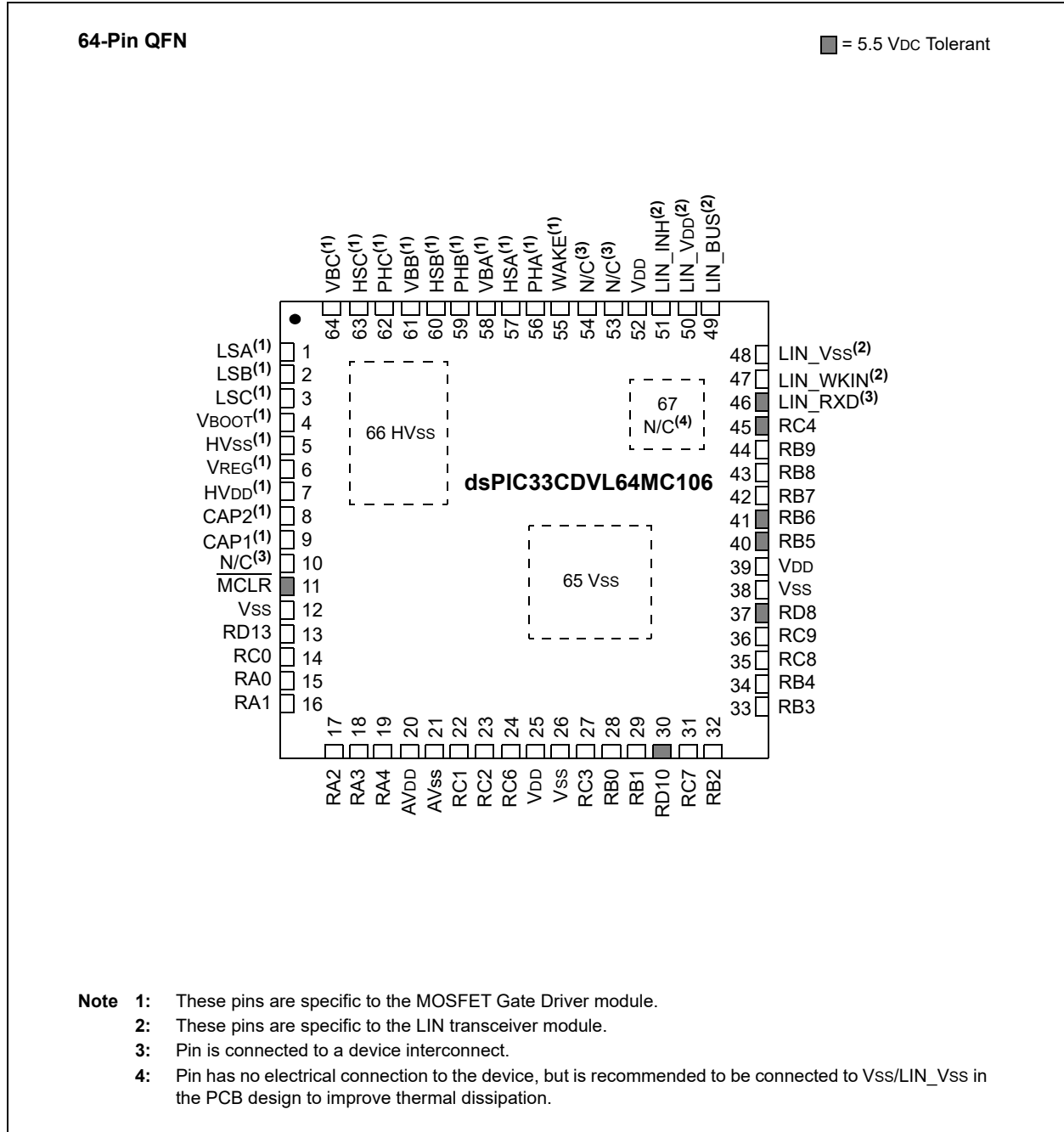
2.3 Pin Diagrams

Figure 2-3 provides the pin diagram for the dsPIC33CDVL64MC106 device. The pins that are required for programming are listed in Table 2-1 and are indicated in bold text in the figures. Refer to the specific device data sheet for complete pin descriptions.

2.3.1 PGECx AND PGEDx PIN PAIRS

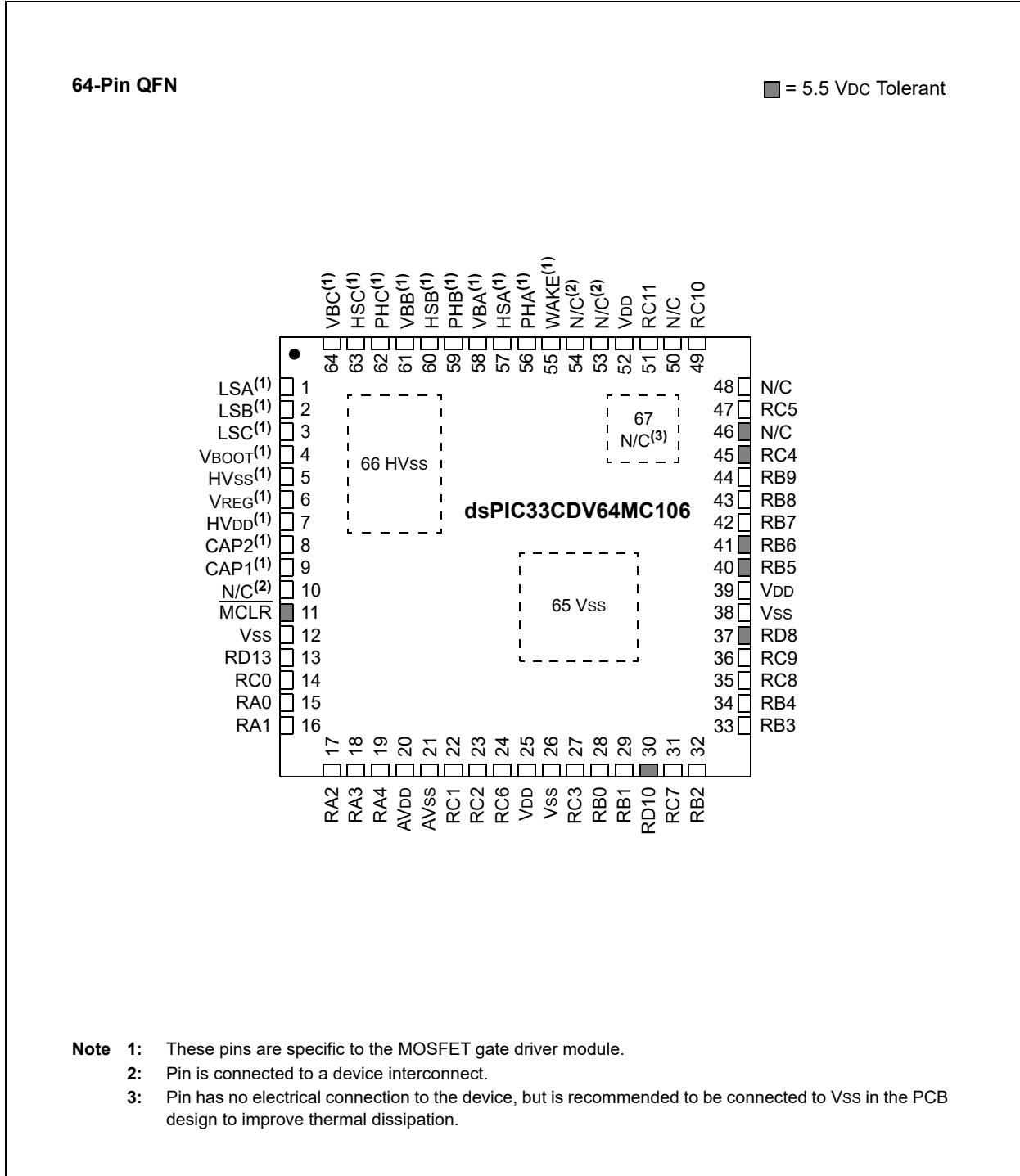
The dsPIC33CDVL64MC106 device has three separate pairs of programming pins, labeled as PGEC1/PGED1, PGEC2/PGED2 and PGEC3/PGED3. Any one of these pin pairs may be used for device programming by either ICSP or Enhanced ICSP. Unlike voltage supply and ground pins, it is not necessary to connect all three pin pairs to program the device. However, the programming method must use both pins of the same pair.

FIGURE 2-3: PIN DIAGRAMS



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FIGURE 2-4: PIN DIAGRAMS (CONTINUED)



2.4 Program Memory Write/Erase Requirements

The program Flash memory has a specific write/erase requirement that must be adhered to for proper device operation. The rule is that any given word in memory must not be written without first erasing the page in which it is located. Thus, the easiest way to conform to this rule is to write all the data in a programming block within one write cycle. The programming methods specified in this document comply with this requirement.

Note: A program memory word can be programmed twice before an erase, but only if (a) the same data are used in both program operations or (b) bits containing '1' are set to '0', but no '0' is set to '1'.

2.5 Memory Map

The program memory map extends from 0x000000 to 0xFFFFFE. User program code storage is located at the base of the memory map. The last row (128 instruction words) of the last page of implemented program memory is reserved for the device Configuration bits. Caution should be used when placing code in the last page of user memory. When programming the device configuration, the whole last page must first be erased, including any code located in the last page.

Table 2-3 lists the user memory address limit and available number of instruction words (including the device configuration area), the number of write blocks and the number of erase blocks present in each device variant.

Locations, 0x800000 through 0x800BFE, are reserved for executive code memory. This region stores the PE and the debugging executive, which is used for device programming. This region of memory cannot be used to store user code. See Section 5.0 “The Programming Executive” for more information.

Locations, 0xFF0000 and 0xFF0002, are reserved for the Device ID Word registers. These bits can be used by the programmer to identify which device type is being programmed. They are described in Section 6.0 “Device ID/Unique ID”. The Device ID registers read out normally, even after code protection is applied.

The locations, 0x801700 to 0x8017FE, are a One-Time-Programmable (OTP) memory area. The user OTP Words can be used for storing product information, such as serial numbers, system manufacturing dates, manufacturing lot numbers and other application-specific information. They are described in Section 2.7 “User One-Time-Programmable (OTP) Memory”.

Figure 2-5 through Figure 2-6 show generic memory maps for the device listed in Table 2-3. See the “Memory Organization” chapter in the specific device data sheet for exact memory addresses.

TABLE 2-2: DEVICE ROW AND PAGE SIZE

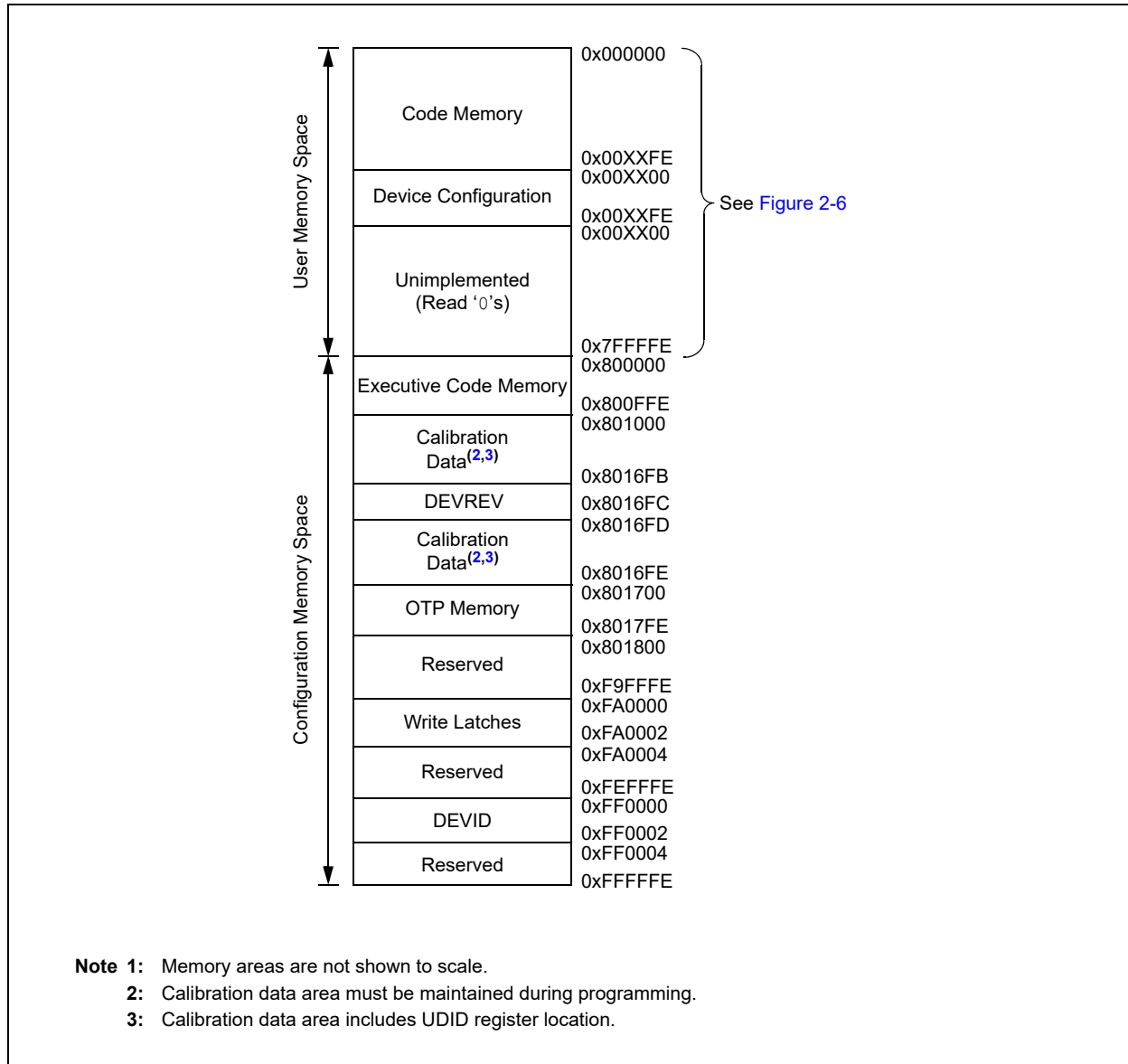
| Instruction Words (24 bits) | |
|-----------------------------|------|
| Row | 128 |
| Page | 1024 |

TABLE 2-3: FLASH CODE MEMORY SIZE

| Device Family | User Memory Limit (Instruction Words) | Write Blocks/ No. of Rows | Erase Blocks/ No. of Pages |
|--------------------|---------------------------------------|---------------------------|----------------------------|
| dsPIC33CDVL64MC106 | 0x00AFFE (22528) | 176 | 22 |
| dsPIC33CDV64MC106 | 0x00AFFE (22528) | 176 | 22 |

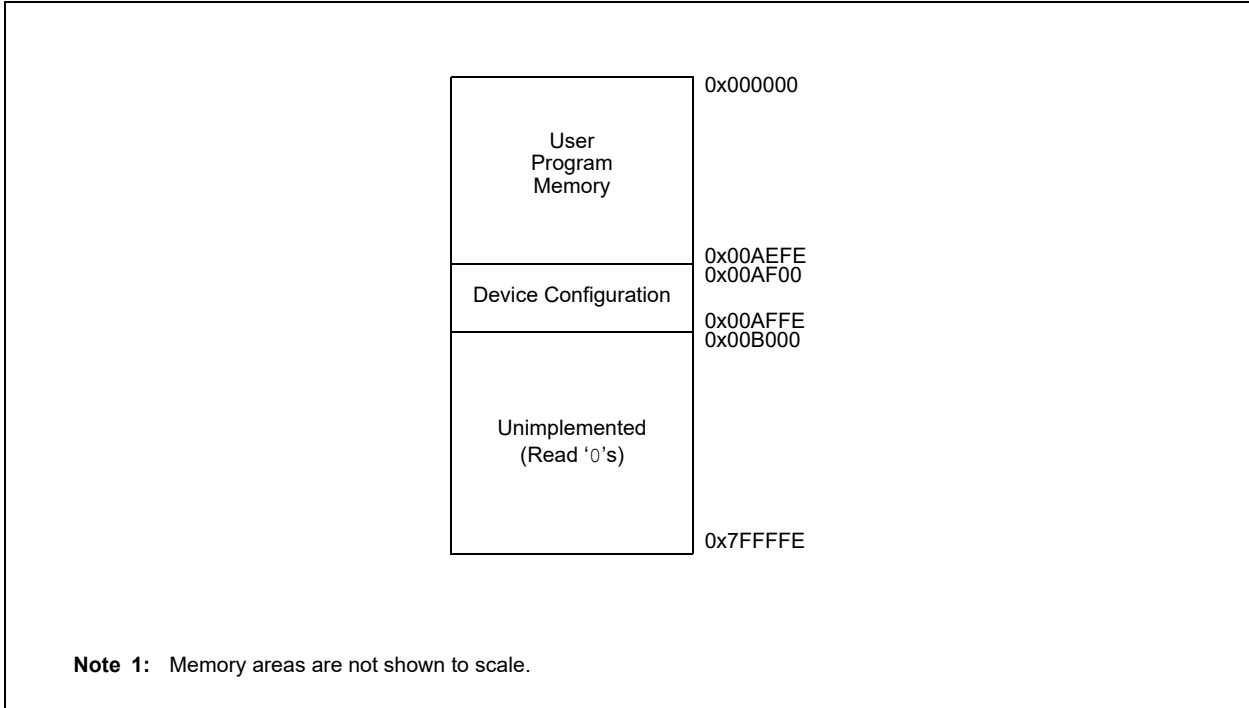
dsPIC33CDVL64MC106 FAMILY

FIGURE 2-5: PROGRAM MEMORY MAP FOR dsPIC33CDVL64MC106 DEVICE⁽¹⁾



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FIGURE 2-6: PROGRAM MEMORY MAP FOR dsPIC33CDVL64MC106 DEVICE⁽¹⁾



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2.6 Configuration Bits

2.6.1 OVERVIEW

The Configuration bits are stored in the last row of the last page location of implemented program memory. These bits can be set or cleared to select various device configurations. There are two types of Configuration bits: system operation bits and code-protect bits. The system operation bits determine the power-on settings for system-level components, such as the oscillator and the Watchdog Timer. The code-protect bits prevent program memory from being read and written.

Table 2-4 list the Configuration register address for each device memory size and partition configuration.

Table 2-5 shows the Configuration register map. Refer to the “**Special Features**” chapter in the specific device data sheet for the full Configuration Word register descriptions for your device.

Special attention should be taken for the reserved Configuration bits. The FSIGN[15], FDEVOPT[8] and FDEVOPT[9] Configuration bits must always be programmed ('0'). The FPOR[5:4], FICD[7], FDEVOPT[7] and FDEVOPT[10] Configuration bits must always stay unprogrammed ('1').

TABLE 2-4: dsPIC33CDVL64MC106 CONFIGURATION ADDRESSES

| Register Name | 64k | 32k |
|---------------|----------|----------|
| FSEC | 0x00AF00 | 0x005F00 |
| FBSLIM | 0x00AF10 | 0x005F10 |
| FSIGN | 0x00AF14 | 0x005F14 |
| FOSCSEL | 0x00AF18 | 0x005F18 |
| FOSC | 0x00AF1C | 0x005F1C |
| FWDT | 0x00AF20 | 0x005F20 |
| FPOR | 0x00AF24 | 0x005F24 |
| FICD | 0x00AF28 | 0x005F28 |
| FDMTIVTL | 0x00AF2C | 0x005F2C |
| FDMTIVTH | 0x00AF30 | 0x005F30 |
| FDMTCNTL | 0x00AF34 | 0x005F34 |
| FDMTCNTH | 0x00AF38 | 0x005F38 |
| FDMT | 0x00AF3C | 0x005F3C |
| FDEVOPT | 0x00AF40 | 0x005F40 |
| FALTREG | 0x00AF44 | 0x005F44 |

TABLE 2-5: CONFIGURATION REGISTERS MAP

| Register Name | Bits 23-16 | Bit 15 | Bit 14 | Bit 13 | Bit 12 | Bit 11 | Bit 10 | Bit 9 | Bit 8 | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | |
|---------------|------------|------------------|------------|---------|-------------|--------------|------------------|------------------|------------------|------------------|------------|------------------|------------------|-------|------------|------------------|--------|---|
| FSEC | — | AIVTDIS | — | — | — | CSS[2:0] | | | CWRP | GSS[1:0] | | GWRP | — | BSEN | BSS[1:0] | | BWRP | |
| FBSLIM | — | — | — | — | BSLIM[12:0] | | | | | | | | | | | | | |
| FSIGN | — | r ⁽²⁾ | — | — | — | — | — | — | — | — | — | — | — | — | — | — | — | |
| FOSCSEL | — | — | — | — | — | — | — | — | — | IESO | — | — | — | — | FNOSC[2:0] | | | |
| FOSC | — | — | — | — | XTBST | XTALCFG[1:0] | | — | PLLKEN | FCKSM[1:0] | | — | — | — | OSCIOFCN | POSCMD[1:0] | | |
| FWDT | — | FWDTEN | SMWDT[4:0] | | | | WDTWIN[1:0] | | WINDIS | WDTCLKS[1:0] | | RMWDT[4:0] | | | | | | |
| FPOR | — | — | — | — | — | — | r ⁽¹⁾ | — | — | — | BISTDIS | r ⁽¹⁾ | r ⁽¹⁾ | — | — | — | — | |
| FICD | — | NOBTSWP | — | — | — | — | — | — | — | r ⁽¹⁾ | — | JTAGEN | — | — | — | ICS[1:0] | | |
| FDMTIVTL | — | DMTIVT[15:0] | | | | | | | | | | | | | | | | |
| FDMTIVTH | — | DMTIVT[31:16] | | | | | | | | | | | | | | | | |
| FDMTCNTL | — | DMTCNT[15:0] | | | | | | | | | | | | | | | | |
| FDMTCNTH | — | DMTCNT[31:16] | | | | | | | | | | | | | | | | |
| FDMT | — | — | — | — | — | — | — | — | — | — | — | — | — | — | — | — | DMTDIS | |
| FDEVOPT | — | — | — | SPI2PIN | — | — | SMB3EN | r ⁽²⁾ | r ⁽²⁾ | r ⁽¹⁾ | — | — | — | — | ALTI2C1 | r ⁽¹⁾ | — | — |
| FALTREG | — | — | CTXT4[2:0] | | | — | CTXT3[2:0] | | | — | CTXT2[2:0] | | | — | CTXT1[2:0] | | | |

Legend: — = unimplemented bit, read as '1'; r = reserved bit.

Note 1: Bit reserved, maintain as '1'.

2: Bit reserved, maintain as '0'.

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2.6.2 CODE-PROTECT CONFIGURATION BITS

The device implements intermediate security features defined by the FSEC register. The Boot Segment (BS) is the highest privileged segment and the General Segment (GS) is the lowest privileged segment. The total user code memory can be split into BS or GS. The size of the segments is determined by the BSLIM[12:0] bits. The relative location of the segments within user space does not change, such that BS (if present) occupies the memory area just after the Vector Space (VS), Interrupt Vector Table (IVT), and the GS occupies the space just after BS (or if the Alternate Interrupt Vector Table (AIVT) is enabled, just after AIVT VS). The Configuration Segment (CS) is a small segment (less than a page, typically just one row) within user Flash address space that contains all user configuration data that are loaded by the NVM controller during the Reset sequence.

2.7 User One-Time-Programmable (OTP) Memory

The dsPIC33CDVL64MC106 device provides 128 words (64 pairs) of One-Time-Programmable (OTP) memory, located at addresses, 0x801700 through 0x8017FE. This memory can be used for persistent storage of application-specific information that will not be erased by reprogramming the device. This includes many types of information, such as:

- Application checksums
- Code revision information
- Product information
- Serial numbers
- System manufacturing dates
- Manufacturing lot numbers

Customer OTP memory may be programmed in any mode, including user RTSP mode, but it cannot be erased. Data are not cleared by a Chip Erase. See [Section 3.9 “Writing OTP Words”](#) for more information regarding OTP programming.

2.8 ICSP Write Inhibit

ICSP Write Inhibit is an access restriction feature that restricts all of Flash memory when activated. Once activated, ICSP Write Inhibit permanently prevents ICSP Flash programming and erase operations, and cannot be deactivated. This feature is intended to prevent alteration of Flash memory contents with behavior similar to One-Time-Programmable (OTP) devices.

RTSP, including erase and programming operations, is not restricted when ICSP Write Inhibit is activated; however, code to perform these actions must be programmed into the device before ICSP Write Inhibit is activated. This allows for a bootloader-type application to alter Flash contents when ICSP Write Inhibit is activated.

Entry into ICSP and Enhanced ICSP modes is not affected by ICSP Write Inhibit. In these modes, it will continue to be possible to read configuration memory space and any user memory space regions, which are not code-protected. With ICSP Write Inhibit, an attempt to set WR (NVMCON[15]) = 1 will maintain WR = 0, and instead, set WRERR (NVMCON[13]) = 1. All Enhanced ICSP erase and programming commands will have no effect, with self-checked programming commands returning a FAIL response opcode (or a PASS if the destination already exactly matched the requested programming data).

Once ICSP Write Inhibit is activated, it is not possible for a device executing in Debug mode to erase/write Flash, nor can a debug tool switch the device to Production mode. ICSP Write Inhibit should therefore only be activated on devices programmed for production.

The JTAG port, when enabled, can be used to map ICSP signals to JTAG I/O pins. All Flash erase/programming operations, initiated via the JTAG port, will therefore also be blocked after activating ICSP Write Inhibit

2.8.1 ACTIVATING ICSP WRITE INHIBIT

Note: It is not possible to deactivate ICSP Write Inhibit.

ICSP Write Inhibit is activated by executing a pair of NVMCON Double-Word Programming commands to save two 16-bit activation values in the configuration memory space. The target NVM addresses and values required for activation are shown in [Table 2-6](#).

TABLE 2-6: ICSP™ WRITE INHIBIT FUSE ADDRESSES AND CODES

| ICSP Write Inhibit Fuse Address | Code Value |
|---------------------------------|------------|
| 0x801028 | 0x006D63 |
| 0x80102C | 0x006870 |

Once both addresses contain their activation values, ICSP Write Inhibit will take permanent effect on the next device Reset. Neither address can be reset, erased or otherwise modified, through any means, after being successfully programmed, even if one of the addresses has not been programmed.

Only the lower 16 data bits stored at the activation addresses are evaluated; the upper 8 bits and second 24-bit word, written by the Double-Word Programming NVMOPx bits, should be '0's. The addresses can be programmed in any order and during separate ICSP/Enhanced ICSP/RTSP sessions, but any attempt to program an incorrect 16-bit value, or use a Row Programming operation to program the values, will be aborted without altering the existing data.

3.0 DEVICE PROGRAMMING – ICSP

ICSP mode is a special programming protocol that allows you to read and write to the device memory of the dsPIC33CDVL64MC106 device. The ICSP mode is the most direct method used to program the device, which is accomplished by applying control codes and instructions, serially to the device, using the PGECx and PGEDx pins. ICSP mode also has the ability to read the contents of the executive memory to determine if the Programming Executive is present, and to write the Programming Executive to executive memory if it is missing, and then, Enhanced ICSP mode will be used.

In ICSP mode, the system clock is taken from the PGECx pin, regardless of the device's Oscillator Configuration bits. All instructions are shifted serially into an internal buffer, then loaded into the Instruction Register (IR) and executed. No program fetching occurs from internal memory. Instructions are fed in 24 bits at a time. PGEDx is used to shift data in, and PGECx is used as both the serial shift clock and the CPU execution clock.

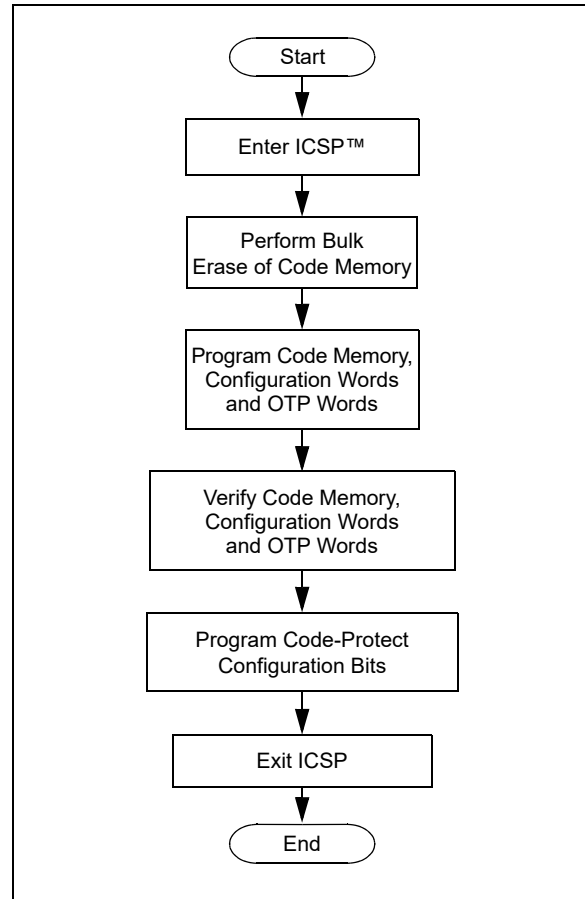
Note 1: During ICSP operation, the operating frequency of PGECx must not exceed 5 MHz.

2: ICSP mode is slower than Enhanced ICSP mode for programming.

3.1 Overview of the Programming Process

Figure 3-1 shows a high-level overview of the ICSP programming process. After entering ICSP mode, the first action is to Bulk Erase the code memory. Next, the code memory is programmed, followed by the device Configuration bits. Code memory (including the Configuration bits) is then verified to ensure that programming was successful. Then, programming the code-protect Configuration bits can be done if required.

FIGURE 3-1: HIGH-LEVEL ICSP™ PROGRAMMING FLOW



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3.2 Entering ICSP Mode

As shown in Figure 3-2, entering ICSP Program/Verify mode requires three steps:

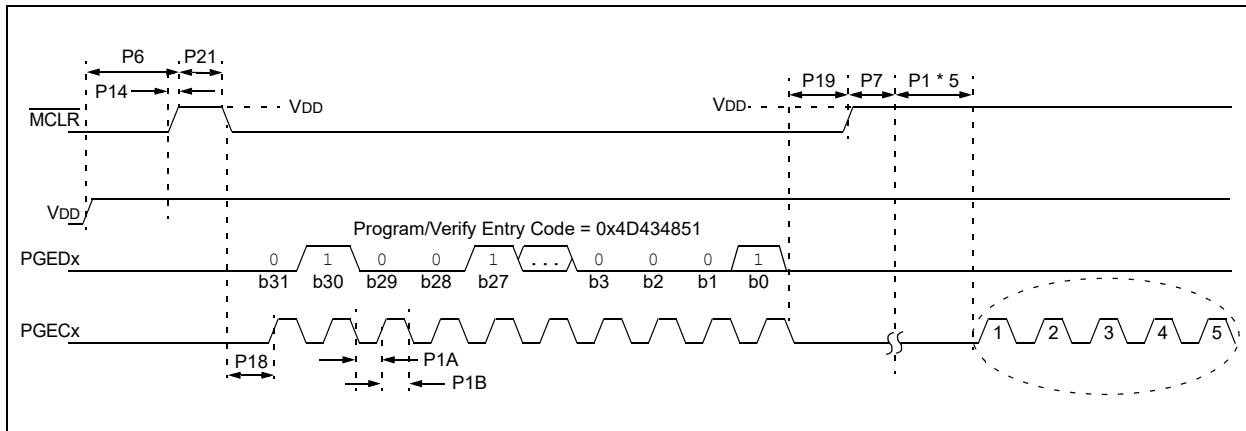
1. $\overline{\text{MCLR}}$ is briefly driven high, then low (P21).
2. A 32-bit key sequence is clocked into PGEDx. An interval of at least P18 must elapse before presenting the key sequence on PGEDx.
3. $\overline{\text{MCLR}}$ is held low during a specified period, P19, and then driven high.
4. After a P7 + 5 * P1 delay, five clock pulses must be generated on the PGECx pin.

The key sequence is a specific 32-bit pattern, '0100 1101 0100 0011 0100 1000 0101 0001' (more easily remembered as 0x4D434851 in hexadecimal). The device will enter ICSP mode only if the sequence is valid. The Most Significant bit (MSb) of the most significant nibble must be shifted in first.

On successful entry, the program memory can be accessed and programmed in serial fashion.

Note: If a capacitor is present on the $\overline{\text{MCLR}}$ pin, the high time for entering ICSP mode can vary.

FIGURE 3-2: ENTERING ICSP™ MODE



3.3 ICSP Operation

Upon entry into ICSP mode, the CPU is Idle. Execution of the CPU is governed by an internal state machine. A 4-bit control code must be clocked in using PGECx and PGEDx, and this control code is used to command the CPU (see Table 3-1).

The SIX control code is used to send instructions to the CPU for execution and the REGOUT control code is used to read data out of the device through the VISI register.

3.3.1 SIX SERIAL INSTRUCTION EXECUTION

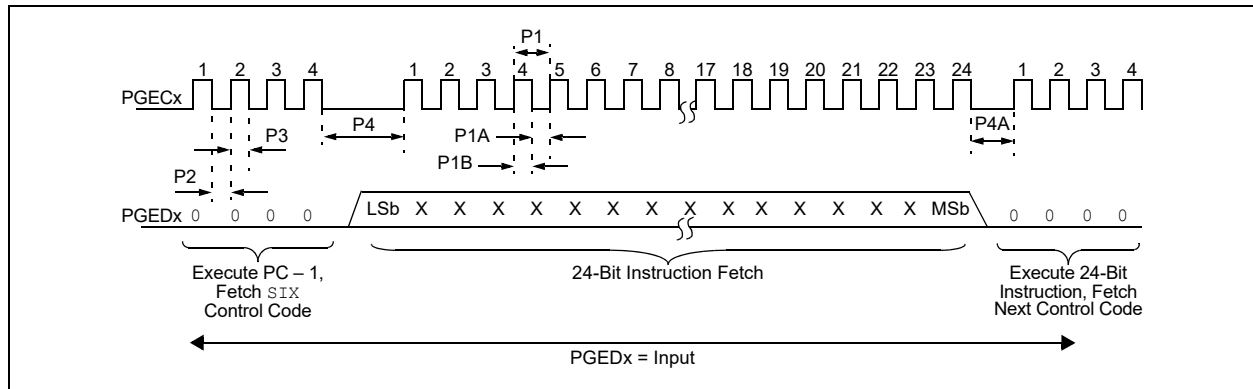
The SIX control code allows execution of the dsPIC33 family assembly instructions. When the SIX code is received, the CPU is suspended for 24 clock cycles, as the instruction is then clocked into the internal buffer. Once the instruction is shifted in, the state machine allows it to be executed over the next four PGECx clock cycles. While the received instruction is executed, the state machine simultaneously shifts in the next 4-bit command (see Figure 3-3).

Note: Data bits on PGEDx are latched on the rising edge of the PGECx clock.

TABLE 3-1: CPU CONTROL CODES IN ICSP™ MODE

| 4-Bit Control Code | Mnemonic | Description |
|--------------------|----------|--|
| 0000 | SIX | Shift in 24-bit instruction and execute. |
| 0001 | REGOUT | Shift out the VISI register. |
| 0010-1111 | N/A | Reserved. |

FIGURE 3-3: SIX SERIAL EXECUTION



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3.3.1.1 Differences Between Execution of SIX and a Normal Instruction

There are some differences between executing instructions normally and using the ICSP `SIX` command. As a result, the code examples in this specification may not match those for performing the same functions during normal device operation.

The important differences are:

- Two-word instructions require two `SIX` operations to clock in all of the necessary data.

Examples of two-word instructions are `GOTO` and `CALL`.

- Two-cycle instructions require two `SIX` operations to complete.

The first `SIX` operation shifts in the instruction and begins to execute it. A second `SIX` operation, which should shift in a `NOP` to avoid losing data, provides the CPU clocks required to finish executing the instruction.

Examples of two-cycle instructions are Table Read (`TBLRD`) and Table Write (`TBLWT`) instructions.

- Must provide `NOP` instruction during Stall to account for pipeline changes.

A CPU Stall occurs when an instruction modifies a register that is used for Indirect Addressing by the instruction immediately following the CPU Stall. During normal operation, the CPU will automatically force a `NOP` while the new data are read. While using ICSP, the CPU stalls under the same conditions, but an instruction needs to be provided to generate the clocks to get through the Stall cycle. Therefore, any indirect references to a recently modified register should be preceded with a `NOP`.

For example, the instructions, `MOV #0x0, W0`, followed by, `MOV[W0], W1`, must have a `NOP` inserted in between.

If a two-cycle instruction modifies a register which is used indirectly, it will require two following `NOPs`: one to execute the second half of the instruction and the other `NOP` stalls the CPU to correct the pipeline.

For example, instructions such as, `TBLWTL [W0++], [W1]`, should be followed by two `NOPs`.

- The device Program Counter (PC) continues to automatically increment during ICSP instruction execution, even though the Flash memory is not being used. As a result, the PC may be incremented so that it points to invalid memory locations.

Examples of invalid memory spaces are unimplemented Flash addresses or the vector space (location: `0x0` to `0x1FF`).

If the PC points to these locations, the device will reset, possibly interrupting the ICSP operation. To prevent this, instructions should be periodically executed to reset the PC to a safe space. The optimal method of achieving this is to perform a "`GOTO 0x200`" instruction.

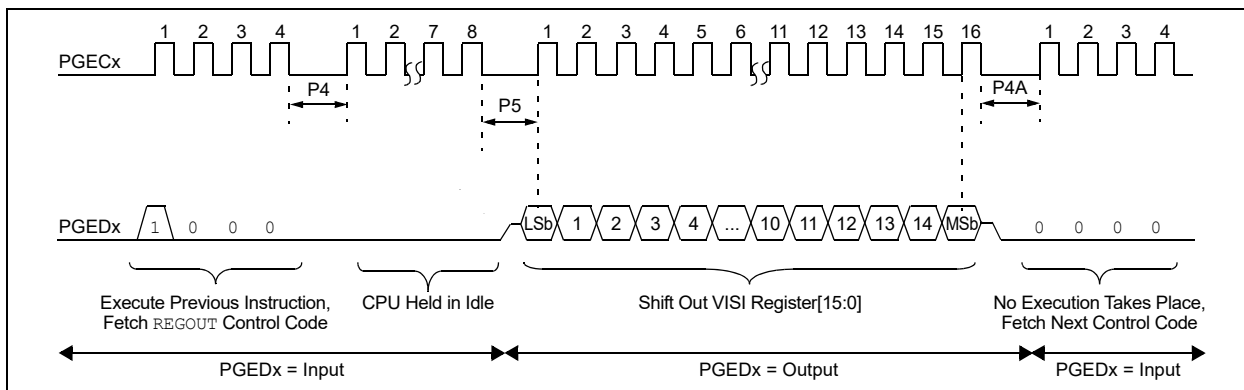
3.3.2 REGOUT SERIAL INSTRUCTION EXECUTION

The `REGOUT` control code allows for data to be extracted from the device in ICSP mode. It is used to clock the contents of the VISI register, out of the device, over the `PGEDx` pin. After the `REGOUT` control code is received, the CPU is held Idle for eight cycles. After these eight cycles, an additional 16 cycles are required to clock the data out (see [Figure 3-4](#)).

The `REGOUT` code is unique as the `PGEDx` pin is an input when the control code is transmitted to the device. However, after the control code is processed, the `PGEDx` pin becomes an output as the VISI register is shifted out.

- Note 1:** After the contents of VISI are shifted out, the dsPIC33CDVL64MC106 device maintains `PGEDx` as an output until the first rising edge of the next clock is received.
- 2:** Data change on the falling edge and latch on the rising edge of `PGEDx`. For all data transmissions, the Least Significant bit (LSb) is transmitted first.

FIGURE 3-4: REGOUT SERIAL EXECUTION



3.4 Flash Memory Programming in ICSP Mode

3.4.1 PROGRAMMING OPERATIONS

Flash memory write and erase operations are controlled by the NVMCON register. Programming is performed by setting NVMCON to select the type of erase operation (Table 3-2) or write operation (Table 3-3) and initiating the programming by setting the WR control bit (NVMCON[15]).

The PGECx clock is required to complete the programming operation. The WR control bit is cleared by hardware when the operation is finished. Refer to Section 8.0 “AC/DC Characteristics and Timing Requirements” for detailed information about the maximum time required for various programming operations.

TABLE 3-2: NVMCON ERASE OPERATIONS

| NVMCON Value | Erase Operation |
|--------------|--|
| 0x400E | Bulk Erase of user memory only (does not erase Device ID, Programming Executive memory and OTP Words). |
| 0x4003 | Page Erase of program or Programming Executive memory. |

TABLE 3-3: NVMCON WRITE OPERATIONS

| NVMCON Value | Write Operation |
|--------------|------------------------------------|
| 0x4001 | Double-Word Programming operation. |
| 0x4003 | Row Programming operation. |

3.4.2 STARTING AND STOPPING A PROGRAMMING CYCLE

For protection against accidental operations, the erase/write initiation sequence must be written to the NVMKEY register to allow any erase or program operation to proceed. The two instructions following the start of the programming sequence should be NOPs. To start an erase or write sequence, the following steps must be completed:

1. Write 0x55 to the NVMKEY register.
2. Write 0xAA to the NVMKEY register.
3. Set the WR bit in the NVMCON register.
4. Execute three NOP instructions.

The WR bit can be polled to generate enough clock cycles for the programming operation and to determine if the erase or write cycle has been completed.

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REGISTER 3-1: NVMCON: NONVOLATILE MEMORY (NVM) CONTROL REGISTER

| | | | | | | | |
|-----------------------|----------------------|----------------------|------------------------|-----|-----|---------------------|----------------------|
| R/SO-0 ⁽¹⁾ | R/W-0 ⁽¹⁾ | R/W-0 ⁽¹⁾ | R/W-0 | U-0 | U-0 | R/W-0 | R/C-0 |
| WR | WREN | WRERR | NVMSIDL ⁽²⁾ | — | — | RPDF ⁽⁶⁾ | URERR ⁽⁶⁾ |
| bit 15 | | | | | | | bit 8 |

| | | | | | | | |
|-------|-----|-----|-----|-----------------------------|----------------------|----------------------|----------------------|
| U-0 | U-0 | U-0 | U-0 | R/W-0 ⁽¹⁾ | R/W-0 ⁽¹⁾ | R/W-0 ⁽¹⁾ | R/W-0 ⁽¹⁾ |
| — | — | — | — | NVMOP[3:0] ^(3,4) | | | |
| bit 7 | | | | | | | bit 0 |

| | | |
|-------------------|-------------------|------------------------------------|
| Legend: | C = Clearable bit | SO = Settable Only bit |
| R = Readable bit | W = Writable bit | U = Unimplemented bit, read as '0' |
| -n = Value at POR | '1' = Bit is set | '0' = Bit is cleared |
| | | x = Bit is unknown |

- bit 15 **WR:** Write Control bit⁽¹⁾
 1 = Initiates a Flash memory program or erase operation; the operation is self-timed and the bit is cleared by hardware once operation is complete
 0 = Program or erase operation is complete and inactive
- bit 14 **WREN:** Write Enable bit⁽¹⁾
 1 = Enables Flash program/erase operations
 0 = Inhibits Flash program/erase operations
- bit 13 **WRERR:** Write Sequence Error Flag bit⁽¹⁾
 1 = An improper program/erase sequence attempt or termination has occurred (bit is set automatically on any set attempt of the WR bit)
 0 = The program/erase operation completed normally
- bit 12 **NVMSIDL:** NVM Stop in Idle Control bit⁽²⁾
 1 = Flash voltage regulator goes into Standby mode during Idle mode
 0 = Flash voltage regulator is active during Idle mode
- bit 11-10 **Unimplemented:** Read as '0'
- bit 9 **RPDF:** Row Programming Data Format Control bit⁽⁶⁾
 1 = Row data to be stored in RAM are in compressed format
 0 = Row data to be stored in RAM are in uncompressed format
- bit 8 **URERR:** Row Programming Data Underrun Error Flag bit⁽⁶⁾
 1 = Row Programming operation has been terminated due to data underrun error
 0 = No data underrun error has occurred
- bit 7-4 **Unimplemented:** Read as '0'

- Note 1:** These bits can only be reset on a POR.
- Note 2:** If this bit is set, there will be minimal power savings (IDLE), and upon exiting Idle mode, there is a delay (TVREG) before Flash memory becomes operational.
- Note 3:** All other combinations of NVMOP[3:0] are unimplemented.
- Note 4:** Execution of the PWRSAV instruction is ignored while any of the NVM operations are in progress.
- Note 5:** Two adjacent words on a 4-word boundary are programmed during execution of this operation.
- Note 6:** Not used in ICSP™ mode.

REGISTER 3-1: NVMCON: NONVOLATILE MEMORY (NVM) CONTROL REGISTER (CONTINUED)

bit 3-0 **NVMOP[3:0]:** NVM Operation Select bits^(1,3,4)

- 1111 = Reserved
- 1110 = User memory Bulk Erase operation
- 1010 = Reserved
- 1001 = Reserved
- 1000 = Reserved
- 0101 = Reserved
- 0100 = Inactive Partition memory erase operation
- 0011 = Memory Page Erase operation
- 0010 = Memory Row Program operation⁽⁶⁾
- 0001 = Memory Double-Word Program operation⁽⁵⁾
- 0000 = Reserved

Note 1: These bits can only be reset on a POR.

- 2:** If this bit is set, there will be minimal power savings (IDLE), and upon exiting Idle mode, there is a delay (TVREG) before Flash memory becomes operational.
- 3:** All other combinations of NVMOP[3:0] are unimplemented.
- 4:** Execution of the PWRSAV instruction is ignored while any of the NVM operations are in progress.
- 5:** Two adjacent words on a 4-word boundary are programmed during execution of this operation.
- 6:** Not used in ICSP™ mode.

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3.5 Erasing Program Memory

Figure 3-5 shows a high-level overview for the Bulk Erase of code memory.

Table 3-4 provides the ICSP programming process for erasing the program memory.

Note: Program memory must be erased before writing any data to program memory.

FIGURE 3-5: BULK ERASE FLOW

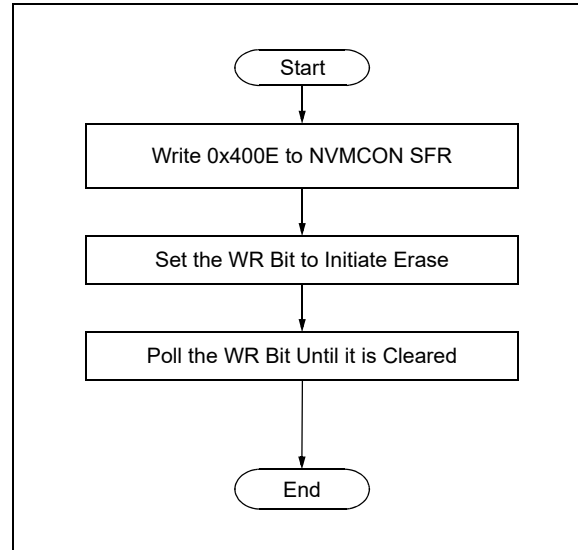


TABLE 3-4: SERIAL INSTRUCTION EXECUTION FOR BULK ERASE OF CODE MEMORY

| Command (Binary) | Data (Hex) | Description |
|--|------------|------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Set the NVMCON register to erase all user program memory. | | |
| 0000 | 2400EA | MOV #0x400E, W10 |
| 0000 | 88468A | MOV W10, NVMCON |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 3: Initiate the erase cycle. | | |
| 0000 | 200551 | MOV #0x55, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | 200AA1 | MOV #0xAA, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | A8E8D1 | BSET NVMCON, #WR |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |

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TABLE 3-4: SERIAL INSTRUCTION EXECUTION FOR BULK ERASE OF CODE MEMORY (CONTINUED)

| Command (Binary) | Data (Hex) | Description |
|--|------------|--|
| Step 4: Generate clock pulses for the code memory Bulk Erase operation to complete until the WR bit is clear. | | |
| 0000 | 000000 | NOP |
| 0000 | 804680 | MOV NVMCON, W0 |
| 0000 | 000000 | NOP |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| — | — | Repeat until the WR bit is clear. |

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3.6 Page Erase

Figure 3-6 shows a high-level overview for erasing a page of code memory.

Table 3-5 provides the ICSP programming details for erasing a page of code memory.

Note: For Page Erase operations, the NVMCON value must be modified as per Table 3-2. The NVMADR/U registers must point to any of the locations of the page to be erased.

FIGURE 3-6: PAGE ERASE FLOW

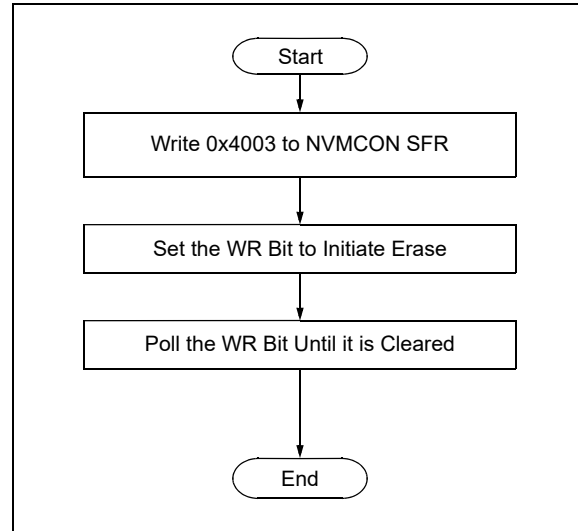


TABLE 3-5: SERIAL INSTRUCTION EXECUTION FOR ERASING A PAGE OF CODE MEMORY

| Command (Binary) | Data (Hex) | Description |
|--|------------|------------------------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Set the NVMADRU/NVMADR register pair to point to the correct page to be erased. | | |
| 0000 | 2xxxx3 | MOV #DestinationAddress<15:0>, W3 |
| 0000 | 2xxxx4 | MOV #DestinationAddress<23:16>, W4 |
| 0000 | 884693 | MOV W3, NVMADR |
| 0000 | 8846A4 | MOV W4, NVMADRU |
| Step 3: Set the NVMCON register to erase the first page of executive memory. | | |
| 0000 | 24003A | MOV #0x4003, W10 |
| 0000 | 88468A | MOV W10, NVMCON |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |

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TABLE 3-5: SERIAL INSTRUCTION EXECUTION FOR ERASING A PAGE OF CODE MEMORY (CONTINUED)

| Command (Binary) | Data (Hex) | Description |
|--|---------------|--|
| Step 4: Initiate the erase cycle. | | |
| 0000 | 200551 | MOV #0x55, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | 200AA1 | MOV #0xAA, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | A8E8D1 | BSET NVMCON, #WR |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 5: Generate clock pulses for the Page Erase operation to complete until the WR bit is clear. | | |
| 0000 | 000000 | NOP |
| 0000 | 804680 | MOV NVMCON, W0 |
| 0000 | 000000 | NOP |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| — | — | Repeat until the WR bit is clear. |

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3.7 Writing Code Memory

Figure 3-8 shows a high-level overview for writing the code memory.

Table 3-6 provides the ICSP programming details for writing the code memory.

Code memory is written two instruction words at a time. Two words are loaded into the write latches, located at 0xFA0000 and 0xFA0002, using the packed data format shown in Figure 3-7. The destination address is loaded into the NVMADR and NVMADRU registers. Next, the write cycle is initiated by setting the WREN bit in the NVMCON register. The WR bit in NVMCON will be cleared in hardware once the double-word write is complete. This process is repeated for all memory locations to be programmed.

FIGURE 3-7: PACKED INSTRUCTION WORD FORMAT

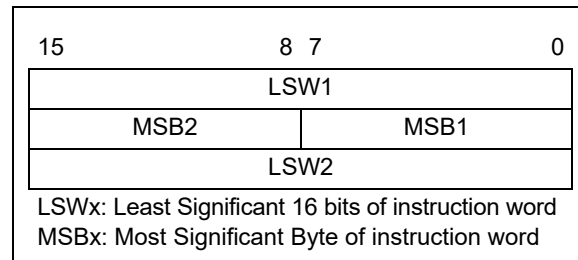
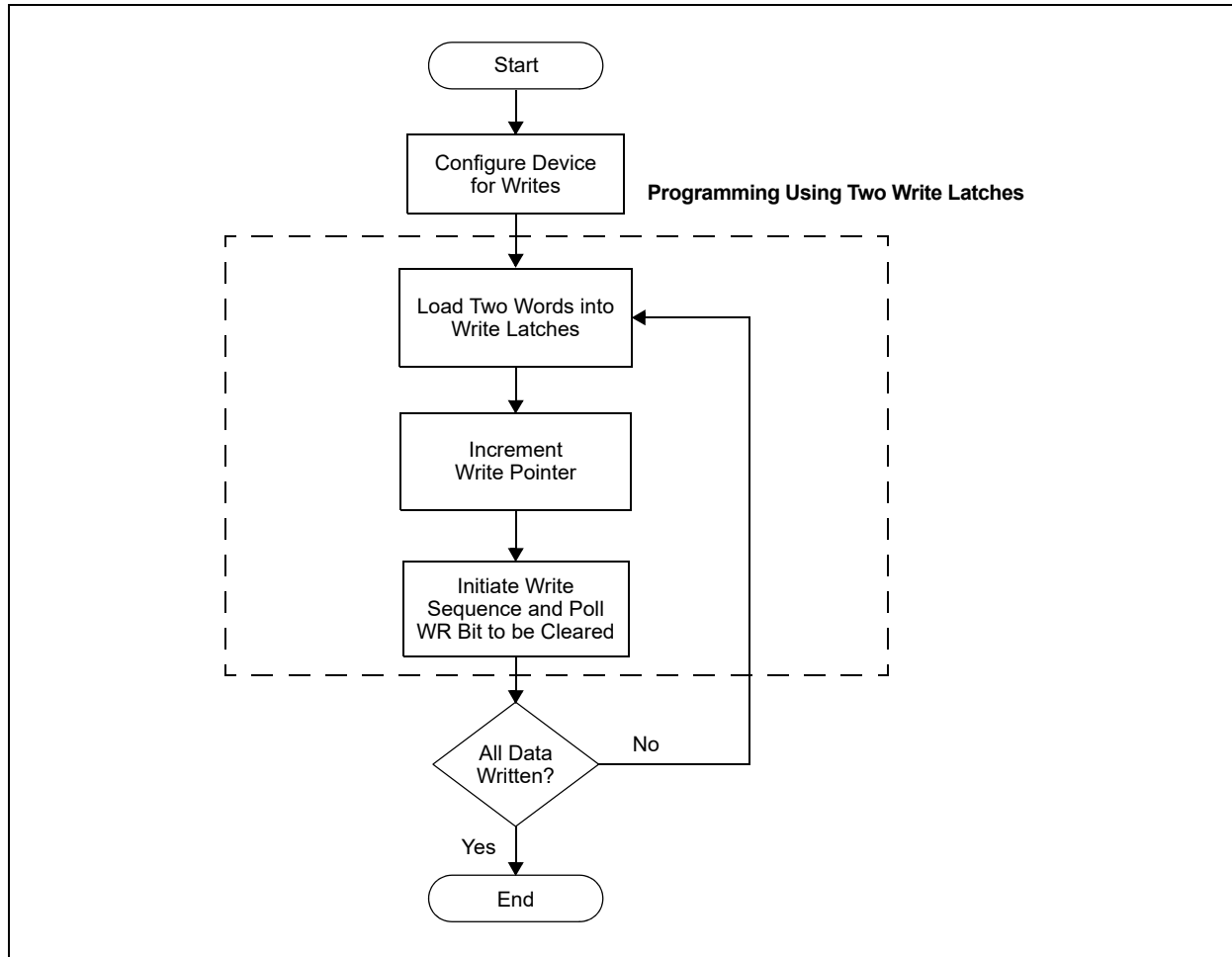


FIGURE 3-8: PROGRAM CODE MEMORY FLOW



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**TABLE 3-6: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CODE MEMORY:
TWO-WORD LATCH WRITES**

| Command (Binary) | Data (Hex) | Description |
|--|------------|------------------------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register for writing to the latches. | | |
| 0000 | 200FAC | MOV #0xFA, W12 |
| 0000 | 8802AC | MOV W12, TBLPAG |
| Step 3: Load W0:W2 with the next two packed instruction words to program. | | |
| 0000 | 2xxxx0 | MOV #<LSW0>, W0 |
| 0000 | 2xxxx1 | MOV #<MSB1:MSB0>, W1 |
| 0000 | 2xxxx2 | MOV #<LSW1>, W2 |
| Step 4: Set the Read Pointer (W6) and Write Pointer (W7), and load the (next set of) write latches. | | |
| 0000 | EB0300 | CLR W6 |
| 0000 | 000000 | NOP |
| 0000 | EB0380 | CLR W7 |
| 0000 | 000000 | NOP |
| 0000 | BB0BB6 | TBLWTL [W6++], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BBDBB6 | TBLWTH.B [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BBEBB6 | TBLWTH.B [W6++], [++W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BB0B96 | TBLWTL.W [W6], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 5: Set the NVMADRU/NVMADR register pair to point to the correct address. | | |
| 0000 | 2xxxx3 | MOV #DestinationAddress<15:0>, W3 |
| 0000 | 2xxxx4 | MOV #DestinationAddress<23:16>, W4 |
| 0000 | 884693 | MOV W3, NVMADR |
| 0000 | 8846A4 | MOV W4, NVMADRU |
| Step 6: Set the NVMCON register to program two instruction words. | | |
| 0000 | 24001A | MOV #0x4001, W10 |
| 0000 | 000000 | NOP |
| 0000 | 88468A | MOV W10, NVMCON |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |

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**TABLE 3-6: SERIAL INSTRUCTION EXECUTION FOR PROGRAMMING CODE MEMORY:
TWO-WORD LATCH WRITES (CONTINUED)**

| Command (Binary) | Data (Hex) | Description |
|---|---------------|---|
| Step 7: Initiate the write cycle. | | |
| 0000 | 200551 | MOV #0x55, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | 200AA1 | MOV #0xAA, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | A8E8D1 | BSET NVMCON, #WR |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 8: Generate clock pulses for the program operation to complete until the WR bit is clear. | | |
| 0000 | 000000 | NOP |
| 0000 | 804680 | MOV NVMCON, W0 |
| 0000 | 000000 | NOP |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| — | — | Repeat until the WR bit is clear. |
| Step 9: Repeat Steps 3-8 until all code memory is programmed. | | |

3.8 Writing Configuration Bits

The procedure for writing Configuration bits is similar to the procedure for writing code memory.

To change the values of the Configuration bits once they have been programmed, the device must be erased, as described in [Section 3.5 “Erasing Program Memory”](#), and reprogrammed to the desired value.

[Table 3-7](#) provides the ICSP programming details for writing the Configuration bits.

The code protection can be enabled by programming ‘0’ in the Code Protection Configuration bits. In order to verify the data by reading the Configuration bits after performing the write, the code protection bits should initially be programmed to ‘1’ to ensure that the verification can be performed properly. After verification is finished, the code protection bits can be programmed to ‘0’ by using a word write to the appropriate Configuration register.

TABLE 3-7: SERIAL INSTRUCTION EXECUTION FOR WRITING CONFIGURATION WORDS

| Command (Binary) | Data (Hex) | Description |
|---|------------|------------------------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register for writing to the latches. | | |
| 0000 | 200FAC | MOV #0xFA, W12 |
| 0000 | 8802AC | MOV W12, TBLPAG |
| Step 3: Load W0:W1 with the next two Configuration Words to program. | | |
| 0000 | 2xxxx0 | MOV #<Config1 lower word data>, W0 |
| 0000 | 2xxxx1 | MOV #<Config1 upper word data>, W1 |
| 0000 | 2xxxx2 | MOV #<Config2 lower word data>, W2 |
| 0000 | 2xxxx3 | MOV #<Config2 upper word data>, W3 |
| Step 4: Set the Write Pointer (W3) and load the write latches. | | |
| 0000 | EB0300 | CLR W6 |
| 0000 | 000000 | NOP |
| 0000 | BB0B00 | TBLWTL W0, [W6] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BB9B01 | TBLWTH W1, [W6++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BB0B02 | TBLWTL W2, [W6] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BB9B03 | TBLWTH W3, [W6++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 5: Set the NVMADRU/NVMADR register pair to point to the correct Configuration Word address. | | |
| 0000 | 2xxxx4 | MOV #DestinationAddress<15:0>, W4 |
| 0000 | 2xxxx5 | MOV #DestinationAddress<23:16>, W5 |
| 0000 | 884694 | MOV W4, NVMADR |
| 0000 | 8846A5 | MOV W5, NVMADRU |

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TABLE 3-7: SERIAL INSTRUCTION EXECUTION FOR WRITING CONFIGURATION WORDS (CONTINUED)

| Command (Binary) | Data (Hex) | Description |
|---|------------|--|
| Step 6: Set the NVMCON register to program two instruction words. | | |
| 0000 | 24001A | MOV #0x4001, W10 |
| 0000 | 000000 | NOP |
| 0000 | 88468A | MOV W10, NVMCON |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 7: Initiate the write cycle. | | |
| 0000 | 200551 | MOV #0x55, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | 200AA1 | MOV #0xAA, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | A8E8D1 | BSET NVMCON, #WR |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 8: Generate clock pulses for the program operation to complete until the WR bit is clear. | | |
| 0000 | 000000 | NOP |
| 0000 | 804680 | MOV NVMCON, W0 |
| 0000 | 000000 | NOP |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| — | — | Repeat until the WR bit is clear. |
| Step 9: Repeat Steps 3-8 until all Configuration registers are programmed. | | |

3.9 Writing OTP Words

The procedure for writing to user OTP memory is similar to the procedure for writing to user program memory, except that each of the 64 OTP double-word pairs can only be written once. Both words in each OTP location must be written together using the same two-word latch write process used to write code memory.

Writing anything, with the exception of all '1's, to an OTP location generates an ECC checksum and renders that location used. Attempting to write to an OTP location that has already been programmed will cause an ECC checksum error the next time that location is read. Care should be taken to avoid writing to OTP locations that have already been programmed or may need to be programmed at a later time. See [Figure 2-5](#) for the location of user OTP memory.

[Figure 3-9](#) shows a high-level overview of the OTP programming process.

3.10 Reading OTP Words

The procedure for reading OTP Words is similar to the procedure for reading code memory. Since there are multiple OTP Words, they are read one at a time.

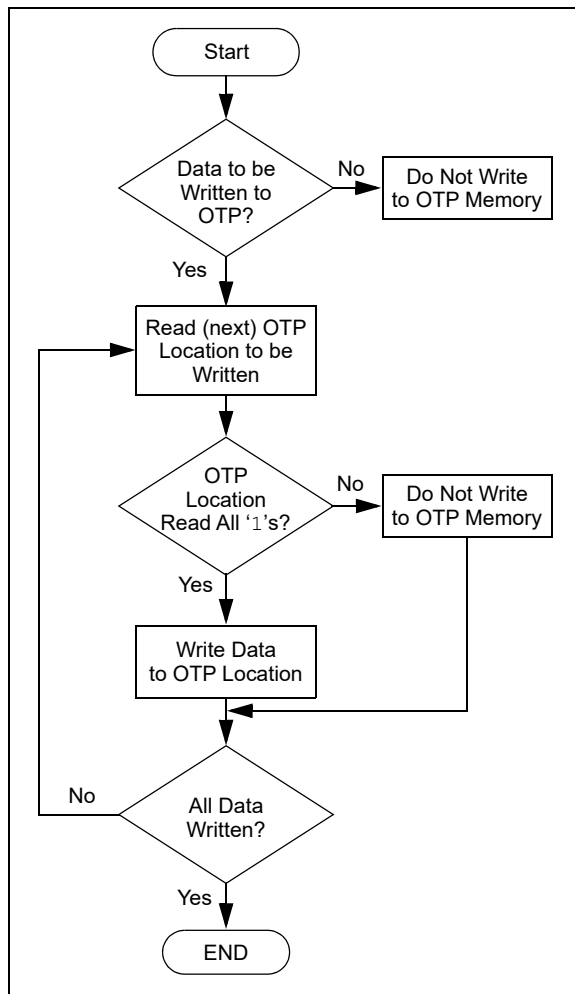
3.11 Reading Code Memory

Reading from code memory is performed by executing a series of `TBLRD` instructions and clocking out the data using the `REGOUT` command.

[Table 3-8](#) provides the ICSP programming details for reading code memory.

To minimize reading time, the same packed data format that the write procedure uses is utilized. See [Section 3.7 "Writing Code Memory"](#) for more details on the packed data format.

FIGURE 3-9: OTP PROGRAMMING PROCESS



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TABLE 3-8: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY

| Command (Binary) | Data (Hex) | Description |
|--|------------|-------------------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register and the Read Pointer (W6) for the TBLRD instruction. | | |
| 0000 | 200xx0 | MOV #<SourceAddress23:16>, W0 |
| 0000 | 8802A0 | MOV W0, TBLPAG |
| 0000 | 2xxxx6 | MOV #<SourceAddress15:0>, W6 |

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TABLE 3-8: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY (CONTINUED)

| Command (Binary) | Data (Hex) | Description |
|---|---------------|-------------------------|
| Step 3: Initialize the Write Pointer (W7) and store the next four locations of code memory to W0:W5. | | |
| 0000 | EB0380 | CLR W7 |
| 0000 | 000000 | NOP |
| 0000 | BA1B96 | TBLRDL [W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBB6 | TBLRDH.B [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBD6 | TBLRDH.B [++W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BA1BB6 | TBLRDL [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BA1B96 | TBLRDL [W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBB6 | TBLRDH.B [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBD6 | TBLRDH.B [++W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BA0BB6 | TBLRDL [W6++], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |

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TABLE 3-8: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY (CONTINUED)

| Command (Binary) | Data (Hex) | Description |
|---|------------|--|
| Step 4: Output W0:W5 using the VISI register and REGOUT command. | | |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E61 | MOV W1, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E62 | MOV W2, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E63 | MOV W3, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E64 | MOV W4, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E65 | MOV W5, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| Step 5: Reset the device's internal PC. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 6: Repeat Steps 3-5 until all desired code memory is read. | | |

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3.12 Reading Configuration Registers

Table 3-9 provides the ICSP programming details for reading the Configuration Words.

The procedure for reading the Configuration Words is similar to the procedure for reading code memory. Since there are multiple Configuration Words, they are read one at a time.

TABLE 3-9: SERIAL INSTRUCTION EXECUTION FOR READING CONFIGURATION WORDS

| Command (Binary) | Data (Hex) | Description |
|--|------------|--|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register, the Write Pointer (W7) and the Read Pointer (W6) for the TBLRD instruction. | | |
| 0000 | 200xx0 | MOV #<Address23:16>, W0 |
| 0000 | 20FCC7 | MOV #VISI, W7 |
| 0000 | 8802A0 | MOV W0, TBLPAG |
| 0000 | 2xxxx6 | MOV #<Address15:0>, W6 |
| Step 3: Store the Configuration register and send the contents of the VISI register. | | |
| 0000 | 000000 | NOP |
| 0000 | BA8B96 | TBLRDH [W6], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | BA0B96 | TBLRDL [W6], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| Step 4: Repeat Steps 1-3 until all Configuration registers are read. | | |

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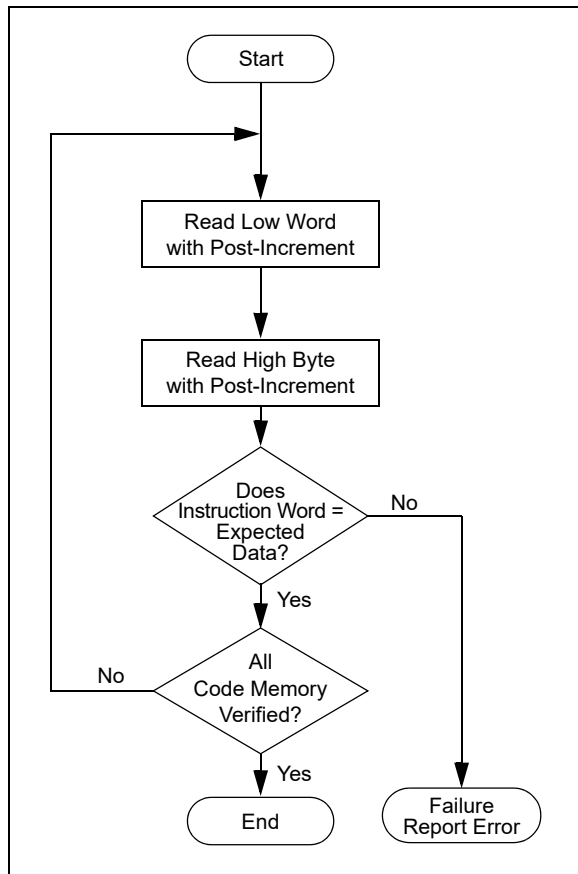
3.13 Verify Code Memory and Configuration Bits

The verify step involves reading back the code memory space and comparing it against the copy held in the programmer's buffer. The Configuration Words are verified with the rest of the code.

The verify process is shown in Figure 3-10. The lower word of the instruction is read, and then the lower byte of the upper word is read and compared against the instruction stored in the programmer's buffer. Refer to Section 3.11 "Reading Code Memory" for implementation details of reading code memory.

Note: Because the Configuration Words include the device code protection bit, code memory should be verified immediately after writing if code protection is to be enabled. This is because the device will not be readable or verifiable if a device Reset occurs after the code-protect bit has been cleared.

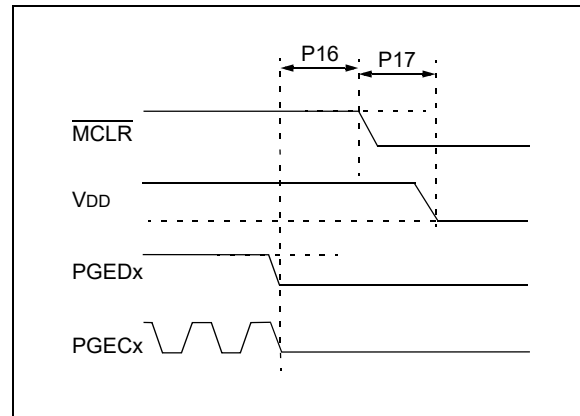
FIGURE 3-10: VERIFY CODE MEMORY FLOW



3.14 Exiting ICSP Mode

Exiting Program/Verify mode is done by removing VDD from MCLR, as shown in Figure 3-11. The only requirement for exit is that an interval, P16, should elapse between the last clock, and the program signals on PGECx and PGEDx before removing VDD.

FIGURE 3-11: EXITING ICSP™ MODE



4.0 DEVICE PROGRAMMING – ENHANCED ICSP

This section discusses programming the device through Enhanced ICSP and the Programming Executive. The Programming Executive resides in executive memory (separate from code memory) and is executed when Enhanced ICSP Programming mode is entered. The Programming Executive provides the mechanism for the programmer (host device) to program and verify the dsPIC33CDVL64MC106 device using a simple command set and communication protocol. There are several basic functions provided by the Programming Executive:

- Read Memory
- Erase Memory
- Program Memory
- Blank Check

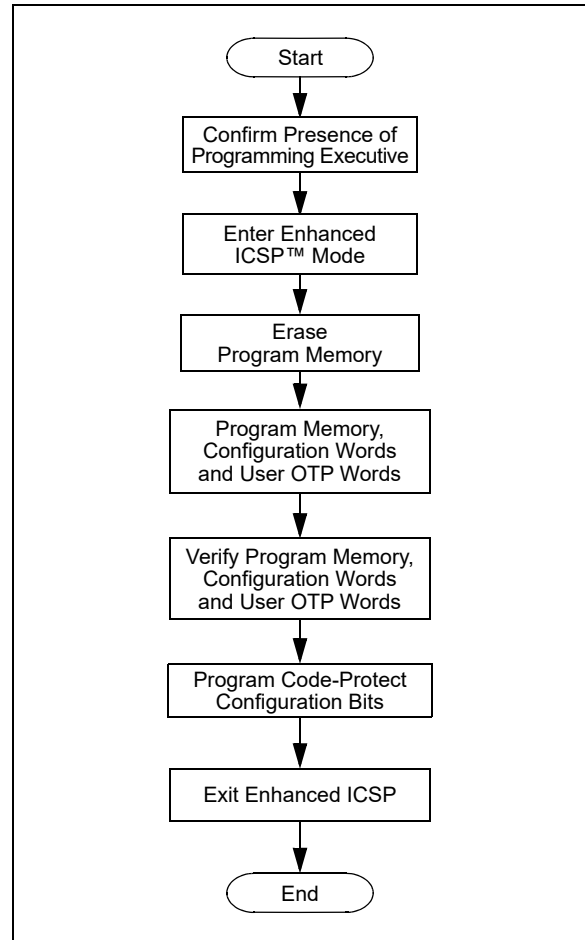
The Programming Executive performs the low-level tasks required for erasing, programming and verifying a device. This allows the programmer to program the device by issuing the appropriate commands and data. A detailed description for each command is provided in [Section 5.2 “Programming Executive Commands”](#).

Note: The PE uses the device’s data RAM for variable storage and program execution. After running the PE, no assumptions should be made about the contents of data RAM.

4.1 Overview of the Programming Process

Figure 4-1 shows the high-level overview of the programming process. First, it must be determined if the Programming Executive is present in executive memory, then the Enhanced ICSP mode is entered. The program memory is then erased, and the program memory and Configuration Words are programmed and verified. Last, the code-protect Configuration bits are programmed (if required) and Enhanced ICSP mode is exited.

FIGURE 4-1: HIGH-LEVEL ENHANCED ICSP™ PROGRAMMING FLOW



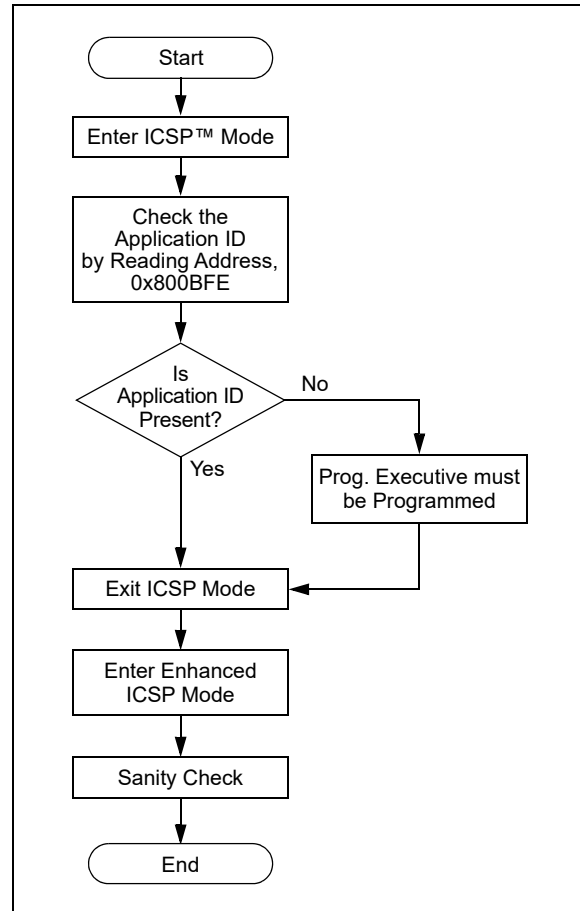
4.2 Confirming the Presence of the Programming Executive

Before programming can begin, the programmer must confirm that the Programming Executive is stored in executive memory. The procedure for this task is shown in Figure 4-2.

First, In-Circuit Serial Programming (ICSP) mode is entered. Then, the unique Application ID Word stored in executive memory is read. If the Programming Executive is resident, the correct Application ID Word, 0xDF, is read and programming can resume as normal. However, if the Application ID Word is not present, the PE must be programmed to executive code memory using the method described in Section 5.0 “The Programming Executive”.

Section 3.0 “Device Programming – ICSP” describes the ICSP programming method. Section 4.3 “Reading the Application ID Word” describes the procedure for reading the Application ID Word in ICSP mode.

FIGURE 4-2: CONFIRMING PRESENCE OF PROGRAMMING EXECUTIVE



4.3 Reading the Application ID Word

The Application ID Word is stored at the address, 0x800BFE, in executive code memory. To read this memory location, you must use the `SIX` control code to move this program memory location to the `VISI` register. Then, the `REGOUT` control code must be used to clock the contents of the `VISI` register out of the device. The corresponding control and instruction codes that must be serially transmitted to the device to perform this operation are provided in [Table 4-1](#).

After the programmer has clocked out the Application ID Word, it must be inspected. If the Application ID has the value, 0xDF, the Programming Executive is resident in memory and the device can be programmed using the mechanism described in [Section 4.0 “Device Programming – Enhanced ICSP”](#). However, if the Application ID has any other value, the Programming Executive is not resident in memory; it must be loaded into memory before the device can be programmed. The procedure for loading the Programming Executive to memory is described in [Section 5.0 “The Programming Executive”](#).

TABLE 4-1: SERIAL INSTRUCTION EXECUTION FOR READING THE APPLICATION ID WORD

| Command (Binary) | Data (Hex) | Description |
|--|------------|--|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register and the Read Pointer (W0) for the TBLRD instruction. | | |
| 0000 | 200800 | MOV #0x80, W0 |
| 0000 | 8802A0 | MOV W0, TBLPAG |
| 0000 | 20BFE0 | MOV #0xBFE, W0 |
| 0000 | 20FCC1 | MOV #VISI, W1 |
| 0000 | 000000 | NOP |
| 0000 | BA0890 | TBLRDL [W0], [W1] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 3: Output the VISI register using the REGOUT command. | | |
| 0001 | <VISI> | Clock out contents of the VISI register. |

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4.4 Entering Enhanced ICSP Mode

As shown in Figure 4-3, entering Enhanced ICSP Program/Verify mode requires three steps:

1. The MCLR pin is briefly driven high, then low.
2. A 32-bit key sequence is clocked into PGEDx. An interval of at least P18 must elapse before presenting the key sequence on PGEDx.
3. MCLR is held within a specified period of time, then driven high.

The key sequence is a specific 32-bit pattern, '0100 1101 0100 0011 0100 1000 0101 0000' (more easily remembered as 0x4D434850 in hexadecimal format). The device will enter Program/Verify mode only if the key sequence is valid. The Most Significant bit (MSb) of the most significant nibble must be shifted in first.

Once the key sequence is complete, VDD must be applied to MCLR and held at that level for as long as Program/Verify mode is to be maintained. An interval time of at least time, P19, P7 and P1 * 5, must elapse before presenting data on PGEDx. Signals appearing on PGEDx before P7 has elapsed will not be interpreted as valid.

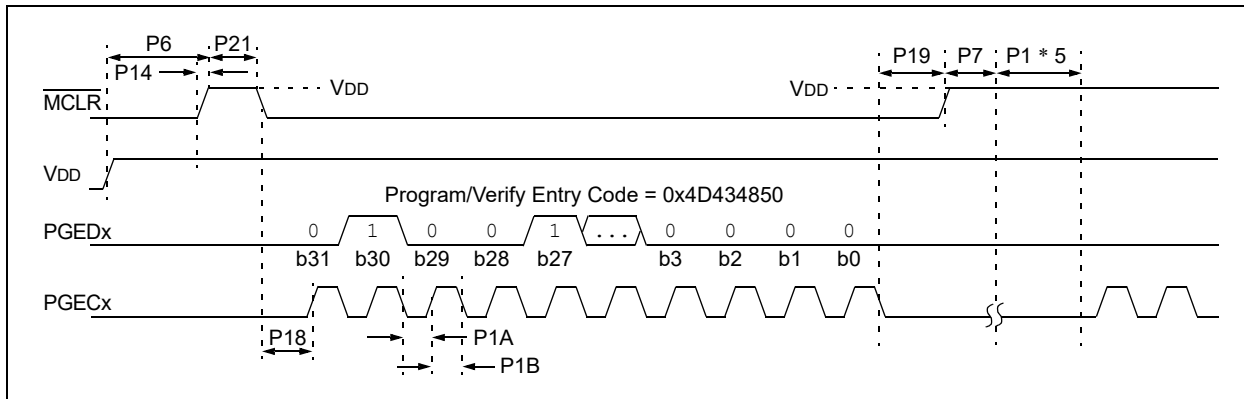
4.5 Blank Check

The term, "Blank Check", implies verifying that the device has been successfully erased and has no programmed memory locations. A blank or erased memory location is always read as '1'.

The Device ID registers (0xFF0000:0xFF0002) can be ignored by the Blank Check since this region stores device information that cannot be erased. Additionally, all unimplemented memory space and Calibration registers should be ignored by the Blank Check.

The QBLANK command is used for the Blank Check. It determines if the code memory is erased by testing these memory regions. A 'BLANK' or 'NOT BLANK' response is returned. If it is determined that the device is not blank, it must be erased before attempting to program the chip.

FIGURE 4-3: ENTERING ENHANCED ICSP™ MODE



4.6 Code Memory Programming

4.6.1 PROGRAMMING METHODOLOGY

There are two commands that can be used for programming code memory when utilizing the Programming Executive. The `PROG2W` command programs and verifies two 24-bit instruction words into the program memory, starting at the specified address. The second and faster command, `PROGP`, programs and verifies an entire row of 128 24-bit instruction words to program memory, starting at the specified address. Please ensure that the starting address is on a row boundary when using Row Programming. See [Section 5.0 “The Programming Executive”](#) for a full description of each of these commands.

[Figure 4-4](#) and [Figure 4-5](#) show a high-level overview of the code memory programming process using the `PROG2W` and `PROGP` commands.

FIGURE 4-4: FLOWCHART FOR DOUBLE-WORD PROGRAMMING

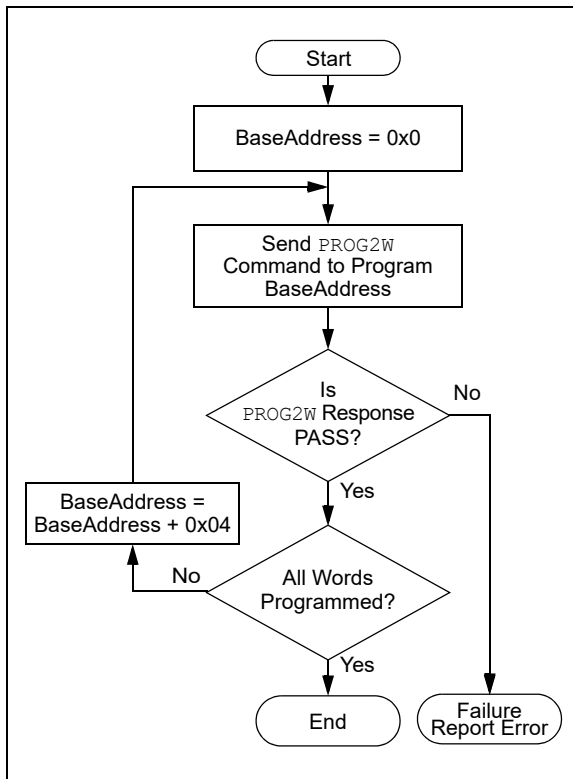
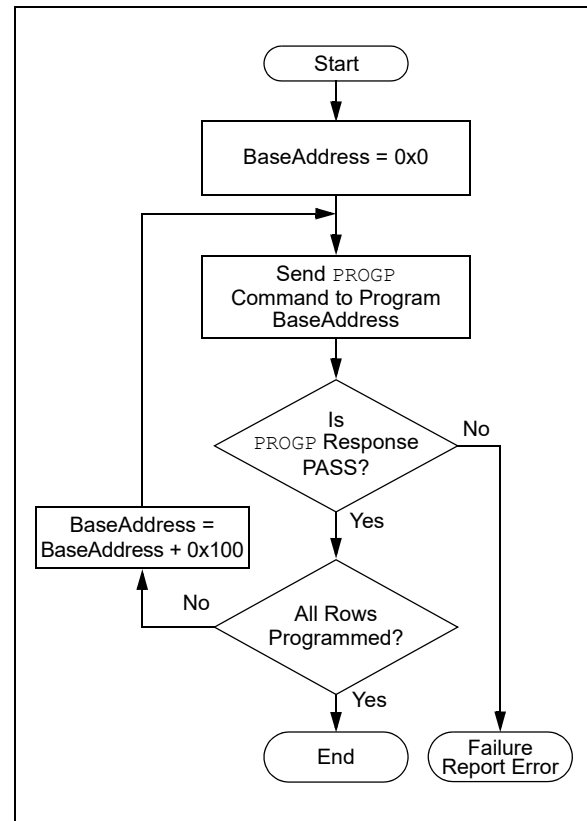


FIGURE 4-5: FLOWCHART FOR ROW PROGRAMMING

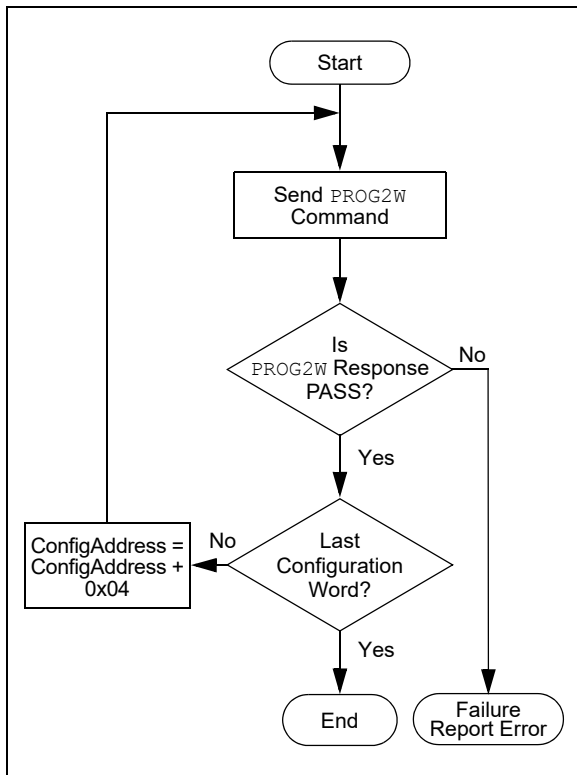


4.7 Configuration Bits Programming

The Configuration bits are programmed, one 16-bit word at a time, using the `PROG2W` command. This command specifies the configuration data and address. When Configuration bits are programmed, any unimplemented bits must be programmed with a '1'.

Multiple `PROG2W` commands are required to program all Configuration bits. A flowchart for Configuration bit programming is shown in [Figure 4-6](#).

FIGURE 4-6: CONFIGURATION BIT PROGRAMMING FLOW



4.8 Programming Verification

After the code memory is programmed, the contents of the memory should be verified to ensure that the programming was successful. Verification requires code memory to be read back and compared against the copy held in the programmer's buffer. The `READP` command can be used to read back all the programmed code memory and Configuration Words.

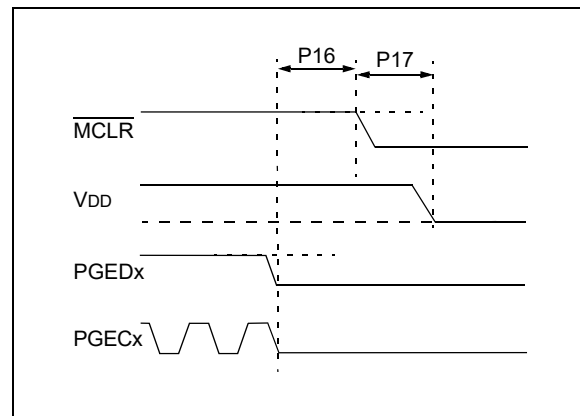
Alternatively, the programmer can perform the verification after the entire device is programmed using a checksum computation.

See [Section 7.0 "CRC Checksum Computation"](#) for more information on calculating the checksum.

4.9 Exiting Enhanced ICSP Mode

Exiting Program/Verify mode is done by removing `VDD` from `MCLR`, as shown in [Figure 4-7](#). The only requirement for exit is that an interval, `P16`, should elapse between the last clock, and program signals on `PGECx` and `PGEDx`, before removing `VDD`.

FIGURE 4-7: EXITING ENHANCED ICSP™ MODE



5.0 THE PROGRAMMING EXECUTIVE

Note: The Programming Executive can be obtained from each device page on the Microchip website: www.microchip.com.

5.1 Programming Executive Communication

The programmer and Programming Executive have a master-slave relationship, where the programmer is the master programming device and the Programming Executive is the slave.

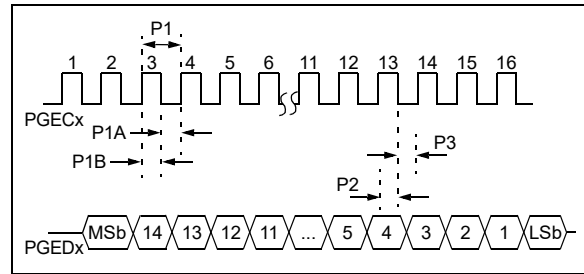
All communication is initiated by the programmer in the form of a command. Only one command at a time can be sent to the Programming Executive. In turn, the PE only sends one response to the programmer after receiving and processing a command. The PE command set is described in [Section 5.2 “Programming Executive Commands”](#). The response set is described in [Section 5.3 “Programming Executive Responses”](#).

5.1.1 COMMUNICATION INTERFACE AND PROTOCOL

The ICSP/Enhanced ICSP interface is a two-wire SPI, implemented using the PGECx and PGEDx pins. The PGECx pin is used as a clock input pin and the clock source must be provided by the programmer. The PGEDx pin is used for sending command data to, and receiving response data from, the PE.

Note: For Enhanced ICSP, all serial data are transmitted on the falling edge of PGECx and latched on the rising edge of PGECx. All data transmissions are sent to the MSb first using 16-bit mode (see [Figure 5-1](#)).

FIGURE 5-1: PROGRAMMING EXECUTIVE SERIAL TIMING



Since a two-wire SPI is used, and data transmissions are bidirectional, a simple protocol is used to control the direction of PGEDx. When the programmer completes a command transmission, it releases the PGEDx line and allows the PE to drive this line high. The PE keeps the PGEDx line high to indicate that it is processing the command.

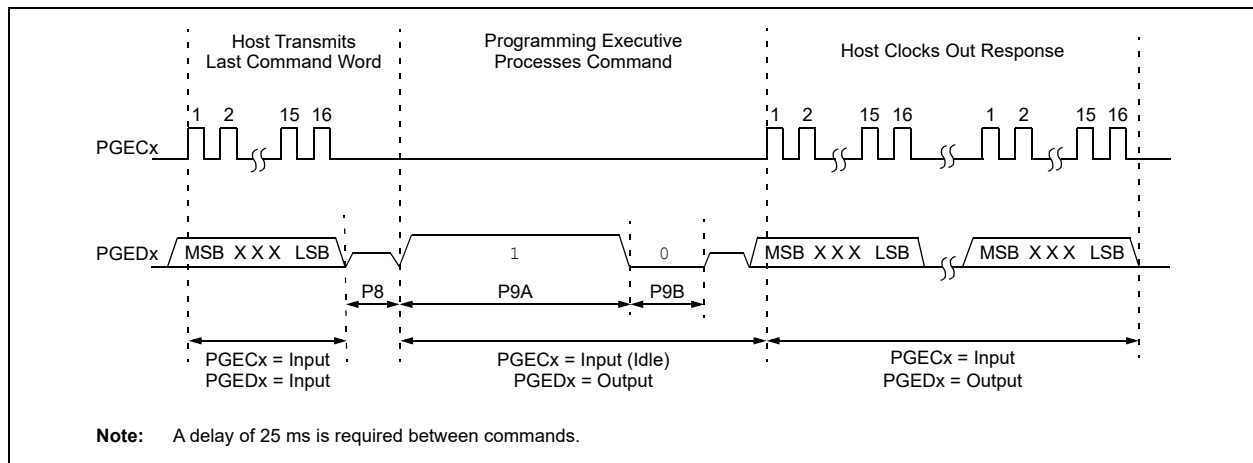
After the PE has processed the command, it brings PGEDx low (P9B) to indicate to the programmer that the response is available to be clocked out. The programmer can begin to clock out the response after a maximum wait (P9B) and the programmer must provide the necessary amount of clock pulses to receive the entire response from the PE.

After the entire response is clocked out, the programmer should terminate the clock on PGECx until it is time to send another command to the Programming Executive. This protocol is shown in [Figure 5-2](#).

5.1.2 SPI RATE

In Enhanced ICSP mode, the dsPIC33CDVL64MC106 device operates from the internal Fast RC (FRC) oscillator, which has a nominal frequency of 8 MHz. This oscillator frequency yields an effective system clock frequency of 4 MHz. To ensure that the programmer does not clock too fast, it is recommended that a 2 MHz clock be provided by the programmer.

FIGURE 5-2: PROGRAMMING EXECUTIVE – PROGRAMMER COMMUNICATION PROTOCOL



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5.1.3 TIME-OUTS

The Programming Executive uses no Watchdog Timer or time-out for transmitting responses to the programmer. If the programmer does not follow the flow control mechanism using PGECx, as described in [Section 5.1.1 “Communication Interface and Protocol”](#), it is possible that the Programming Executive will behave unexpectedly while trying to send a response to the

programmer. Since the PE has no time-out, it is imperative that the programmer correctly follows the described communication protocol.

As a safety measure, the programmer should use the command time-outs identified in [Table 5-1](#). If the command time-out expires, the programmer should reset the PE and start programming the device again.

TABLE 5-1: PROGRAMMING EXECUTIVE COMMAND SET

| Opcode | Mnemonic | Length (16-bit words) | Time-out | Description |
|--------|----------|--------------------------|----------|--|
| 0x0 | SCHECK | 1 | 1 ms | Sanity check. |
| 0x1 | Reserved | N/A | N/A | — |
| 0x2 | READP | 4 | 1 ms/row | Read 'N' 24-bit instruction words of the user Flash memory, Configuration Word or Device ID register, starting from the specified address. |
| 0x3 | PROG2W | 6 | 5 ms | Program a double instruction word of code memory at the specified address and verify. |
| 0x4 | Reserved | N/A | N/A | This command is reserved; it will return a NACK. |
| 0x5 | PROGP | 195 | 5 ms | Program 128 words of program memory at the specified starting address, then verify. |
| 0x6 | Reserved | N/A | N/A | This command is reserved; it will return a NACK. |
| 0x7 | ERASEB | 1 | 125 ms | Bulk Erase user memory. |
| 0x8 | Reserved | N/A | N/A | This command is reserved; it will return a NACK. |
| 0x9 | ERASEP | 3 | 25 ms | Command to erase a page. |
| 0xA | Reserved | N/A | N/A | This command is reserved; it will return a NACK. |
| 0xB | QVER | 1 | 1 ms | Query the Programming Executive software version. |
| 0xC | CRCP | 5 | 1s | Perform a CRC-16 on the specified range of memory. |
| 0xD | Reserved | N/A | N/A | This command is reserved; it will return a NACK. |
| 0xE | QBLANK | 5 | 700 ms | Query to check whether the code memory is blank. |

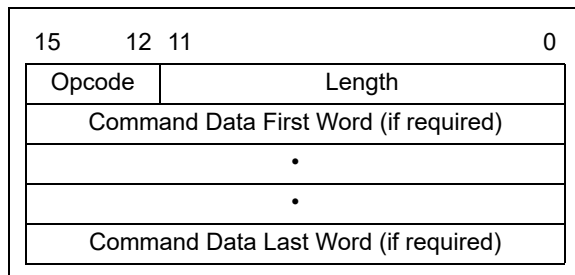
5.2 Programming Executive Commands

The Programming Executive command set is shown in [Table 5-1](#). This table contains the opcode, mnemonic, length, time-out and description for each command. Functional details on each command are provided in the command descriptions (see [Section 5.2.4 “Command Descriptions”](#)).

5.2.1 COMMAND FORMAT

All Programming Executive commands have a general format, consisting of a 16-bit header and any required data for the command (see [Figure 5-3](#)). The 16-bit header consists of a 4-bit opcode field, which is used to identify the command, followed by a 12-bit command length field.

FIGURE 5-3: COMMAND FORMAT



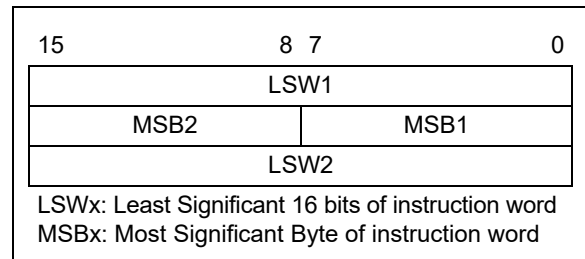
The command opcode must match one of those in the command set. Any command that is received which does not match the list in [Table 5-1](#) will return a “NACK” response (see [Section 5.3.1.1 “Opcode Field”](#)).

The command length is represented in 16-bit words since the SPI operates in 16-bit mode. The Programming Executive uses the command length field to determine the number of words to read from the SPI port. If the value of this field is incorrect, the command will not be properly received by the Programming Executive.

5.2.2 PACKED DATA FORMAT

When 24-bit instruction words are transferred across the 16-bit SPI interface, they are packed to conserve space using the format shown in [Figure 5-4](#). This format minimizes traffic over the SPI and provides the Programming Executive with data that are properly aligned for performing Table Write operations.

FIGURE 5-4: PACKED INSTRUCTION WORD FORMAT



5.2.3 PROGRAMMING EXECUTIVE ERROR HANDLING

The Programming Executive will “NACK” all unsupported commands. Additionally, due to the memory constraints of the Programming Executive, no checking is performed on the data contained in the programmer command. It is the responsibility of the programmer to command the Programming Executive with valid command arguments or the programming operation may fail. Additional information on error handling is provided in [Section 5.3.1.3 “QE_Code Field”](#).

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5.2.4 COMMAND DESCRIPTIONS

All commands supported by the Programming Executive are described in [Section 5.2.4.1 “SCHECK Command”](#) through [Section 5.2.4.9 “QBLANK Command”](#).

5.2.4.1 SCHECK Command

| | | |
|--------|--------|---|
| 15 | 12 11 | 0 |
| Opcode | Length | |

[Table 5-2](#) provides the description for the SCHECK command.

TABLE 5-2: COMMAND DESCRIPTION

| Field | Description |
|--------|-------------|
| Opcode | 0x0 |
| Length | 0x1 |

The SCHECK command instructs the Programming Executive to do nothing but generate a response. This command is used as a “Sanity Check” to verify that the Programming Executive is operational.

Expected Response (2 words):

0x1000
0x0002

Note: This instruction is not required for programming but is provided for development purposes only.

5.2.4.2 READP Command

| | | | |
|----------|--------|----------|---|
| 15 | 12 11 | 8 7 | 0 |
| Opcode | Length | | |
| N | | | |
| Reserved | | Addr_MSB | |
| Addr_LS | | | |

[Table 5-3](#) provides the description for the READP command.

TABLE 5-3: COMMAND DESCRIPTION

| Field | Description |
|----------|--|
| Opcode | 0x2 |
| Length | 0x4 |
| N | Number of 24-bit instructions to read (maximum of 32768) |
| Reserved | 0x0 |
| Addr_MSB | MSB of 24-bit source address |
| Addr_LS | Least Significant 16 bits of 24-bit source address |

The READP command instructs the Programming Executive to read N 24-bit words of code memory, Flash Configuration Words or Device ID registers, starting from the 24-bit address specified by Addr_MSB and Addr_LS. This command can only be used to read 24-bit data. All data returned in response to this command use the packed data format described in [Section 5.2.2 “Packed Data Format”](#).

Expected Response (2 + 3 * N/2 words for N even):

0x1200
2 + 3 * N/2
Least Significant Program Memory Word 1
...
Least Significant Data Word N

Expected Response (4 + 3 * (N – 1)/2 words for N odd):

0x1200
4 + 3 * (N – 1)/2
Least Significant Program Memory Word 1
...
MSB of Program Memory Word N (zero-padded)

Note 1: Reading unimplemented memory will cause the Programming Executive to reset. Please ensure that only memory locations present on a particular device are accessed.

2: When the READP command is used to read Device ID registers, the upper byte (bits[23:16]) of each word returned by the Programming Executive should be ignored.

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5.2.4.3 PROG2W Command

| | | | | | |
|-----------|----|-----------|---|---|---|
| 15 | 12 | 11 | 8 | 7 | 0 |
| Opcode | | Length | | | |
| Reserved | | Addr_MSB | | | |
| Addr_LS | | | | | |
| DataL_LS | | | | | |
| DataH_MSB | | DataL_MSB | | | |
| DataH_LS | | | | | |

Table 5-4 provides the description for the PROG2W command.

TABLE 5-4: COMMAND DESCRIPTION

| Field | Description |
|-----------|--|
| Opcode | 0x3 |
| Length | 0x6 |
| DataL_MSB | MSB of 24-bit data for low instruction word |
| DataH_MSB | MSB of 24-bit data for high instruction word |
| Addr_MSB | MSB of 24-bit destination address |
| Addr_LS | Least Significant 16 bits of 24-bit destination address |
| DataL_LS | Least Significant 16 bits of 24-bit data for low instruction word |
| DataH_LS | Least Significant 16 bits of 24-bit data for high instruction word |

The PROG2W command instructs the Programming Executive to program two instruction words of code memory (6 bytes) to the specified memory address.

After the words have been programmed to code memory, the Programming Executive verifies the programmed data against the data in the command.

Expected Response (2 words):

0x1300
0x0002

5.2.4.4 PROGP Command

| | | | | | |
|----------|----|----------|---|---|---|
| 15 | 12 | 11 | 8 | 7 | 0 |
| Opcode | | Length | | | |
| Reserved | | Addr_MSB | | | |
| Addr_LS | | | | | |
| D_1 | | | | | |
| D_2 | | | | | |
| ... | | | | | |
| D_N | | | | | |

Table 5-5 provides the description for the PROGP command.

TABLE 5-5: COMMAND DESCRIPTION

| Field | Description |
|----------|---|
| Opcode | 0x5 |
| Length | 0xC3 |
| Reserved | 0x0 |
| Addr_MSB | MSB of 24-bit destination address |
| Addr_LS | Least Significant 16 bits of 24-bit destination address |
| D_1 | 16-bit Data Word 1 |
| D_2 | 16-bit Data Word 2 |
| ... | 16-bit Data Word 3 through 191 |
| D_192 | 16-bit Data Word 192 |

The PROGP command instructs the Programming Executive to program one row of code memory (128 instruction words) to the specified memory address. Programming begins with the row address specified in the command. The destination address should be a multiple of 0x100.

The data to program the memory, located in command words, D_1 through D_192, must be arranged using the packed instruction word format illustrated in Figure 5-4.

After all data have been programmed to code memory, the Programming Executive verifies the programmed data against the data in the command.

Expected Response (2 words):

0x1500
0x0002

Note: Refer to Table 2-3 for code memory size information.

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5.2.4.5 ERASEB Command

| | | | | | |
|--------|----|----|--------|---|---|
| 15 | 12 | 11 | 8 | 7 | 0 |
| Opcode | | | Length | | |

Table 5-6 provides the description for the ERASEB command.

TABLE 5-6: COMMAND DESCRIPTION

| Field | Description |
|--------|-------------|
| Opcode | 0x7 |
| Length | 0x1 |

The ERASEB command instructs the Programming Executive to perform a Bulk Erase of the user Flash memory.

Expected Response (2 words):

0x1700
0x0002

5.2.4.6 ERASEP Command

| | | | | | |
|-----------|----|--------|----------|---|---|
| 15 | 12 | 11 | 8 | 7 | 0 |
| Opcode | | Length | | | |
| NUM_PAGES | | | Addr_MSB | | |
| Addr_LS | | | | | |

Table 5-7 provides the description for the ERASEP command.

TABLE 5-7: COMMAND DESCRIPTION

| Field | Description |
|-----------|---|
| Opcode | 0x9 |
| Length | 0x3 |
| NUM_PAGES | Up to 255 |
| Addr_MSB | Most Significant Byte of the 24-bit address |
| Addr_LS | Least Significant 16 bits of the 24-bit address |

The ERASEP command instructs the Programming Executive to Page Erase [NUM_PAGES] of code memory. The code memory must be erased at an “even” 1024 instruction word address boundary.

Expected Response (2 words):

0x1900
0x0002

5.2.4.7 QVER Command

| | | | |
|--------|----|--------|---|
| 15 | 12 | 11 | 0 |
| Opcode | | Length | |

Table 5-8 provides the description for the QVER command.

TABLE 5-8: COMMAND DESCRIPTION

| Field | Description |
|--------|-------------|
| Opcode | 0xB |
| Length | 0x1 |

The QVER command queries the version of the Programming Executive software stored in test memory. The “version.revision” information is returned in the response’s QE_Code, using a single byte with the following format: main version in upper nibble and a revision in the lower nibble (i.e., 0x23 means Version 2.3 of the Programming Executive software).

Expected Response (2 words):

0x1BMN (where “MN” stands for Version M.N)
0x0002

5.2.4.8 CRCP Command

| | | | | | |
|----------|----|--------|----------|---|---|
| 15 | 12 | 11 | 8 | 7 | 0 |
| Opcode | | Length | | | |
| Reserved | | | Addr_MSB | | |
| Addr_LSW | | | | | |
| Reserved | | | Size_MSB | | |
| Size_LSW | | | | | |

Table 5-9 provides the description for the CRCP command.

TABLE 5-9: COMMAND DESCRIPTION

| Field | Description |
|----------|---|
| Opcode | 0xC |
| Length | 0x5 |
| Addr_MSB | Most Significant Byte of 24-bit address |
| Addr_LSW | Least Significant 16 bits of 24-bit address |
| Size | Number of 24-bit locations (address range divided by two) |

The CRCP command performs a CRC-16 on the range of memory specified. This command can substitute for a full chip verify. Data are shifted in a packed method as shown in Figure 5-4, bitwise, Least Significant Byte (LSB) first.

Example:

CRC-CCITT-16 with test data of “123456789” becomes 0x29B1

Expected Response (3 words):

QE_Code: 0x1C00
Length: 0x0003
CRC Value: 0xXXXX

5.2.4.9 QBLANK Command

| | | | | | |
|----------|--|----------|----|--|---|
| 15 | | 12 | 11 | | 0 |
| Opcode | | Length | | | |
| Reserved | | Size_MSB | | | |
| Size_LSW | | | | | |
| Reserved | | Addr_MSB | | | |
| Addr_LSW | | | | | |

Table 5-10 provides the description for the QBLANK command.

TABLE 5-10: COMMAND DESCRIPTION

| Field | Description |
|----------|---|
| Opcode | 0xE |
| Length | 0x5 |
| Size | Length of program memory to check (in 24-bit words) + Addr_MS |
| Addr_MSB | Most Significant Byte of the 24-bit address |
| Addr_LSW | Least Significant 16 bits of the 24-bit address |

The QBLANK command queries the Programming Executive to determine if the contents of code memory are blank (contains all '1's). The size of code memory to check must be specified in the command.

The Blank Check for code memory begins at [Addr] and advances toward larger addresses for the specified number of instruction words.

QBLANK returns a QE_Code of 0xF0 if the specified code memory is blank; otherwise, QBLANK returns a QE_Code of 0x0F.

Expected Response (2 words for blank device):

0x1EF0
0x0002

Expected Response (2 words for non-blank device):

0x1E0F
0x0002

Note: The QBLANK command does not check the system operation Configuration bits, since these bits are not set to '1' when a Chip Erase is performed.

5.3 Programming Executive Responses

The Programming Executive sends a response to the programmer for each command that it receives. The response indicates if the command was processed correctly. It includes any required response data or error data.

The Programming Executive response set is shown in Table 5-11. This table contains the opcode, mnemonic and description for each response. The response format is described in Section 5.3.1 "Response Format".

TABLE 5-11: PROGRAMMING EXECUTIVE RESPONSE OPCODES

| Opcode | Mnemonic | Description |
|--------|----------|----------------------------------|
| 0x1 | PASS | Command successfully processed |
| 0x2 | FAIL | Command unsuccessfully processed |
| 0x3 | NACK | Command not known |

5.3.1 RESPONSE FORMAT

All Programming Executive responses have a general format, consisting of a two-word header and any required data for the command.

| | | | | | | | | |
|---------------------|--|----------|----|---------|---|---|--|---|
| 15 | | 12 | 11 | | 8 | 7 | | 0 |
| Opcode | | Last_Cmd | | QE_Code | | | | |
| Length | | | | | | | | |
| D_1 (if applicable) | | | | | | | | |
| ... | | | | | | | | |
| D_N (if applicable) | | | | | | | | |

Table 5-12 provides the description of the response format.

TABLE 5-12: RESPONSE FORMAT DESCRIPTION

| Field | Description |
|----------|---|
| Opcode | Response opcode |
| Last_Cmd | Programmer command that generated the response |
| QE_Code | Query code or error code |
| Length | Response length in 16-bit words (includes two header words) |
| D_1 | First 16-bit data word (if applicable) |
| D_N | Last 16-bit data word (if applicable) |

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5.3.1.1 Opcode Field

The opcode is a 4-bit field in the first word of the response. The opcode indicates how the command was processed (see [Table 5-11](#)). If the command was processed successfully, the response opcode is PASS. If there was an error in processing the command, the response opcode is FAIL and the QE_Code indicates the reason for the failure. If the command sent to the Programming Executive is not identified, the Programming Executive returns a NACK response.

5.3.1.2 Last_Cmd Field

The Last_Cmd is a 4-bit field in the first word of the response and indicates the command that the Programming Executive processed. Since the Programming Executive can only process one command at a time, this field is technically not required. However, it can be used to verify that the Programming Executive correctly received the command that the programmer transmitted.

5.3.1.3 QE_Code Field

The QE_Code is a byte in the first word of the response. This byte is used to return data for query commands and error codes for all other commands.

When the Programming Executive processes one of the two query commands (QBLANK or QVER), the returned opcode is always PASS and the QE_Code holds the query response data. The format of the QE_Code for both queries is shown in [Table 5-13](#).

TABLE 5-13: QE_Code FOR QUERIES

| Query | QE_Code |
|--------|---|
| QBLANK | 0x0F = Code memory is NOT blank 0xF0 = Code memory is blank |
| QVER | 0xMN, where Programming Executive Software Version = M.N (i.e., 0x32 means Software Version 3.2) |

When the Programming Executive processes any command other than a query, the QE_Code represents an error code. Supported error codes are shown in [Table 5-14](#). If a command is successfully processed, the returned QE_Code is set to 0x0, which indicates that there is no error in the command processing. If the verify of the programming for the PROGW command fails, the QE_Code is set to 0x1. For all other Programming Executive errors, the QE_Code is 0x2.

TABLE 5-14: QE_Code FOR NON-QUERY COMMANDS

| QE_Code | Description |
|---------|---------------|
| 0x0 | No error |
| 0x1 | Verify failed |
| 0x2 | Other error |

5.3.1.4 Response Length

The response length indicates the length of the Programming Executive's response in 16-bit words. This field includes the two words of the response header.

With the exception of the response for the read commands, the length of each response is only two words.

The response to the READP commands uses the packed instruction word format described in [Section 5.2.2 "Packed Data Format"](#). When reading an odd number of Program Memory Words (N odd), the response to the READP command is $(3 * (N + 1)/2 + 2)$ words. When reading an even number of Program Memory Words (N even), the response to the READP command is $(3 * N/2 + 2)$ words.

5.4 Programming the Programming Executive to Memory

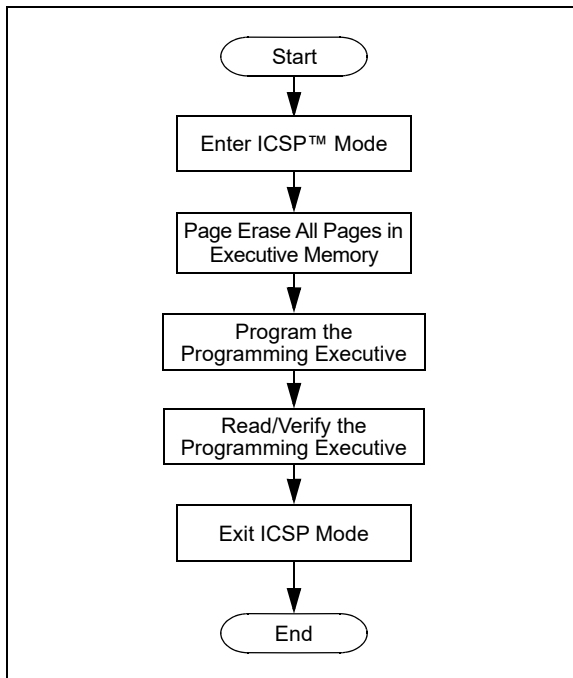
Note: The Programming Executive can be obtained from each device page on the Microchip website: www.microchip.com.

5.4.1 OVERVIEW

If it is determined that the Programming Executive is not present in the executive memory (as described in [Section 4.2 “Confirming the Presence of the Programming Executive”](#)), the Programming Executive must be programmed to executive memory.

[Figure 5-5](#) shows the high-level process of programming the Programming Executive into executive memory. First, the ICSP mode must be entered and the executive memory must be erased. Then, the Programming Executive is programmed and verified. Finally, ICSP mode is exited.

FIGURE 5-5: HIGH-LEVEL PROGRAMMING EXECUTIVE PROGRAM FLOW



5.4.2 ERASING EXECUTIVE MEMORY

The procedure for erasing each page of executive memory is similar to that of erasing program memory and is shown in [Figure 3-6](#). It consists of setting NVMCON to 0x4003 and then executing the programming cycle.

[Table 5-15](#) shows the ICSP programming process for erasing the executive code memory.

Note: The Programming Executive memory must always be erased before it is programmed, as described in [Figure 5-5](#).

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TABLE 5-15: SERIAL INSTRUCTION EXECUTION FOR ERASING ALL PAGES OF EXECUTIVE MEMORY

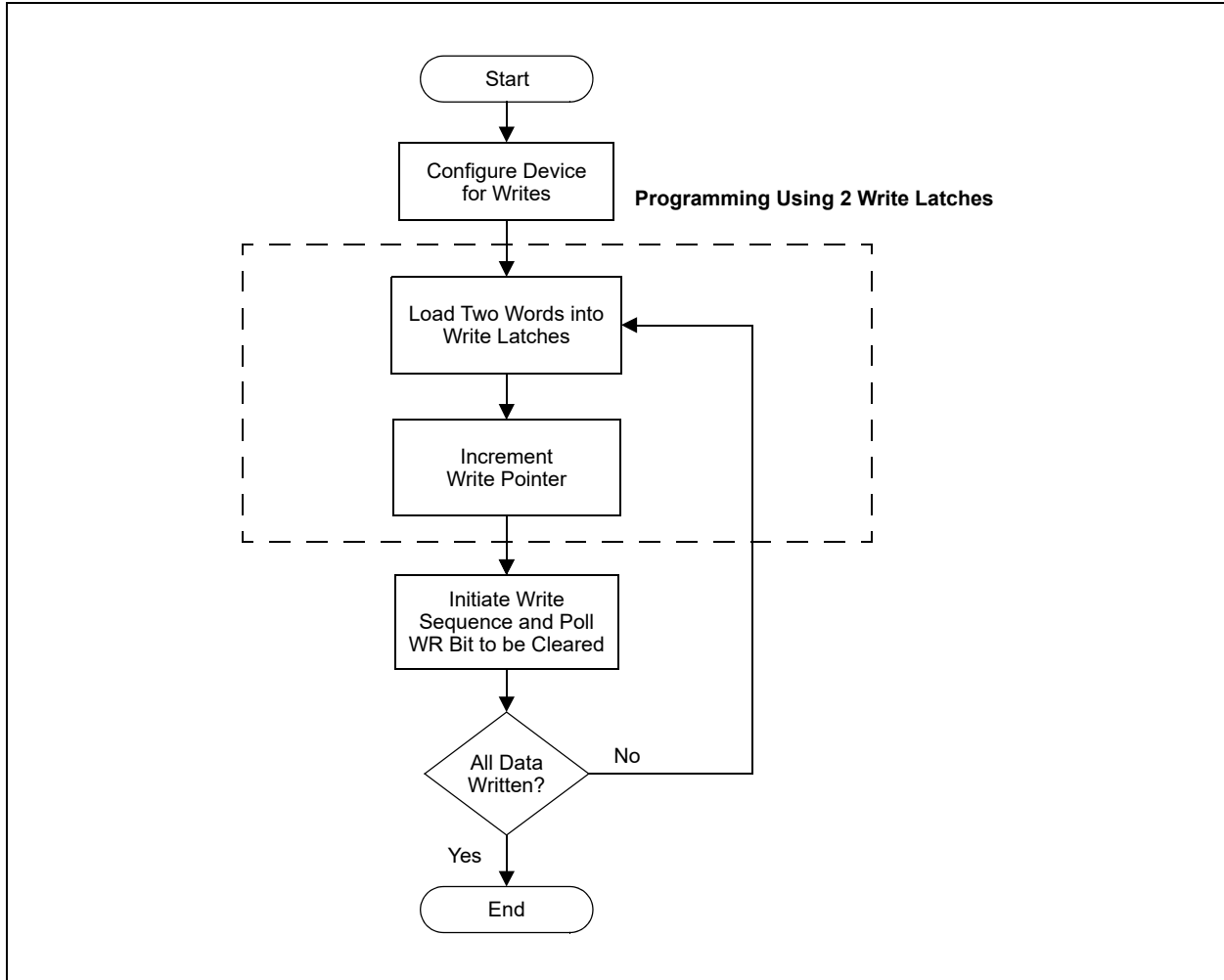
| Command (Binary) | Data (Hex) | Description |
|--|------------|--|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Set the NVMADRU/NVMADR register pair to point to the correct page of executive memory to be erased. | | |
| 0000 | 2xxxx3 | MOV #DestinationAddress<15:0>, W3 |
| 0000 | 2xxxx4 | MOV #DestinationAddress<23:16>, W4 |
| 0000 | 884693 | MOV W3, NVMADR |
| 0000 | 8846A4 | MOV W4, NVMADRU |
| Step 3: Set the NVMCON register to erase the first page of executive memory. | | |
| 0000 | 24003A | MOV #0x4003, W10 |
| 0000 | 88468A | MOV W10, NVMCON |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 4: Initiate the erase cycle. | | |
| 0000 | 200551 | MOV #0x55, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | 200AA1 | MOV #0xAA, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | A8F1A1 | BSET NVMCON, #WR |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 5: Generate clock pulses for the Page Erase operation to complete until the WR bit is clear. | | |
| 0000 | 000000 | NOP |
| 0000 | 804680 | MOV NVMCON, W0 |
| 0000 | 000000 | NOP |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| — | — | Repeat until the WR bit is clear. |
| Step 6: Repeat Steps 2-5 for all pages of executive memory. | | |

5.4.3 PROGRAM THE PROGRAMMING EXECUTIVE

Storing the PE to executive memory is similar to normal programming of code memory. The executive memory must first be erased and then programmed using two-word writes (two instruction words). The control flow for this method is summarized in [Figure 5-6](#).

[Table 5-16](#) provides the ICSP programming processes for PE memory. To minimize programming time, the same packed data format that the PE uses is utilized. See [Section 5.2 “Programming Executive Commands”](#) for more details on the packed data format.

FIGURE 5-6: PROGRAMMING EXECUTIVE PROGRAM FLOW



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**TABLE 5-16: PROGRAMMING THE PROGRAMMING EXECUTIVE
(TWO-WORD LATCH WRITES)**

| Command (Binary) | Data (Hex) | Description |
|--|---------------|------------------------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register for writing to the latches. | | |
| 0000 | 200FAC | MOV #0xFA, W12 |
| 0000 | 8802AC | MOV W12, TBLPAG |
| Step 3: Load W0:W2 with the next two packed instruction words to program. | | |
| 0000 | 2xxxx0 | MOV #<LSW0>, W0 |
| 0000 | 2xxxx1 | MOV #<MSB1:MSB0>, W1 |
| 0000 | 2xxxx2 | MOV #<LSW1>, W2 |
| Step 4: Set the Read Pointer (W6) and the Write Pointer (W7), and load the write latches. | | |
| 0000 | EB0300 | CLR W6 |
| 0000 | 000000 | NOP |
| 0000 | EB0380 | CLR W7 |
| 0000 | 000000 | NOP |
| 0000 | BB0BB6 | TBLWTL [W6++], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BBDBB6 | TBLWTH.B [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BBEBB6 | TBLWTH.B [W6++], [++W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BB0B96 | TBLWTL.W [W6], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 5: Set the NVMADRU/NVMADR register pair to point to the correct row. | | |
| 0000 | 2xxxx3 | MOV #DestinationAddress<15:0>, W3 |
| 0000 | 2xxxx4 | MOV #DestinationAddress<23:16>, W4 |
| 0000 | 884693 | MOV W3, NVMADR |
| 0000 | 8846A4 | MOV W4, NVMADRU |
| Step 6: Set the NVMCON register to program two instruction words. | | |
| 0000 | 24001A | MOV #0x4001, W10 |
| 0000 | 000000 | NOP |
| 0000 | 88468A | MOV W10, NVMCON |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |

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**TABLE 5-16: PROGRAMMING THE PROGRAMMING EXECUTIVE
(TWO-WORD LATCH WRITES) (CONTINUED)**

| Command (Binary) | Data (Hex) | Description |
|--|---------------|--|
| Step 7: Initiate the write cycle. | | |
| 0000 | 200551 | MOV #0x55, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | 200AA1 | MOV #0xAA, W1 |
| 0000 | 8846B1 | MOV W1, NVMKEY |
| 0000 | A8F1A1 | BSET NVMCON, #WR |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 8: Wait for program operation to complete and make sure the WR bit is clear. | | |
| 0000 | 000000 | NOP |
| 0000 | 804680 | MOV NVMCON, W0 |
| 0000 | 000000 | NOP |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| — | — | Repeat until the WR bit is clear. |
| Step 9: Repeat Steps 3-8 until all code memory is programmed. | | |

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5.4.4 READING EXECUTIVE MEMORY

Reading from executive memory is performed by executing a series of TBLRD instructions and clocking out the data using the REGOUT command.

Table 5-17 provides the ICSP programming details for reading executive memory.

To minimize reading time, the same packed data format that the PE uses is utilized. See Section 5.2 “Programming Executive Commands” for more details on the packed data format.

TABLE 5-17: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY

| Command (Binary) | Data (Hex) | Description |
|--|------------|-------------------------------|
| Step 1: Exit the Reset vector. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 2: Initialize the TBLPAG register and the Read Pointer (W6) for the TBLRD instruction. | | |
| 0000 | 200xx0 | MOV #<SourceAddress23:16>, W0 |
| 0000 | 8802A0 | MOV W0, TBLPAG |
| 0000 | 2xxxx6 | MOV #<SourceAddress15:0>, W6 |

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TABLE 5-17: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY (CONTINUED)

| Command (Binary) | Data (Hex) | Description |
|---|---------------|-------------------------|
| Step 3: Initialize the Write Pointer (W7) and store the next four locations of code memory to W0:W5. | | |
| 0000 | EB0380 | CLR W7 |
| 0000 | 000000 | NOP |
| 0000 | BA1B96 | TBLRDL [W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBB6 | TBLRDH.B [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBD6 | TBLRDH.B [++W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BA1BB6 | TBLRDL [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BA1B96 | TBLRDL [W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBB6 | TBLRDH.B [W6++], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BADBD6 | TBLRDH.B [++W6], [W7++] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | BA0BB6 | TBLRDL [W6++], [W7] |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |

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TABLE 5-17: SERIAL INSTRUCTION EXECUTION FOR READING CODE MEMORY (CONTINUED)

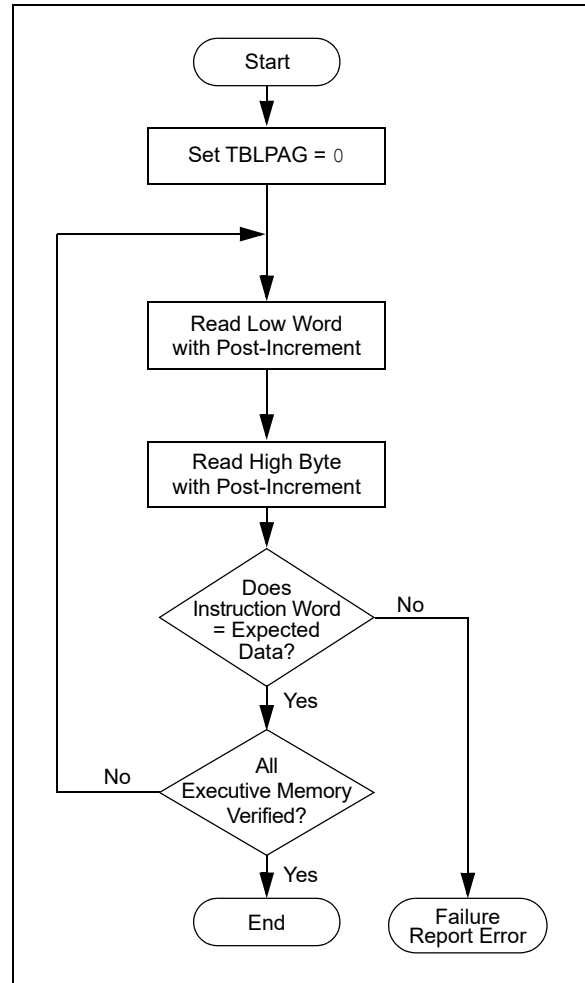
| Command (Binary) | Data (Hex) | Description |
|---|------------|--|
| Step 4: Output W0:W5 using the VISI register and REGOUT command. | | |
| 0000 | 887E60 | MOV W0, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E61 | MOV W1, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E62 | MOV W2, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E63 | MOV W3, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E64 | MOV W4, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| 0000 | 887E65 | MOV W5, VISI |
| 0000 | 000000 | NOP |
| 0001 | <VISI> | Clock out contents of the VISI register. |
| 0000 | 000000 | NOP |
| Step 5: Reset the device's internal PC. | | |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 040200 | GOTO 0x200 |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| 0000 | 000000 | NOP |
| Step 6: Repeat Steps 3-5 until all desired code memory is read. | | |

5.4.5 PROGRAMMING VERIFICATION

The verification is performed by reading back the executive memory space and comparing it against the copy held in the programmer's buffer.

The verification process is shown in [Figure 5-7](#). The lower word of the instruction is read, and then the lower byte of the upper word is read and compared against the instruction stored in the programmer's buffer. Refer to [Section 5.4.4 "Reading Executive Memory"](#) for implementation details of reading executive memory.

FIGURE 5-7: VERIFY PROGRAMMING EXECUTIVE MEMORY FLOW



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6.0 DEVICE ID/UNIQUE ID

The Device ID region of memory can be used to determine variant and manufacturing information about the chip. This region of memory is read-only and can be read when code protection is enabled.

Table 6-1 lists the identification information for the device. Table 6-2 shows the Device ID registers. Table 6-3 shows that the JTAG ID is generated given a device DEVREV and DEVID.

TABLE 6-1: DEVICE ID

| Device | DEVID |
|--------------------|--------|
| dsPIC33CDVL64MC106 | 0x991A |
| dsPIC33CDV64MC106 | 0x991B |

TABLE 6-2: DEVICE ID REGISTERS

| Address | Name | Bit | | | | | | | | | | | | | | | |
|----------|--------|--------------|----|----|----|----|----|---|---|---|---|---|---|---|---|---|---|
| | | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| 0xFF0000 | DEVID | DEVID Value | | | | | | | | | | | | | | | |
| 0x8016FC | DEVREV | DEVREV Value | | | | | | | | | | | | | | | |

TABLE 6-3: JTAG ID

| 31 | 28 | 27 | 12 | 11 | 0 | | |
|-------------|----|----|----|-------------|---|-------------------------|--|
| DEVREV[3:0] | | | | DEVID[15:0] | | Manufacturer ID (0x053) | |
| 4 bits | | | | 16 bits | | 12 bits | |

6.1 Unique Device ID (UDID)

The dsPIC33CDVL64MC106 device is individually encoded during final manufacturing with a Unique Device Identifier or UDID. The UDID cannot be erased by a Bulk Erase command or any other user-accessible means. This feature allows for manufacturing traceability of Microchip Technology devices in applications where this is a requirement. It may also be used by the application manufacturer for any number of things that may require unique identification, such as:

- Tracking the device
- Unique serial number
- Unique security key

The UDID comprises five 24-bit program words. When taken together, these fields form a unique 120-bit identifier. The UDID is stored in five read-only locations in the device configuration space. Table 6-4 lists the addresses of the Identifier Words and shows their contents.

TABLE 6-4: UDID ADDRESSES

| UDID | Address | Content |
|-------|----------|-------------|
| UDID1 | 0x801200 | UDID Word 1 |
| UDID2 | 0x801202 | UDID Word 2 |
| UDID3 | 0x801204 | UDID Word 3 |
| UDID4 | 0x801206 | UDID Word 4 |
| UDID5 | 0x801208 | UDID Word 5 |

7.0 CRC CHECKSUM COMPUTATION

Unlike older PIC® devices, the Microchip toolchain runs a 32-bit CRC calculation on the entire Hex file to calculate its checksum. The checksum uses the standard CRC-32 algorithm with the polynomial, 0x4C11DB7.

EXAMPLE 7-1: CRC CHECKSUM COMPUTATION

$$(x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1)$$

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8.0 AC/DC CHARACTERISTICS AND TIMING REQUIREMENTS

Table 8-1 lists the AC/DC characteristics and timing requirements.

TABLE 8-1: AC/DC CHARACTERISTICS AND TIMING REQUIREMENTS

| Standard Operating Conditions Operating Temperature: -40°C to +125°C. Programming at +25°C is recommended. | | | | | | |
|---|--------|---|---------|---------|-------|-----------------------------------|
| Param. No. | Symbol | Characteristic | Min. | Max. | Units | Conditions |
| D111 | VDD | Supply Voltage During Programming | 3.0 | 3.60 | V | Normal programming ⁽¹⁾ |
| D113 | IDDP | Supply Current During Programming | — | 10 | mA | |
| D114 | IPEAK | Instantaneous Peak Current During Start-up | — | 200 | mA | |
| D031 | VIL | Input Low Voltage | VSS | 0.2 VDD | V | |
| D041 | VIH | Input High Voltage | 0.7 VDD | VDD | V | |
| D080 | VOL | Output Low Voltage | — | 0.4 | V | |
| D090 | VOH | Output High Voltage | 2.4 | — | V | |
| D012 | CIO | Capacitive Loading on I/O Pin (PGEDx) | — | 50 | pF | |
| P1 | TPGC | Serial Clock (PGECx) Period (ICSP™) | 200 | — | ns | |
| P1 | TPGC | Serial Clock (PGECx) Period (Enhanced ICSP) | 500 | — | ns | |
| P1A | TPGCL | Serial Clock (PGECx) Low Time (ICSP) | 80 | — | ns | |
| P1A | TPGCL | Serial Clock (PGECx) Low Time (Enhanced ICSP) | 200 | — | ns | |
| P1B | TPGCH | Serial Clock (PGECx) High Time (ICSP) | 80 | — | ns | |
| P1B | TPGCH | Serial Clock (PGECx) High Time (Enhanced ICSP) | 200 | — | ns | |
| P2 | TSET1 | Input Data Setup Time to Serial Clock ↓ | 15 | — | ns | |
| P3 | THLD1 | Input Data Hold Time from PGECx ↓ | 15 | — | ns | |
| P4 | TDLY1 | Delay Between 4-Bit Command and Command Operand | 40 | — | ns | |
| P4A | TDLY1A | Delay Between Command Operand and Next 4-Bit Command | 40 | — | ns | |
| P5 | TDLY2 | Delay Between Last PGECx ↓ of Command to First PGECx ↑ of Read of Data Word | 20 | — | ns | |
| P6 | TSET2 | VDD ↑ Setup Time to MCLR ↑ | 100 | — | ns | |
| P7 | THLD2 | Input Data Hold Time from MCLR ↑ | 50 | — | ms | |
| P8 | TDLY3 | Delay Between Last PGECx ↓ of Command Byte to PGEDx ↑ by PE | 12 | — | μs | |

Note 1: VDD must also be supplied to the AVDD pins during programming. AVDD and AVSS should always be within ±0.3V of VDD and VSS, respectively.

2: Time depends on the FRC accuracy and the value of the FRC Oscillator Tuning register. Refer to the “**Electrical Characteristics**” chapter in the specific device data sheet.

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TABLE 8-1: AC/DC CHARACTERISTICS AND TIMING REQUIREMENTS (CONTINUED)

| Standard Operating Conditions | | | | | | |
|--|--------|---|------|------|-------|----------------------------|
| Operating Temperature: -40°C to +125°C. Programming at +25°C is recommended. | | | | | | |
| Param. No. | Symbol | Characteristic | Min. | Max. | Units | Conditions |
| P9A | TDLY4 | Programming Executive Command Processing Time | 10 | — | μs | |
| P9B | TDLY5 | Delay Between PGEDx ↓ by Programming Executive to PGEDx Released by Programming Executive | 15 | 23 | μs | |
| P10 | TDLY6 | PGECx Low Time After Programming | 400 | — | ns | |
| P11 | TDLY7 | Bulk Erase Time | — | 16.0 | ms | See Note 2 |
| P12 | TDLY8 | Page Erase Time | — | 4.2 | ms | See Note 2 |
| P13 | TDLY9 | Double-Word Write Time | — | 34.5 | μs | See Note 2 |
| | | Row Write Time | | 1.1 | ms | See Note 2 |
| P14 | TR | MCLR Rise Time to Enter ICSP mode | — | 1.0 | μs | |
| P15 | TVALID | Data Out Valid from PGECx ↑ | 10 | — | ns | |
| P16 | TDLY10 | Delay Between Last PGECx ↓ and MCLR ↓ | 0 | — | s | |
| P17 | THLD3 | MCLR ↓ to VDD ↓ | 100 | — | ns | |
| P18 | TKEY1 | Delay from First MCLR ↓ to First PGECx ↑ for Key Sequence on PGEDx | 1 | — | ms | |
| P19 | TKEY2 | Delay from Last PGECx ↓ for Key Sequence on PGEDx to Second MCLR ↑ | 25 | — | ns | |
| P21 | TMCLRH | MCLR High Time | — | 500 | μs | |

Note 1: VDD must also be supplied to the AVDD pins during programming. AVDD and AVSS should always be within ±0.3V of VDD and VSS, respectively.

2: Time depends on the FRC accuracy and the value of the FRC Oscillator Tuning register. Refer to the “**Electrical Characteristics**” chapter in the specific device data sheet.

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APPENDIX A: REVISION HISTORY

Revision A (September 2020)

Original version of the programming specification created for the dsPIC33CDVL64MC106 device.

Revision B (August 2021)

Added the dsPIC33CDV64MC106 device.

Updated [Figure 2-3](#) and added [Figure 2-4](#).

Updated [Table 2-3](#) and [Table 6-1](#).

Revision C (January 2023)

Updated [Figure 2-5](#).

Updated [Table 6-2](#).

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